Scheme Primitives and Function Definitions

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Course Overview

- Functional program design in Scheme
 - Data-directed (functional) program design 2-10
 - Algorithm design 11-14
 - Applied functional programming 15-17
- Object-oriented (OO) program design in Java 18-45

Today's Goals

- Common basic types
- Common primitive operations
- Rules for reducing programs
- Simple programs =

Variable definitions (Constants)

+ Function definitions

- The design recipe
- Errors
- Data definitions

Basic (primitive) types of data

numbers:

- naturals: 0, 1, 2, ...
 // number theory
- integers: ..., -1, 0, 1, ...
 - , ... // include negatives
- rational numbers: 3/4, 0, -1/3, ... // include fractions
- inexact numbers: #i0.123, #i0, ... // floating point numbers
 Operations: +, -, *, /, expt, remainder

Scheme computes exact answers on exact inputs if possible booleans: false, true

Operations: not, and, or, ...

Symbols: 'A, 'a, 'Aa, 'Corky, ...

Operations: ... // none important for now

Other basic types: strings, lists , ... // none important for now

Mixed-type Operations and Primitive Computation

- Basic relational operators
 - equal? // all data values
 - =, <, >, <=, >= // only on numbers
- Primitive computation \equiv application of a basic operation to constants
 - Basic operation \equiv basic function
 - Soon, we will see how to define our own (non-primitive) functions
- Function application in Scheme: parenthesized prefix notation
 - Scheme uses parenthesized prefix notation uniformly for everything
 - (+ 2 2), (sqrt 25), (remainder 7 3)
 - Bigger example: (* (+ 1 2) (+ 3 4))
 - How does this compare to writing 1+2*3+4 ?
- Scheme syntax is simple, uniform, and avoids possible ambiguity

Computation is repeated reduction

- Every Scheme program execution is the evaluation of a given expression constructed from primitive or defined functions and variables (names for constants).
- Evaluation proceeds by repeatedly performing the leftmost possible reduction (simplification) until the resulting expression is a value.
- A **value** is any constant. We will identify all of the expressions that are values as we explicate the language. Numbers, booleans, symbols are all values.

Reduction for primitive functions

- A reduction is an atomic computational step that replaces some expression by a simpler expression as specified by a Scheme evaluation rule (law). Every application of a basic operation to values yields a value (where run-time error is a special kind of value).
- Example

```
(* (+ 1 2) (+ 3 4))
```

```
=> (reduces to) (* 3 (+ 3 4))
```

```
=> (* 3 7) => 21
```

- Always perform leftmost reduction
- The following is **not** an atomic step, and so **not** a reduction
 (- (+ 1 3) (+ 1 3)) => 0

Programs = Variable Definitions + Function Definitions

Variables are simply names for values

• pi, my-SSN, album-name, tax-rate, x

Variable definitions

- (define freezing 32)
- (define boiling 212)
- Function definitions
 - (define (area-of-box x) (* x x))
 - (define (half x) (/ x 2))

Function applications (just as we saw before)

- (area-of-box 2)
- (half (area-of-box 3))
- Almost **any** function f used in a program can be written in the form

(define (f v1 ... vn) <expression>)

where <*expression*> is constructed from constants, variables, function applications, and a few other constructs TBN.

Reductions for defined functions

- Assume we have declared the two functions
 - (define (area-of-box x) (* x x))
 - (define (half x) (/ x 2))
- Then Scheme can perform these reductions

(half (area-of-box 3))

=> (half (* 3 3))

=> (half 9)

=> 4.5

Reduction stops when we get to a value or an error

Example: Solve quadratic equation

	Contract solve-quadratic: number number number \rightarrow number Purpose : (solve-quadratic a b c) finds the larger root of a*x*x + b*x + c = 0 given it has real roots and a != 0	Step 2		
;; ;; ;; ;;	<pre>Examples: (solve-quadratic 1 0 -25) = 5 (solve-quadratic 5 0 -20) = 2 (solve-quadratic 1 -10 25) = -4 and other examples</pre>	Step 3		
	;; Template instantiation: (degenerate) ;; (define (solve-quadratic a b c))			
	;; Code ;; (define (solve-quadratic a b c) ;; (/ (+ (- b) (sqrt (- (* b b) (* 4 a c)))) (* 2 a)))			
	Tests for solve-quadratic (check-expect exp ans) reports error if exp != ans (check-expect (solve-quadratic 1 0 -25) 5) (check-expect (solve-quadratic 5 0 -20) 2) (check-expect (solve-quadratic 1 -10 25) 5)	Step 6		

Syntax Errors

- A syntactically correct expression can be
 - An atomic expression, like
 - a number 17, 4.5, #i0.34
 - a variable radius
 - · A compound **expression**,
 - starting with (
 - \cdot followed by basic or program-defined operation such as + or ${\rm f}$
 - one or more **expression**s separated by spaces
 - ending with)
- Syntax errors:

Runtime Errors

- Happen when basic operations are applied to illegal arguments
- Consider the following examples:
 - (sqrt 1 2 3 4) => error: sqrt applied to more than one argument
 - (18 17) => error: 18 applied as function ;;
 - · (/ 1 0) => error: division by zero
 - · (+ 1 'a) => error: second argument in application of + is not a number
- If a reduction produces an error, the computation is aborted and the error is returned as the result.
- Try things like that in DrScheme, and make a mental note of the error messages you get back.

Conditional Expressions

- An expression that distinguishes different forms of data
- Form:

```
(cond [question-1 result-1]
  [question-2 result-2]
  ...
  [question-n result-n]
  [else default-result])
```

- Square brackets are used above for clarity. In Scheme, they are synonymous with parentheses, but balancing brackets must match.
- else is optional. If omitted and none of the questions are true, the result is a run-time error (like division by zero).

Reduction of Conditional Expressions

	(cond	[true	result-1]	
	_	[])	
=> result-1				
	(cond	[false	result-1]	
		[question-2	result-2]	
		• • •		
		[else	default-result])	
=>	(cond	[question-2	result-2]	
		• • •		
		[else	default-result])	
		[foloo		
	(cond	[false	result-1]	
		[else	default-result])	
=> default-result				

Conditional Expression Examples

```
(cond [(> 12 0) 5] [else -5])
  \Rightarrow (cond [true 5] [else -5])
  => (cond [true 5] [else -5])
Given
 (define (abs x)
     (cond [(>= x 0) x])
            [else (- x)]))
      (abs -10)
  => (cond [(>= -10 0) -10] [else (- -10)])
  => (cond [false -10] [else (- -10)])
  \Rightarrow (cond [else (- -10)]) \Rightarrow (- -10) \Rightarrow 10
```

The Design Recipe

How should I go about writing programs?

- 1. Analyze problem and define any requisite data types.
- 2. State contract (type) and purpose for *function* that solves the problem.
- 3. Give examples of function use and result.
- 4. Select and instantiate a template for the function body.
- 5. Write the function itself.
- 6. Test it, and confirm that tests succeeded.

The order of the steps of the recipe is important

The Design Recipe (Big Picture)

- Encourages systematic problem solving
- Works best if keep our functions small
- We will learn how to repeatedly decompose problems into simpler problems until we reach problems that can be solved by simple expressions as in solve-quadratic
- Decomposition driven by structure of data being processed: *data-directed* design

Reminders

- New homework (HW1) is posted online
 - Due next Friday, so you will get to check it over in lab; don't wait until your lab to get started.
 - Sign up for mailing list to get any updates, discussions
 - Make absolutely sure you follow the recipe in writing Scheme programs.
 - Partners: Talk to people after class, at lab, etc.
 - For Scheme programs, follow format of the sample solution in the Scheme HW Guide.
 - For hand evaluations, follow the format of the hand evaluation problems posted in the Scheme HW Guide.
 - Submit your assignment using Owlspace.

Epilog

- Reminder: continue digesting chs. 1-10 in HTDP Section 8.3 is particularly important and it is not wordy.
- Next class
 - Inductive Data definitions
 - Amplified design recipe
- Challenge problem: What happens if we use rightmost reduction instead of leftmost? Can you devise a program using the Scheme subset given in this lecture such that some invocation of that program (expression composed from constants and and basic and program-defined operations defined in the program) behaves differently (either in terms the result produced by the computation or lack thereof) under rightmost evaluation than leftmost evaluation. Hint: focus on pathological behavior and note that two different errors are not equivalent.