

Comp 311

Functional Programming

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Announcements

- Homework 2 is due two weeks from today
- Assignment description PDF on Piazza
- No provided “skeleton” code
- Simple interface (compilation/linking) check provided

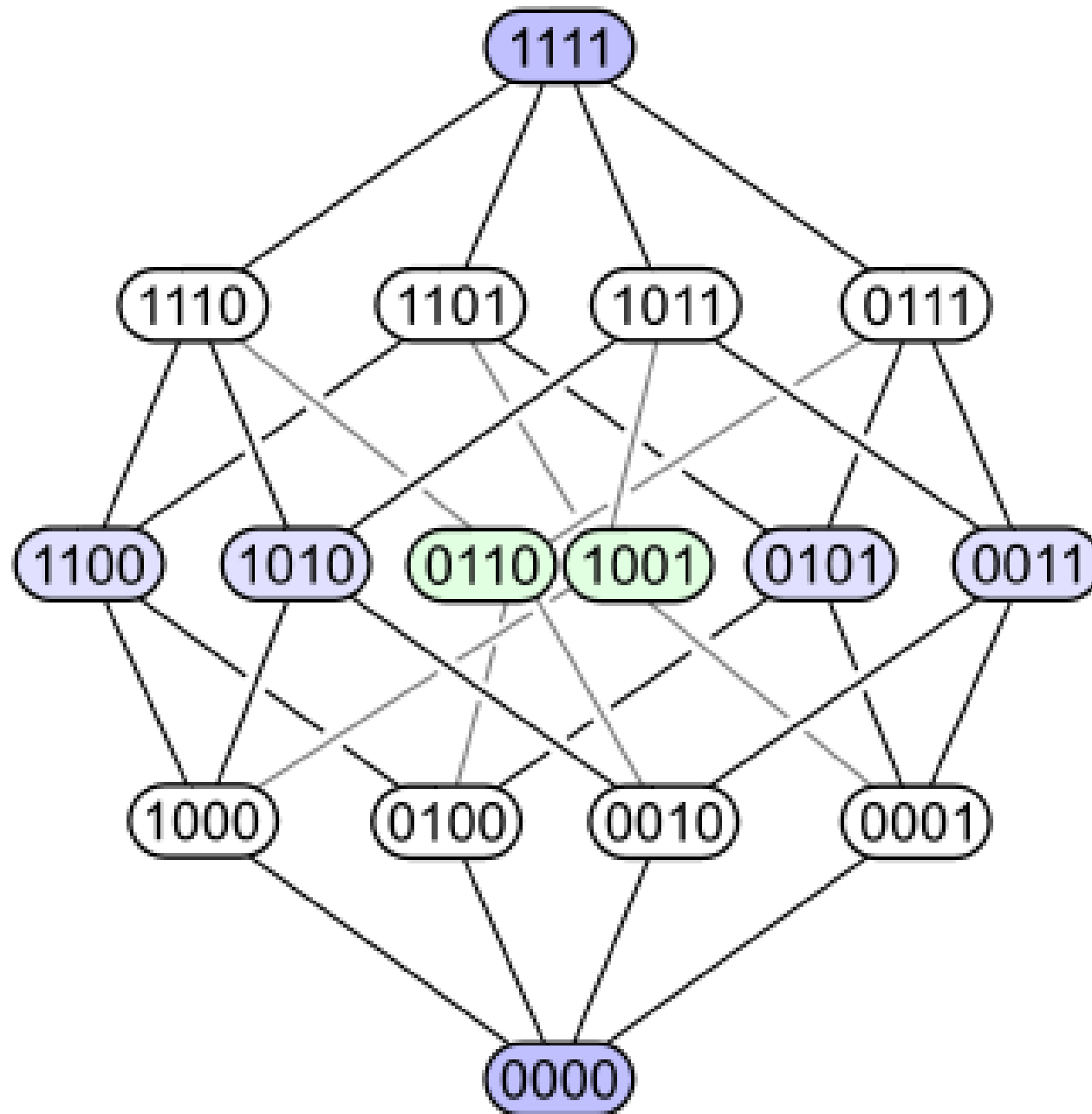
Scala Type Hierarchy

Type Hierarchies

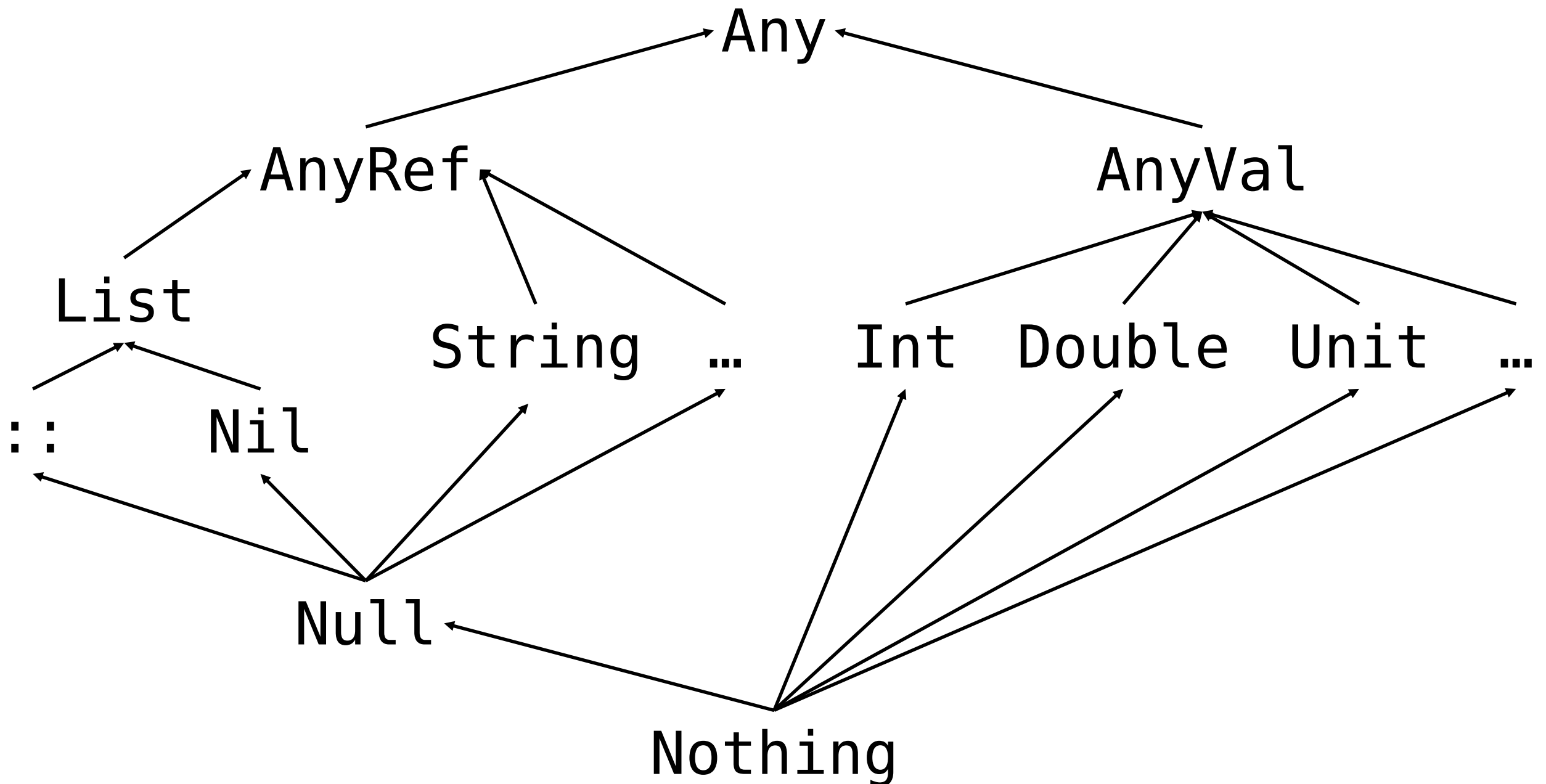
Inheritance (subclass / superclass relationships) form a *complete lattice* in the Scala type system:

- Each pair of classes has exactly one:
 - *Least upper-bound*
 - *Greatest lower-bound*
- The same applies to all value types

Hasse Diagrams



Scala Type Lattice



Parametric Polymorphism (Parametric/Generic Types)

Parametric Types

- We have defined two forms of lists: lists of ints and lists of shapes
- Many computations useful for one are useful for the other:
 - Map, reduce, filter, etc.
- It would be better to define lists and their operations once for all of these cases

Parametric Types

- Higher-order functions take functions as arguments and return functions as results
- Likewise, *parametric types*, a.k.a., a *generic types*, takes types as arguments and return types as results

Parametric Lists

- Every application of this parametric type to an argument yields a new type:

```
abstract class List[T] {  
    def ++(ys: List[T]): List[T]  
}
```

Parametric Lists

- Every application of this parametric type to an argument yields a new type:

```
abstract class List[T <: Any] {  
  def ++(ys: List[T]): List[T]  
}
```



- We augment the declarations of type parameters to permit an upper bound on all instantiations of a parameter
- By default, the bound is *Any*

Syntax of Parametric Class

Definitions

```
<modifiers> class C[T1 <: N, ..., TN <:
N] extends N {
    <ordinary class body>
}
```

- We denote type parameters as T1, T2, etc.
- We denote all other types with N, M, etc.

Syntax of Parametric Class

Definitions

```
<modifiers> class C[T1 <: N, ..., TN <:
N] extends N {
    <ordinary class body>
}
```

- Declared type parameters T1, ..., TN are in scope throughout the entire class definition, including:
 - The bounds of type parameters
 - The `extends` clause
- Object definitions must not be parametric

Parametric Lists

- Every application of this parametric type yields a new type:

```
List[Int]  
List[String]  
List[List[Double]]  
etc.
```

Parametric Lists

- Every application (a.k.a., *instantiation*) of this parametric type yields a new type:

```
abstract class List[T] {  
  def ++(ys: List[T]): List[T]  
}
```



Note that our parametric type can be instantiated with type parameters, including its own!

Parametric Lists

```
case class Empty[S]() extends List[S] {  
  def ++(ys: List[S]) = ys  
}
```

```
case class Cons[T](head: T, tail: List[T]) extends List[T] {  
  def ++(ys: List[T]) = Cons[T](head, tail ++ ys)  
}
```

Our definition requires a separate type `Empty[S]` for every instantiation of `S`. Thus we must define `Empty` as a class rather than an object.

Covariance

- Can one instantiation of a parametric type be a subtype of another?
- Currently our rules allow this only in the reflexive case:

`List[Int] <: List[Int] in E`

Covariance

- It would be useful to allow some instantiations to be subtypes of another
- For example, we would like it to be the case that:

```
List[Int] <: List[Any]
```

Covariance

- In general, we say that a parametric type C is covariant with respect to its type parameter S if:

$$S <: T \text{ in } E$$

implies

$$C[S] <: C[T] \text{ in } E$$

- We must be careful that such relationships do not break the soundness of our type system

Covariance

- For a parametric type such as:

```
abstract class List[T <: Any] {  
    def ++(ys: List[T]): List[T]  
}
```

- And types S and T , such that $S <: T$ in some environment E :
 - What must we check about the body of class `List` to allow for `List[S] <: List[T]` in E ?

Covariance

- Consider instantiations for types `String` and `Any`:

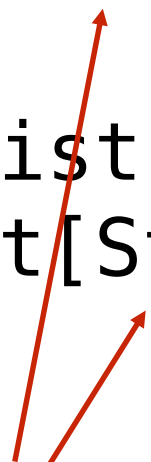
```
abstract class List[Any] {  
  def ++(ys: List[Any]): List[Any]  
}  
abstract class List[String] {  
  def ++(ys: List[String]): List[String]  
}
```

Covariance

- If these were ordinary classes connected by an `extends` class:
 - We would need to ensure that the overriding definition of `++` in class `List[String]` was compatible with the overridden definition in `List[Any]`

Covariance

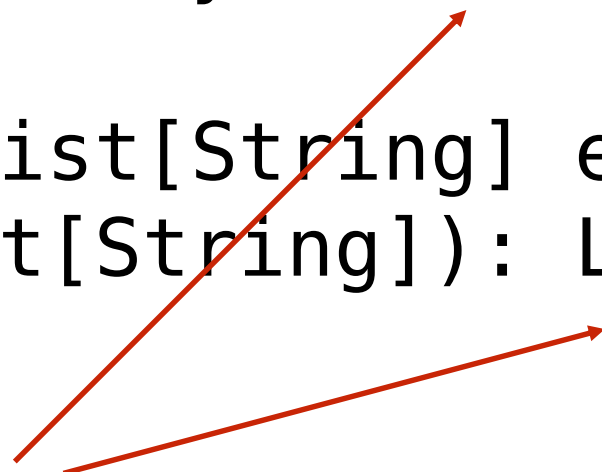
```
abstract class List[Any] {  
  def ++(ys: List[Any]): List[Any]  
}  
abstract class List[String] extends List[Any] {  
  def ++(ys: List[String]): List[String]  
}
```



But if `List[String] <: List[Any]` in E
then this is not a valid override

Covariance

```
abstract class List[Any] {  
  def ++(ys: List[Any]): List[Any]  
}  
abstract class List[String] extends List[Any] {  
  def ++(ys: List[String]): List[String]  
}
```



On the other hand, the return types
are not problematic

Covariance

- From our example, we can glean the following rule:
 - We allow a parametric class C to be covariant with respect to a type parameter T so long as T does not appear in the types of the method parameters of C

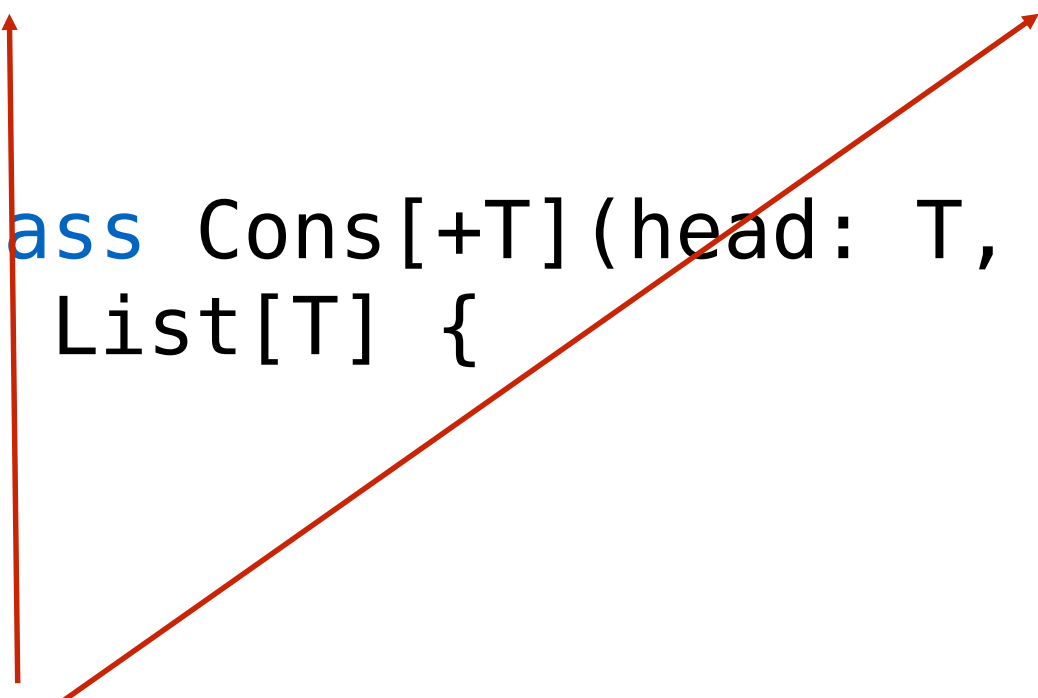
Covariance

```
abstract class List[+T] {}
```

- We stipulate that a parametric type is covariant in a parameter T by prefixing a $+$ at the definition of T
- (We will return to our definition of `append` later)

Covariance

```
case object Empty extends List[Nothing] {  
}  
  
case class Cons[+T](head: T, tail: List[T])  
extends List[T] {  
}
```



Now we can define Empty as an object that extends the bottom of the List types

Covariance and Append

- The problem with our original declaration of `append` was that it was not general enough:
 - There is no reason to require that we always append lists of identical type
 - Really, we can append a `List[S]` for any supertype of our `List[T]`
 - The result will be of type `List[S]`

Lower Bounds on Type Parameters

- Thus far, we have allowed type parameters to include upper bounds:

$$T <: S$$

- They can also include lower bounds:

$$T >: U$$

- Or they can include both:

$$T >: S <: U$$

Parametric Functions

- Just as we can add type parameters to a class definition, we can also add them to a function definition
- The type parameters are in scope in the header and body of the function

Covariance and Append

```
abstract class List[+T] {  
  def ++[S >: T](ys: List[S]): List[S]  
}  
  
case object Empty extends List[Nothing] {  
  def ++[S](ys: List[S]) = ys  
}  
  
case class Cons[+T](head: T, tail: List[T])  
extends List[T] {  
  def ++[S >: T](ys: List[S]) = Cons(head, tail ++ ys)  
}
```

Map Revisited

```
abstract class List[+T] {  
  ...  
  def map[U](f: T => U): List[U]  
}
```