

Comp 311

Functional Programming

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The Substitution Rule of Computation

- To reduce an application of a function to a set of arguments:
 - Reduce the arguments, left to right
 - Reduce the body of the function, with each parameter replaced by the corresponding argument

Using the Substitution Rule

$$f(x, y) = x^2 + y^2$$

$$f(4 - 5, 3 + 1) \mapsto$$

$$f(-1, 3 + 1) \mapsto$$

$$f(-1, 4) \mapsto$$

$$-1^2 + 4^2 \mapsto$$

$$1 + 16 \mapsto$$

$$17$$

What About Types?

- Eventually, we learn that our functions need to include rules indicating the required types of their arguments, and the types of applications
- You might have seen notation like this in a math class:

$$f: \mathbf{Z} \rightarrow \mathbf{Z}$$

Typing Rules for Functions

$$f: \mathbf{Z} \rightarrow \mathbf{Z}$$

What does this rule mean?

Typing Rules for Functions

$$f: \mathbf{Z} \rightarrow \mathbf{Z}$$

- We can interpret the arrow as denoting data flow:

The function f consumes arguments with value type \mathbf{Z} and produces values with value type \mathbf{Z}

(or one of a well-defined set of exceptional events occurs).

Typing Rules for Functions

$$f: \mathbf{Z} \rightarrow \mathbf{Z}$$

- We can also interpret the arrow as logical implication:

If f is applied to an argument expression with static type \mathbf{Z} then the application expression has static type \mathbf{Z} .

What are The Exceptional Events in Algebra?

- A “division by zero” error
- We run out of some finite resource
- What else?

The Substitution Rule Allows for Computations that Never Finish

$$f: \mathbf{Z} \times \mathbf{Z} \rightarrow \mathbf{Z}$$

$$f(x, y) = f(x, y)$$

$$f(4 - 5, 3 + 1) \mapsto$$

$$f(-1, 3 + 1) \mapsto$$

$$f(-1, 4) \mapsto$$

$$f(-1, 4) \mapsto$$

...

The Substitution Rule Allows for Computations that Keep Getting Larger

$$f: \mathbf{Z} \times \mathbf{Z} \rightarrow \mathbf{Z}$$

$$f(x, y) = f(f(x, y), f(x, y))$$

$$f(4 - 5, 3 + 1) \mapsto$$

$$f(-1, 3 + 1) \mapsto$$

$$f(-1, 4) \mapsto$$

$$f(f(-1, 4), f(-1, 4)) \mapsto$$

$$f(f(f(-1, 4), f(-1, 4)), f(f(-1, 4), f(-1, 4))) \mapsto$$

...

But We Need at Least Limited Recursion to
Define Common Algebraic Constructs

$$\begin{array}{c} !: \\ \mathbf{N} \rightarrow \mathbf{N} \end{array}$$

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ n(n-1)! & \text{if } n > 0 \end{cases}$$

What are The Exceptional Events in Algebra?

- A “division by zero” error
- We run out of some finite resource
- The computation never stops
(unbounded time)
- The computation keeps getting larger
(unbounded space)

*Our third exposure
to computation:
Core Scala*

Core Scala

- We will continue to use algebra as our model of computation
- We will switch to Scala syntax
- We will introduce new value types

Value Types in Core Scala

Int: -3, -2, -1, 0, 1, 2, 3

Double: 1.414, 2.718, 3.14, ∞

Boolean: false, true

String: "Hello, world!"

Primitive Operators on Ints and Doubles in Core Scala

Algebraic operators:

$e + e'$ $e - e'$ $e * e'$ e / e'

- For each operator:
 - If both arguments to an application of an operator are of type `Int` then the application is of type `Int`
 - If both arguments to an application of an operator are of type `Double` then the application is of type `Double`

Primitive Operators on Ints and Doubles in Core Scala

Comparison operators:

$$e == e' \quad e \leq e' \quad e \geq e' \quad e \neq e'$$
$$e > e' \quad e < e'$$

- For each operator:
 - If both arguments to an application of an operator are of type Int then the application is of type Boolean
 - If both arguments to an application of an operator are of type Double then the application is of type Boolean

Some Primitive Operators on Booleans in Core Scala

Conjunction, Disjunction:

$$e \ \& \ e' \quad e \ | \ e'$$

- In both cases:
 - If both arguments to an application are of type Boolean then the application is of type Boolean

More Primitive Operators on Booleans in Core Scala

Negation:

`!e`

- If the argument to an application is of type `Boolean` then the application is of type `Boolean`

Yet More Primitive Operators on Booleans in Core Scala

Conditional Expressions:

$\text{if } (e) \ e' \ \text{else } e''$

- If the first argument is of type `Boolean` and the second and third argument are of the same type T then the application is of type T

Primitive Operators on Strings in Core Scala

String Concatenation:

$$e + e'$$

- If both arguments are of type `String` then the application is of type `String`

An Example Function Definition in Core Scala

```
def square(x: Double) = x * x
```

Syntax for Defining Functions

```
def fnName(arg0: type0, ..., argk: typek): returnType =  
    expr
```

- If there is no recursion, we may elide the return type:

```
def fnName(arg0: type0, ..., argk: typek) =  
    expr
```

The Substitution Rule Works as Before

```
def square(x: Double) = x * x
```

```
square(2.0 * 3.0) ↪  
  square(6.0) ↪  
    6.0 * 6.0 ↪  
      36.0
```


The Nature of Ints

Fixed Size Ints

- Unlike the integers we might write on a sheet of paper, the values of type `Int` are of a fixed size.
- For every `n: Int`,

$$-2^{31} \leq n \leq 2^{31} - 1$$

Fixing the Size of Numbers Has Many Benefits

- The time needed to compute the application of an operation on two numbers is bounded.
- The space needed to store a number is bounded.
- We can easily reuse the space used for one number to store another.

But We Need to Concern Ourselves with Overflow

- If we compute a value larger than $2^{31} - 1$, our representation will “wrap around” (i.e., overflow):

$$2147483647 + 1 \mapsto -2147483648$$

The Moral of Computing with Ints

- If possible, determine the range of potential results of a computation
 - Ensure that this range is no larger than the range of representable values of type Int
- Otherwise, include in your computation a check for overflow

The Nature of Doubles

Scientific Notation

- Numeric values in scientific computations can span enormous ranges, from the very large to the very small
- At the same time, scientific measurements are of limited precision
- “Scientific notation” was devised in order to efficiently represent approximate values that span a large range

Scientific Notation

$$6.022 \times 10^{23}$$

mantissa

exponent

Scientific Notation and Efficient Computation

- We normalize the mantissa so that its value is at least 1 but less than 10
- If we
 - Set the number of digits in the mantissa to a fixed precision, and
 - Set the number of digits in the exponent to a fixed precision
- Then all numbers in our notation are of a fixed size

Doubles

- Values of type Double are stored as with fixed sized numbers in scientific notation, but with a few differences:
- Finite, nonzero numeric values can be expressed in the form:

$$\pm m 2^e$$

Doubles

$$\pm m 2^e$$

- $1 \leq m \leq 2^{53}-1$
- $-2^{10}-53+3 \leq e \leq 2^{10}-53$

Doubles

$$\pm m 2^e$$

- $1 \leq m \leq 2^{53}-1$
- $-2^{10}-53+3 \leq e \leq 2^{10}-53$
- $-1074 \leq e \leq 971$

The Nature of Doubles

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Scientific Notation

$$\underbrace{6.022}_{\text{mantissa}} \times \underbrace{10}_{\text{base}} \underbrace{23}_{\text{exponent}}$$

Scientific Notation and Efficient Computation

- We normalize the mantissa so that its value is at least 1 but less than 10
- If we
 - Set the number of digits in the mantissa to a fixed precision, and
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- Values of type Double are stored as with fixed sized numbers in scientific notation, but with a few differences:
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Doubles

$$\pm m \times 2^e$$

- $1 \leq m \leq 2^{53} - 1$
- $-1022 \leq e \leq 971$

For more details, you can read about double-precision binary representation:
https://en.wikipedia.org/wiki/Double-precision_floating-point_format

Representations of Doubles

- Many quantities have more than one representation in this format:

$$1024 \times 2^{500}$$

$$512 \times 2^{501}$$

Distances Between Doubles

- The distance between adjacent values of type Double is not constant
 - The values are most dense near zero
 - They grow sparser exponentially as one moves away from zero

Operations and Rounding

- Arithmetic operations round to the closest representable value
- Ties are broken by choosing the value with the smaller absolute value

Overflow with Doubles

- Computations on Doubles that result in values larger than the largest finite Double are represented with special values:

Double.PositiveInfinity

Double.NegativeInfinity

Underflow with Doubles

- Computations on Doubles that result in values with magnitudes smaller than the smallest non-zero Double are represented with special values:

0.0 -0.0

Division By Zero

- Division of a non-zero finite value by a zero value results in an infinite value:

$1.0 / 0.0 \mapsto \text{Double.PositiveInfinity}$

$1.0 / -0.0 \mapsto \text{Double.NegativeInfinity}$

Division By Zero

- As does division of an infinite value by a zero value:

`Double.PositiveInfinity / 0.0` \mapsto `Double.PositiveInfinity`

Division By Zero

- Division of a zero value by a zero value results in another special value NaN (for “Not a Number”):

$0.0 / 0.0 \mapsto \text{Double.NaN}$

$-0.0 / 0.0 \mapsto \text{Double.NaN}$

Doubles Break Common Algebraic Properties

- Equality is not reflexive:

`Double.NaN != Double.NaN`

- Multiplication does not distribute over addition:

`100.0 * (0.1 + 0.2) ↪
30.000000000000000000000004`

`100.0 * 0.1 + 100.0 * 0.2 ↪
30.0`

Morals of Floating Point Computation

- Avoid floating point computation whenever you need to compute precise numeric values (such as monetary values)
- Use floating point values only when calculating with inexact measurements over a range larger than can be represented with precise arithmetic

Morals of Floating Point Computation

- Try to bound the margin of error in your calculation
- Don't test for equality directly
- Instead of writing:

`x == y`

- Write:

`abs(x - y) <= tolerance`

Defining Absolute Value

```
def abs(x: Double) = if (x >= 0) x else -x
```

What's wrong here?

```
abs(-0.0) ↦
```

```
if (-0.0 >= 0) -0.0 else -(-0.0) ↦
```

```
if (true) -0.0 else -(-0.0) ↦
```

```
-0.0
```

Defining Absolute Value

```
def abs(x: Double) = if (x > 0) x else 0.0 - x
```

Does it work now?

```
abs(-0.0) ↦
```

```
if (-0.0 > 0) -0.0 else 0.0 - -0.0 ↦
```

```
if (false) -0.0 else 0.0 - -0.0 ↦
```

```
0.0 - -0.0
```

```
0.0
```


What are The Exceptional Events in Core Scala?

- A “division by zero” error on Ints (but not Doubles)
- We run out of some finite resource
 - The computation never stops
 - The computation uses too much memory