COMP 322: Fundamentals of Parallel Programming

Lecture 13: Forall Statements & Barriers (contd)

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https://wiki.rice.edu/confluence/display/PARPROG/COMP322



HJ's forall statement = finish + forasync + barriers (next)

```
AtomicInteger rank = new AtomicInteger();
forall (point[i] : [0:m-1]) {
  int r = rank.getAndIncrement();
  System.out.println("Hello from task ranked " + r);
  rext; // Acts as barrier between phases 0 and 1
  System.out.println("Goodbye from task ranked " + r)  Phase 1
}
```

- next → each forall iteration suspends at next until all iterations arrive (complete previous phase), after which the phase can be advanced
 - If a forall iteration terminates before executing "next", then the other iterations do not wait for it
 - Scope of synchronization is the closest enclosing for all statement
 - Special case of "phaser" construct (will be covered in following lectures)



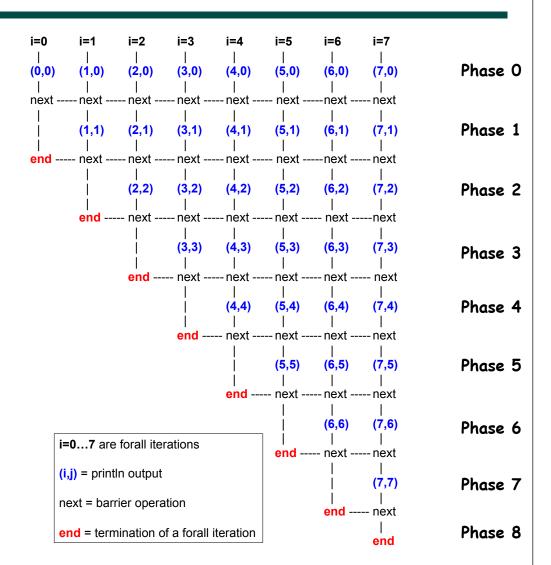
Recap of Observation 2 (Lecture 12): If a forall iteration terminates before "next", then other iterations do not wait for it

- Outer forall-i loop has m iterations, 0...m-1
- Inner sequential j loop has i+1 iterations, 0...i
- Line 4 prints (task, phase) = (i, j) before performing a next operation.
- Iteration i = 0 of the forall-i loop prints (0, 0), performs a next, and then terminates. Iteration i = 1 of the forall-i loop prints (1,0), performs a next, prints (1,1), performs a next, and then terminates. And so on.



Illustration of Observation 2

- Iteration i=0 of the forall-i loop prints (0, 0) in Phase 0, performs a next, and then ends Phase 1 by terminating.
- Iteration i=1 of the forall-i loop prints (1,0) in Phase 0, performs a next, prints (1,1) in Phase 1, performs a next, and then ends Phase 2 by terminating.
- And so on until iteration i=8 ends an empty Phase 8 by terminating





Recap of Observation 3 (Lecture 12): Different forall iterations may perform "next" at different program points

```
forall (point[i] : [0:m-1]) {
2.
     if (i % 2 == 1) { // i is odd
3.
       oddPhaseO(i);
4.
      next;
5.
  oddPhase1(i);
6.
  } else { // i is even
  evenPhase0(i);
7.
8. next;
9.
  evenPhase1(i);
10. } // if-else
11. } // forall
```

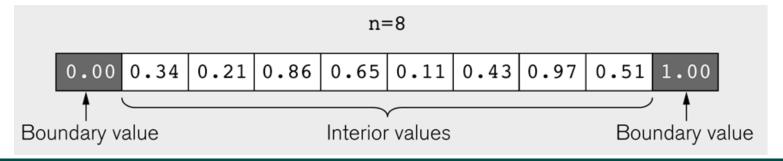
- Barrier operation synchronizes odd-numbered iterations at line 4 with even-numbered iterations in line 8
- next statement may even be in a method such as oddPhase1()



One-Dimensional Iterative Averaging Example (Lecture 10)

- Initialize a one-dimensional array of (n+2) double's with boundary conditions, myVal[0] = 0 and myVal[n+1] = 1.
- In each iteration, each interior element myVal[i] in 1..n is replaced by the average of its left and right neighbors.
 - Two separate arrays are used in each iteration, one for old values and the other for the new values
- After a sufficient number of iterations, we expect each element of the array to converge to myVal[i] = i/(n+1)
 - In this case, myVal[i] = (myVal[i-1]+myVal[i+1])/2, for all i in 1..n

Illustration of an intermediate step for n = 8 (source: Figure 6.19 in Lin-Snyder book)





HJ code for One-Dimensional Iterative Averaging with nested for-forall structure

```
1. double[] myVal=new double[n+2]; double[] myNew=new double[n+2];
2. myVal[n+1] = 1; // Boundary condition
3. for (point [iter] : [0:numIters-1]) {
    // Compute MyNew as function of input array MyVal
4.
    forall (point [j] : [1:n]) { // Create n tasks
5.
       myNew[j] = (myVal[j-1] + myVal[j+1])/2.0;
6.
7.
   } // forall
8.
   // Swap myVal and myNew
    double[] temp=myVal; myVal=myNew; myNew=temp;
9.
10. // myNew becomes input array for next iteration
11.} // for
```

- Replace "finish async" from Lecture 10 by "forall"
- Overhead issue --- this version creates (numIters * n) async tasks



HJ code for One-Dimensional Iterative Averaging with chunked for-forall-for structure

```
double[] myVal=new double[n+2]; double[] myNew=new double[n+2];
2. myVal[n+1] = 1; // Boundary condition
3. // Set desired number of chunks for j loop to total number of workers
  int Cj = Runtime.getNumOfWorkers();
  for (point [iter] : [0:numIters-1]) {
     // Compute MyNew as function of input array MyVal
6.
     forall (point [jj]:[0:Cj-1]) // Iterate over chunks
7.
       for (point [j]:getChunk([1:n],[Cj],[jj])) // Iterate within chunk
8.
9.
          myNew[j] = (myVal[j-1] + myVal[j+1])/2.0;
     // Swap myVal and myNew
10.
11.
     double[] temp=myVal; myVal=myNew; myNew=temp;
     // myNew becomes input array for next iteration
12.
13.} // for iter
```

- Chunked forall version creates numIters*Cj async tasks
- Can we do better with barriers?



One-Dimensional Iterative Averaging with Barrier Synchronization

```
double[] qVal=new double[n+2]; double[] qNew=new double[n+2]; qVal[n+1] = 1;
   int Cj = Runtime.getNumOfWorkers();
   forall (point [jj]:[0:Cj-1]) { // Chunked forall is now the outermost loop
     double[] myVal = gVal; double[] myNew = gNew; // Local copy of myVal/myNew pointers
4.
5.
     for (point [iter] : [0:numIters-1]) {
       // Compute MyNew as function of input array MyVal
6.
       for (point [j]:getChunk([1:n],[Cj],[jj])) // Iterate within chunk
7.
8.
          myNew[j] = (myVal[j-1] + myVal[j+1])/2.0;
       next; // Barrier before executing next iteration of iter loop
9.
       // Swap myVal and myNew (each forall iterations swaps its pointers in local vars)
10.
        double[] temp=myVal; myVal=myNew; myNew=temp;
11.
       // myNew becomes input array for next iter
12.
                                                                                     Thread 3
                                                                     Thread 2
                                                      Thread 1
13.
     } // for
                                         i iteration
14. } // forall
                                                           (b) Barrier synchronization
```

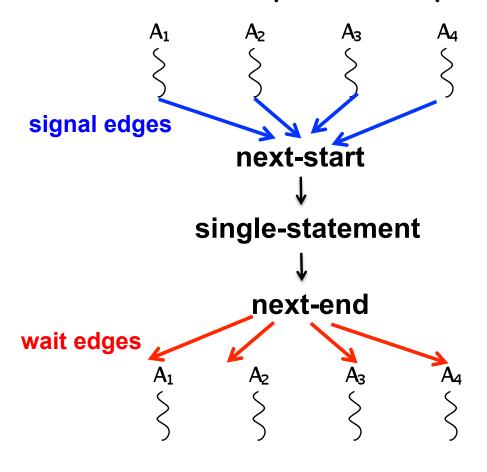
- · Use of barrier reduces number of async tasks created to just Cj
- · However, these Cj tasks perform Cj*numIters barrier operations
 - Good trade-off since, barrier operations have lower overhead than task creation if number of chunks <= number of workers</p>



Next-with-Single Statement

next <single-stmt > is a barrier in which single-stmt is performed exactly once after all tasks have completed the previous phase and before any task begins its next phase.

Modeling next-with-single in the Computation Graph





Use of next-with-single to print a log message between Hello and Goodbye phases

```
1. AtomicInteger rank = new AtomicInteger();
2. forall (point[i] : [0:m-1]) {
3. // Start of Hello phase
4. int r = rank.getAndIncrement();
5.
    System.out.println("Hello from task ranked " + r);
6.
   next single {
    System.out.println("LOG: Between Hello & Goodbye Phases");
7.
8.
    // Start of Goodbye phase
9.
10.
    System.out.println("Goodbye from task ranked " + r);
11.} // forall
```

