

# Comp 311

# Functional Programming

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# Homework 1

- Please submit your homework via the **turnin** system, in a folder named **hw\_1**
- The specific files to submit are defined in the assignments
- For each section, please turn in only your final program resulting from completion of the section
- Think about overflow!

Please Restrict Your Homework  
Submission to Features Covered in Class

# Please Restrict Your Homework Submission to Features Covered in Class

- These should be the only import statements in your program:

```
import junit.framework.TestCase
```

```
import junit.framework.Assert._
```

# Type Checking

- So far, we have been rigorous about computation of programs, but we have relied on intuition for static type checking
- We can provide a *static semantics* for Core Scala along with our *dynamic semantics*

# The Substitution Model of Type Checking

- To type check a value **v**, replace **v** with its value type

`1.003`  $\Rightarrow$  `Double`

- To type check a constant **c**, reduce the defining expression of **c** to a static type **T**, then replace all occurrences of **c** with **T**

`pi = 3.14`  $\Rightarrow$

`pi : Double`

`pi * radius * radius`  $\Rightarrow$

`Double * radius * radius`

# The Substitution Model of Type Checking

- To type check a function definition:
  - Type check the body of the definition, replacing all occurrences of each parameter with the corresponding parameter type
- To type check the occurrence of a function name:
  - Replace the name with an *arrow type*, where the parameter types of the function are to the left of the arrow and the return type is to the right

`square(x: Double): Double = x * x`

`square(3.14) ⇒`

`(Double → Double)(3.14)`

# The Substitution Model of Type Checking

- To type check the application of a function to arguments:
  - Reduce the function to an arrow type
  - Reduce the arguments, left to right, to static types
  - If the argument types match the corresponding parameter types, reduce the application to the return type

# The Substitution Model of Type Checking

`square(3.14) ⇒`

`(Double → Double)(3.14) ⇒`

`(Double → Double)(Double) ⇒`

`Double`

# Methods and Operators

# Syntactic Sugar For Binary Methods

- We refer to methods that take one parameter (in addition to the receiver) as *binary methods*

```
case class Coordinate(x: Int, y: Int) {  
  def magnitude() = x*x + y*y  
  
  def add(that: Coordinate) =  
    Coordinate(x + that.x, y + that.y)  
}
```

# Syntactic Sugar For Binary Methods

```
Coordinate(1,2).add(Coordinate(3,4))
```

↳

```
Coordinate(4,6)
```

# Syntactic Sugar For Binary Methods

- With binary methods, we can elide the dot in a method call
- We can also elide the enclosing parentheses around the sole argument

# Syntactic Sugar For Binary Methods

Coordinate(1,2) add Coordinate(3,4)

↳

Coordinate(4,6)

# Operator Symbols

- Scala allows the use of operator symbols in method names
- In fact, operators are simply methods in Scala

`1.+(2) → 3`

# Coordinates Revisited

```
case class Coordinate(x: Int, y: Int) {  
  def magnitude() = x*x + y*y  
  
  def +(that: Coordinate) =  
    Coordinate(x + that.x, y + that.y)  
}
```

# Coordinates Revisited

Coordinate(1,2) + Coordinate(3,4)

↳

Coordinate(4,6)

# Requires Clauses on Class Constructors

```
case class Name(field1: Type1, ..., fieldN: TypeN)  
  require (boolean-expression)
```

- Checked on every constructor call
- Because case class instances are immutable, this ensures the property holds for the lifetime of an instance

# Equals on Case Classes

- The equals method on a case class instance checks for structural equality with its argument:

```
Rational(4,6).equals(Rational(4,6)) ↪
```

```
true
```

# Equals on Case Classes

- Note that equals is a binary method, and so we can also write this expression as:

`Rational(4,6) equals Rational(4,6) ⇨`

`true`

# Equals on Case Classes

- Of course, the built in equals method does not check for mathematical equality:

`Rational(4,6) equals Rational(2,3) ↪`

`false`

# Equals on Case Classes

- Why is this definition of equality acceptable on case classes?
- What other definition is available to us?

`Rational(4,6) equals Rational(2,3) ↪`

`false`

# Short-Circuiting And and Or Operators

- Just as we have defined a short-circuiting if-then-else operator, we can define short-circuiting and/or operators:

&&      ||

- How do we define the static and dynamic semantics of these operators?
- When are they useful?

# Calling and Defining Parameterless Methods Without Parentheses

```
def toString() = { ... }
```

vs.

```
def toString = { ... }
```

# Calling and Defining Parameterless Methods Without Parentheses

`Rational(4,6).toString()`

vs.

`Rational(4,6).toString`

# The Uniform Access Principle

- Client code should not be affected by whether an attribute is defined as a field or a method
- Only applies to immutable methods
- Can be strange even for some immutable methods (what are some examples?)

# Abstract Datatypes

# Abstract Datatypes

- Often, we wish to abstract over a collection of compound datatypes that share common properties
- For example, we might wish to define an abstract datatype for shapes, with separate case classes for each of several shapes
- For this purpose, we define an *abstract class* and use *subclassing*

# Abstract Datatypes

```
abstract class Shape  
case class Circle(radius: Double) extends Shape  
case class Square(side: Double) extends Shape  
case class Rectangle(height: Double, width: Double) extends Shape
```

# Recall Our Design Recipe

- **Analysis:** What are the objects in the problem domain? What data types we will use to represent them?
- **Contract:** What is name of our functions and their parameters? What are the requirements of the data they consume and produce? What is the meaning of what our program computes?
- **Repeat** until we are confident in our program's correctness
  - Write some **tests**
  - Sketch a function **template**
  - **Define** the function

# Recall Our Design Recipe

- **Analysis:** This is the stage where we would discover we wish to model our problem domain with functions over an abstract datatype
- **Contract:** What contract holds for each function? Do additional constraints and assurances hold for specific subclasses?
- **Repeat** until we are confident in our program's correctness
  - Write some **tests:** Same as before
  - Sketch a function **template:** This needs re-examination
  - **Define** the function

# The Design Recipe for Abstract Datatypes

- Our Function Template for computing with abstract datatypes depends on answering the following questions:
  - Do I expect to eventually add more subclasses?
  - Do I expect to eventually add more functions?

# **Case 1**

We Expect Few New Functions  
But Many New Variants

# Case 1: We Expect Few New Functions But Many New Variants

- This is a case that object-oriented programming handles well
- Classic example domains: GUI Programming, Productivity Apps, Graphics, Games
- Declare an abstract method in our superclass and provide a concrete definition for each sub-class

*a.k.a.,*

*The Union Pattern (for the datatype definitions)*

*The Template Method Pattern (for the function definitions)*

# Abstract Datatypes

```
abstract class Shape {  
    def area(): Double  
}
```

# Abstract Datatypes

```
case class Circle(radius: Double) extends Shape {  
  val pi = 3.14  
  
  def area() = pi * radius * radius  
  
}
```

# Abstract Datatypes

```
case class Square(side: Double) extends Shape {  
  def area() = side * side  
}
```

# Abstract Datatypes

```
case class Rectangle(length: Double, width: Double)
extends Shape {
  def area() = length * width
}
```

# How Do Abstract Classes Affect Our Type Checking Rules?

- When type checking a class definition, ensure that all abstract methods declared in the superclass are actually defined, with *compatible* method types
- When type checking a collection of class definitions, ensure that there are no cycles in the class hierarchy!

# How Do Abstract Classes Affect Our Type Checking Rules?

- If a method is called on a receiver whose static type is an abstract class, extract an arrow type from the declaration (just as with a definition in a concrete class)

`expr.area()`  $\mapsto$

`Shape.area()`  $\mapsto$

`()`  $\mapsto$  `Double`

# Type Checking Arguments to a Method Call

- The static types of an argument might no longer be an exact match:

```
abstract class Shape {  
  def area(): Double  
  
  def makeLikeMe(that: Shape): Shape  
}
```

(Let us set aside the concrete definitions of `makeLikeMe` for awhile)

# Now Consider a Call to Matcher With Concrete Types

`Circle(1).makeLikeMe(Circle(2)) ⇒`

`Circle.makeLikeMe(Circle) ⇒`

`(Shape → Shape)(Circle)`

And now we are stuck...

# Recall The Substitution Model of Type Checking

- To type check the application of a function to arguments:
  - Reduce the function to an arrow type
  - Reduce the arguments, left to right, to static types
  - If the argument types **match** the corresponding parameter types, reduce the application to the return type

# Subtyping

- We need to widen our definition of matching a type to include subtyping:
- A class is a subtype of the class it extends
- Subtyping is Reflexive:

$$A <: A$$

- Subtyping is Transitive:

$$\text{If } A <: B \text{ and } B <: C \text{ then } A <: C$$

# Subtyping

- All types are a subtype of type **Any**
- Type **Nothing** is a subtype of all types
  - There is no value with value type **Nothing**

# Recall The Substitution Model of Type Checking

- To type check the application of a function to arguments:
  - Reduce the function to an arrow type
  - Reduce the arguments, left to right, to static types
  - If the argument types **are subtypes of** the corresponding parameter types, reduce the application to the return type

# Applying a Class Method Revisited

- To reduce the application of a method:

$C(v_1, \dots, v_k).m(\text{arg1}, \dots, \text{argN})$

- Reduce the receiver and arguments, left to right
- **Reduce the body of  $m$** , replacing constructor parameters with constructor arguments and method parameters with method arguments

# Applying a Class Method Revisited

- To reduce the application of a method:

$C(v_1, \dots, v_k).m(\text{arg}_1, \dots, \text{arg}_N)$

- Reduce the receiver and arguments, left to right
- **Find the body of m in C and reduce to that,** replacing constructor parameters with constructor arguments and method parameters with method arguments

# The Body of $m$

- To find the body of method  $m$  in type  $C$ :
  - Find the definition of  $m$  in the body of  $C$ , if it exists
  - Otherwise, find the body of  $m$  in the immediate superclass of  $C$