

Corky Cartwright
Vivek Sarkar
Department of Computer Science
Rice University



Outline

- Vectors in Scheme
- Functional vs. Imperative views of
 - Iteration
 - Arrays
- Today's lecture is all bonus material!
 - Will not be covered in test or homework



A First Look at Vectors (Section 29.3)

Goal: array-like data structure with O(1) lookup time for a given index

- Vector creation
 - (vector V-0 ... V-n) creates a vector with n+1 elements, V-0 through V-n
 - (build-vector n f) creates a vector with n elements, (f 0) through (f (- n 1))
 - · Simple case of an array comprehension



Vector Operations (contd)

- (vector-length V) returns the number of items in vector V
 - Results in an error if V is not a vector
- (vector-ref V i) returns the ith item in vector V
 - Results in an error if V is not a vector or i is not a number or i < 0 or i >= (vector-length V)
- (vector? V) returns true if V is a vector

Simple example: sum-of-3

- Example: (vector-sum-of-3 (vector 2 4 6 8 10))
- NOTE: vector is like cons, and vector-ref is like first/rest

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Binary Search on a Sorted Vector of Numbers

```
;; bin-srch: asvon number number number -> number
  For input vector V, value X, lower & upper bounds
   lo and hi, return index i in lo ... hi such that
  (\text{vector-ref V i}) = X, else return -1 if X not found
;; NOTE: use Advanced Student setting to use vectors
(define (bin-srch V X lo hi)
  (let ((mid (floor (/ (+ lo hi) 2))))
       (cond
         ((> lo hi) -1)
         ((= (vector-ref V mid) X) mid)
         ((> (vector-ref V mid) X) (bin-srch V X lo(- mid 1)))
         ((< (vector-ref V mid) X) (bin-srch V X (+ mid 1) hi)))</pre>
  ))
```



- What is the execution time complexity of binary search using a vector?
- How would the complexity of binary search change if we replaced the vector by a list of pairs (and used list-ref instead of vector-ref)?



- Pro: vector-ref can be used to access any element in a vector in O(1) time
 - Multiple first/rest operations may be needed to traverse a list
- Con: extending a vector or extracting from a vector takes O(n) time
 - Constructing a list with a new element at the start of an existing list takes O(1) time (cons)
 - Extracting the tail of a list takes O(1) time (rest)



Iteration

- Iterating over a vector/list in a functional language is usually accomplished by (tail) recursion
- Iterating over a vector/list in an imperative language is usually accomplished by iteration
 - e.g., while-loops and for-loops in Java
- Does this mean that iteration is inherently non-functional?

Functional Language with Iteration

- Sisal stands for Streams and Iteration in a Single Assignment Language
- Defined in 1983, revised and frozen in 1985
- Original collaborators were LLNL, Colorado State U, University of Manchester, and DEC
 - Used for research at many other institutions, including Stanford University
- Language design strongly influenced by dataflow computation model

Sisal Objectives

- to define a general-purpose functional language
- to define a language independent intermediate form for dataflow graphs
- to develop optimization techniques for high performance parallel applicative computing
- to develop a microtasking environment that supports dataflow on conventional computer systems
- to achieve execution performance comparable to imperative languages
- to validate the functional style of programming for large-scale scientific applications



% Hello world!

end for

end function

```
define main
function main(returns array[character])
 "hello world"
end function
% Simple arrays
define main
function main(A: array[integer] returns integer, array[integer])
 for element in A
                   % parallel loop with independent iterations
   sqr := element * element
returns
   value of sum sqr % reduce operation
                    % array comprehension, like vector-build
   array of sqr
```

Sequential iteration with for-initial loop expressions

- Not all loops are implicitly data parallel
- Sisal supports an iterative form that supports the idea of "loop carried dependencies"
- The loop body is allowed to reference both the "new" and the "old" value of a definition (variable)
- An separate body defines the initial values

Example #1: Iterated Function Composition

```
for initial % Initializer body is like the zeroth iteration
    i := 0;
     accum := 0;
while i \le n repeat
    i := old i + 1;
                          % Note the use of "old" to denote previous value
    accum := f(old accum);
returns
    value of accum
end for
Scheme equivalent:
(local
 ((define (g i accum)
  (cond [(i < n) (g (+ i 1) (f accum))] [else accum])))
  (g 0 0)
```

stencil #2: 3-point stencil w/ Array Replace Operation

for initial

```
A := some_value(); % This is the zeroth "iteration"
i := array_liml(A); % Lower bound of array A's indices
while i < array_limh(A) repeat
i := old i + 1;
A := old A[i: (old A[i-1] + old A[i] + old A[i+1]) / 3.0 ];
returns
value of A
end for
```

- Array replace operation --- A[i:X] returns a new array A' identical to A, except that element I is replaced by X
 - Functional alternative to A[i] = X; in Java or C
- Semantically, A' is a copy of A, but implementations try to make their best effort to eliminate as many copies as possible.



Announcements

 Midterm to be distributed on Friday (Feb 19th)