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COMP 211, Spring 2010

Recap of Previous Lecture

Data-directed design

- Start with *data definition*
 - Specifies structure of data
- Derive *function template* from data definition
 - Model for any function that can be performed on data
 - Use generic name (e.g., f) for function
- Create template instantiation for a specific function and primary argument
 - Use specific name (e.g., sort) for function
 - Define separate auxiliary functions for other arguments if needed
- Develop code based on template instantiation

Today's Goals

- Loose Ends
 - Catching mistakes and raising errors
 - and & or Operations
- Trees
 - Significantly more expressive type
 - "Lists with many tails"
- Examples:
 - Family tree
 - Binary search tree

```
Using and & or
  Scheme and abbreviates a conditional and takes an arbitrary number
  of arguments (our student dialects require at least 2)
  (and arg1 arg2 ... argn) abbreviates
  (cond [(not arg1) false]
         [(not arg2) false]
         [else argn])
  Hence,
•
  (and true true false (zero? (/ 2 0)) ...) => false
  Called "short-circuit" or "non-strict" semantics for and

    What does or do? It is the dual of and.

  (or arg1 arg2 ... argn) abbreviates
  (cond [arg1 true]
         [arg2 true]
         . . .
         [else argn])
```

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and & or cont.

- What are the reduction rules (laws) for and?
 - (and false ... argn) => false
 - (and true arg2 ... argn) => (and arg2 ... argn)
 - (and v) => v
- What are the reduction rules (laws) for or?
 - (or true ... argn) => true
 - (or false arg2 ... argn) => (or arg2 ... argn)
 - (or v) => v

Error Reporting

To report an error in (a student dialect of) Scheme invoke:

(error msg)

where msg is a string enclosed in quotation marks. In full Scheme, error takes additional arguments. See Help Desk.

HTDP (our online book) describes an *obsolete* version of error.

Semantics: the entire computation is aborted and an error message **msg** is printed.

```
Example:
(define (len aloa
  (cond
    [(empty? Aloa) 0]
    [(cons? aloa) (add1 (len (rest aloa)))]
    [else (error "len: expects argument of type <list>")]))
```

Error Reporting (cont.)

Questions:

- Is error reporting a good idea.
- Should error behavior be documented?
- Answers to questions are surprising subtle and lacking in consensus. In the case above, it is probably a good idea but it is often not done because it clutters the code and adds overhead. Moreover, the error is caught anyway with a slightly less informative diagnostic. On the other hand, DrScheme libraries do perform such checks, partly because Scheme does not perform static type checking.
- Error checking should not be included in a contract (purpose) unless the client code can depend on it and use it (by "catching" the error). We will cover error catching in Java.

Reductions for Errors

 First, note how errors work for functions you already know. In any context, erroneous primitive function applications like (/ 1 0) abort the computation and return an error at the top level:

(/ 1 0) => /: division by zero

This aborting behavior is unique among our rules.

 The error construct gives program text access to this mechanism. In any context

(error "append: expects <list> as first argument") =>

append: expects <list> as first argument [at top level]

Use errors only as required by the problem or recipe.

Another Inductive Type: Trees (Structures in Structures)

- Labeled trees
- Organizational charts
- Decision trees
- Search trees and many more!



Approved by



From Lists to Trees

Example of a List Data Definition

- ; A list-of-symbols is
- ; empty, or
- ; (cons s los)
- ; where s is a symbol and los is a list-of-symbols

A list has one embedded structure (rest)

(Ancestor) Family Trees

Example of a Family Tree Data Definition

- ; A child is
- ; empty // Represents "unknown"
- ; (make-child n m f) // Two self-references
- ; where n is a symbol, m is a child and f is a child

(define-struct child (name mother father))

A child has two embedded structures (mother, father) which belong to type child. Perhaps child is misnamed.

Tree Depth (in class ex.)

- Consider the following problem
 - Given an ancestry tree, compute the maximum number of generations for which we know something about this person.
- Contract (or "type") is
 - child -> natural
- Examples (next slide)

Tree Depth Examples

(define cat (make-child 'Cat empty empty)) (define tom (make-child 'Tom cat empty)) (define jane (make-child empty tom)) (define johnny (make-child 'Johnny empty empty)) (define ray (make-child 'Ray empty johnny)) (define sue (make-child 'Ray empty ray)) (define rob (make-child 'Rob empty sue)) (define bob (make-child 'Bob jane rob))

(check-expect (max-depth cat) 1) (check-expect (max-depth tom) 2) (check-expect (max-depth jane) 3) (check-expect (max-depth johnny) 1) (check-expect (max-depth ray) 2) (check-expect (max-depth sue) 3) (check-expect (max-depth rob) 4) (check-expect (max-depth bob) 5)

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Tree Depth Template Instantiation

```
; Template Instantiation (trivial)
(define (max-depth c)
  (cond
    [(empty? c) ...]
    [else ...
    ... (max-depth (child-mother c)) ...
    ... (max-depth (child-father c)) ...)))
```

Tree Depth Code

; Code (define (max-depth c) (cond [(empty? c) 0] [else (add1 (max (max-depth (child-mother c)) (max-depth (child-father c)))]))

Examples can help in writing code. Work through simple examples by hand.

Binary Trees and Binary Search Trees

- ; A binary tree (BT) is either
- ; empty, or
- (make-node n l r)
- ; where n is a number, 1 and r are BTs.

(define-struct node (num left right))

- ; A binary search tree (BST) is a binary tree where
- ; the following invariants hold:
- ; 1. Numbers in 1 are less than or equal to n
- ; 2. Numbers in r are greater than n

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Binary Search Trees

Which binary tree satisfies the invariants of a binary search tree?





