

Comp 311

Functional Programming

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October 2, 2018

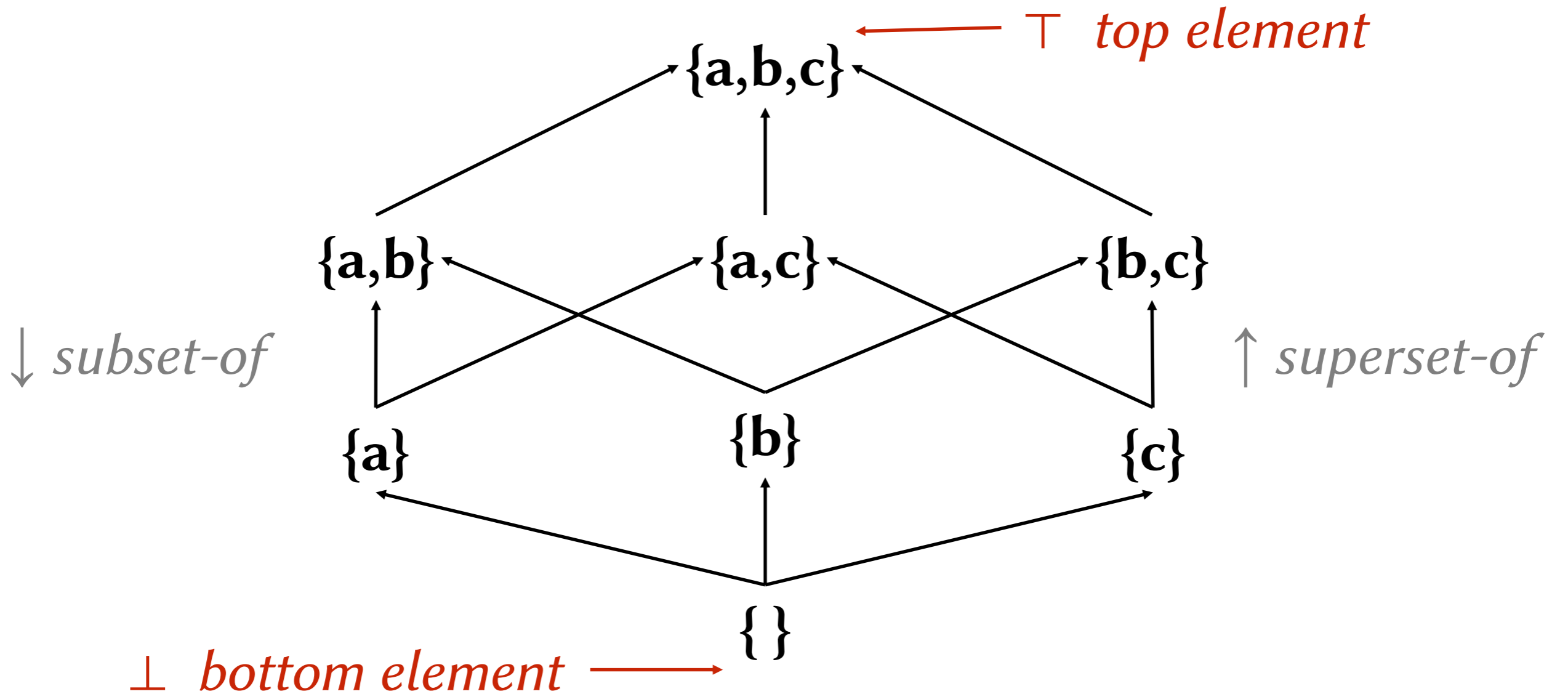
Scala Type Hierarchy

Type Hierarchies

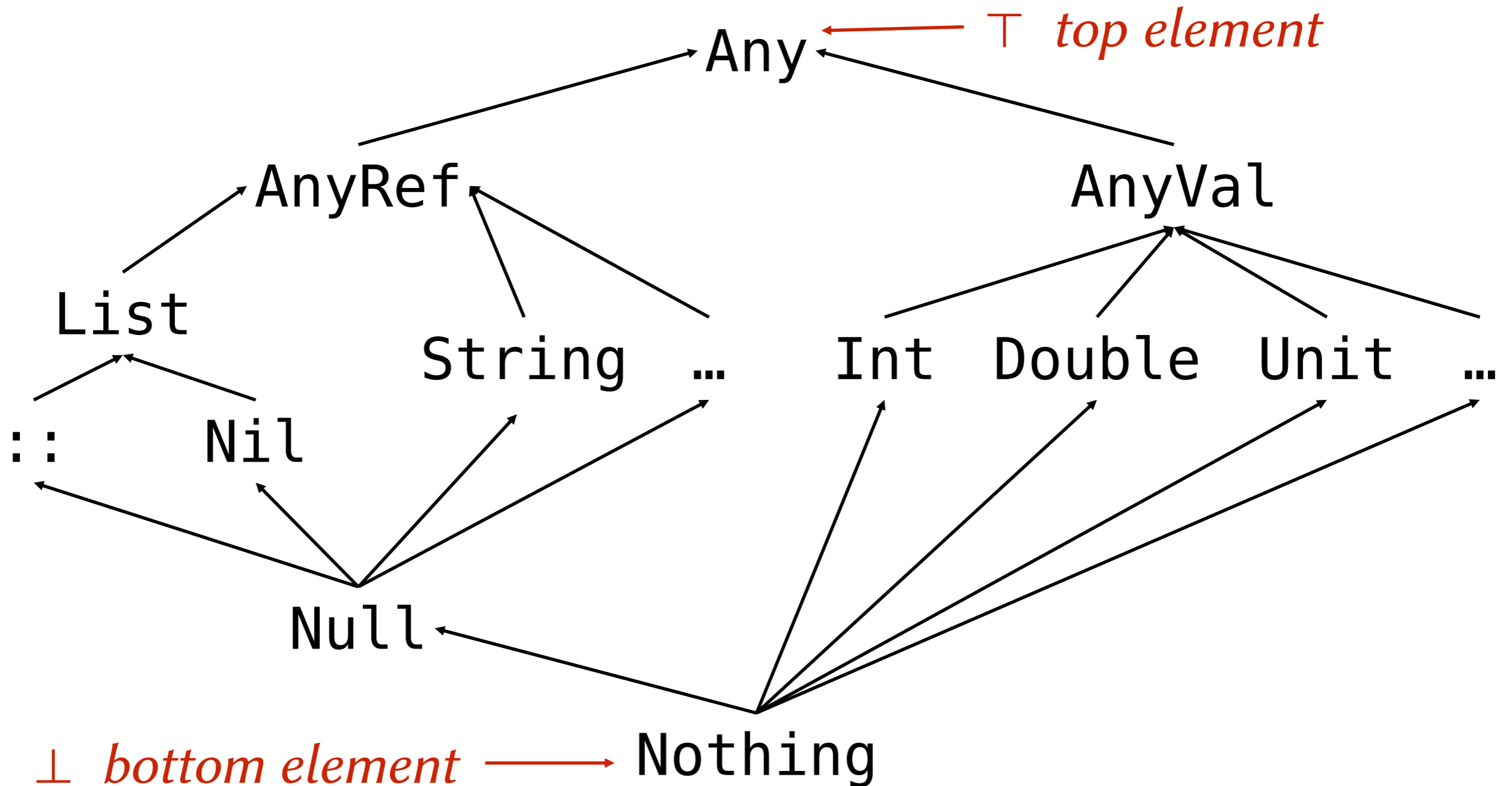
Inheritance (subclass / superclass relationships) form a *complete lattice* in the Scala type system:

- Each pair of classes has exactly one:
 - *Least upper-bound*
 - *Greatest lower-bound*
- The same applies to all value types

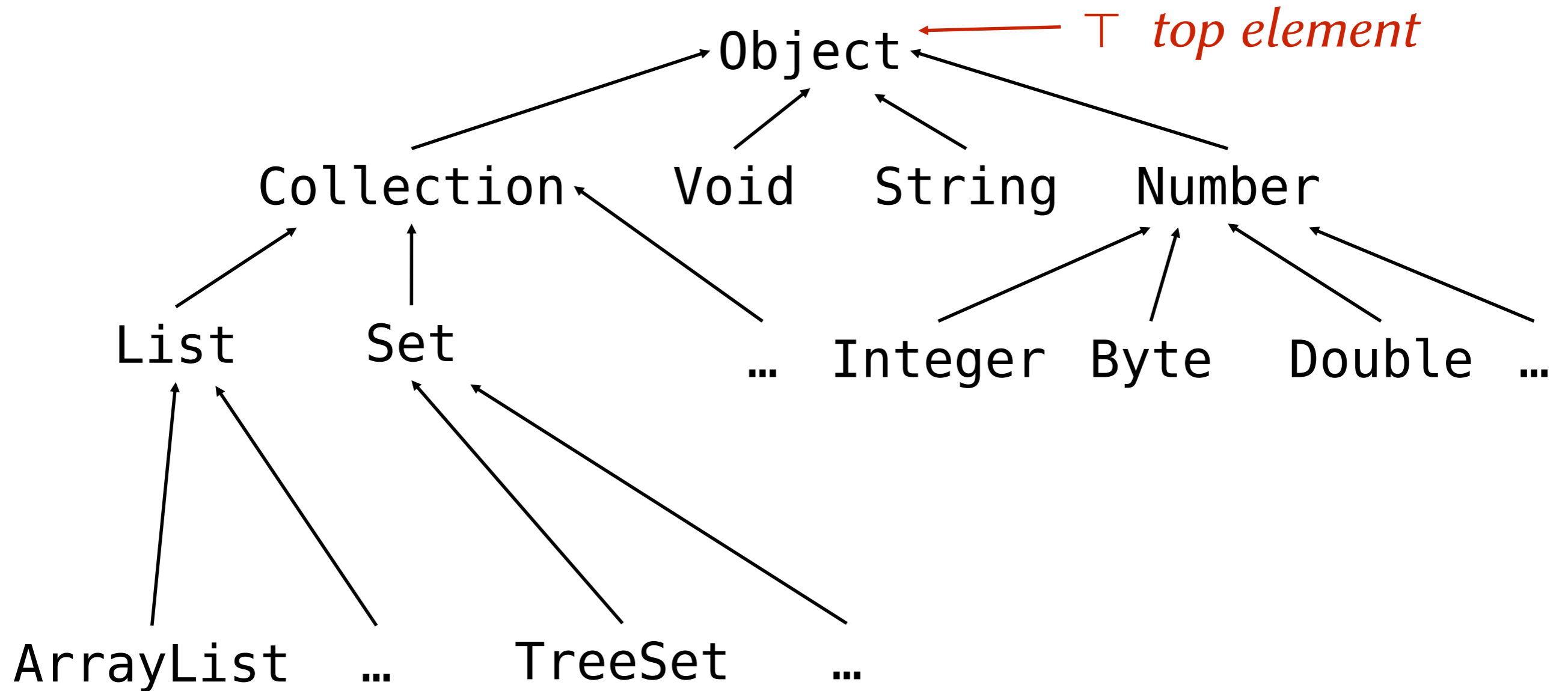
Hasse Diagrams



Scala Type Lattice



Java Type Semi-Lattice



no bottom element

Variance of Parametric Types

Covariance

In general, we say that a parametric type C is covariant with respect to its type parameter S if:

$$S <: T$$

implies

$$C[S] <: C[T]$$

Contravariance

In general, we say that a parametric type C is contravariant with respect to its type parameter S if:

$$S <: T$$

implies

$$C[T] <: C[S]$$

Invariance

In general, we say that a parametric type C is invariant with respect to its type parameter S if:

$$S <: T$$

implies neither

$$C[S] <: C[T]$$

nor

$$C[T] <: C[S]$$

Syntax for Variance

Syntactically:

- Covariant type parameter declarations are annotated with a *plus* sign.
- Contravariant type parameter declarations are annotated with a *minus* sign.
- Invariant type parameter declarations are not annotated with an extra symbol.

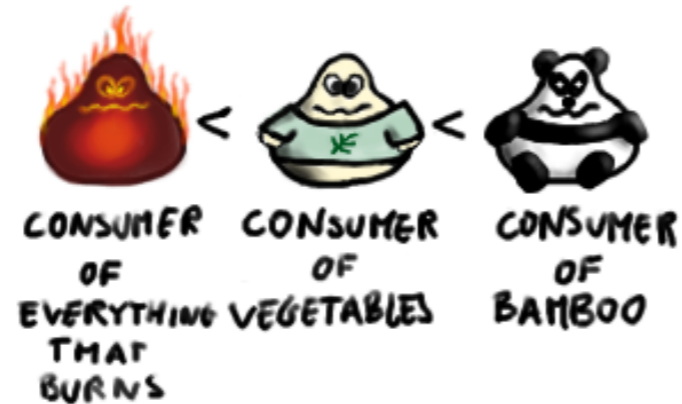
```
case class X[+A, -B, C]
```

CONTRAVARIANCE:

HIERARCHY OF X:



CONSUMERS [-X]:



... YOU CAN GIVE IT TO:

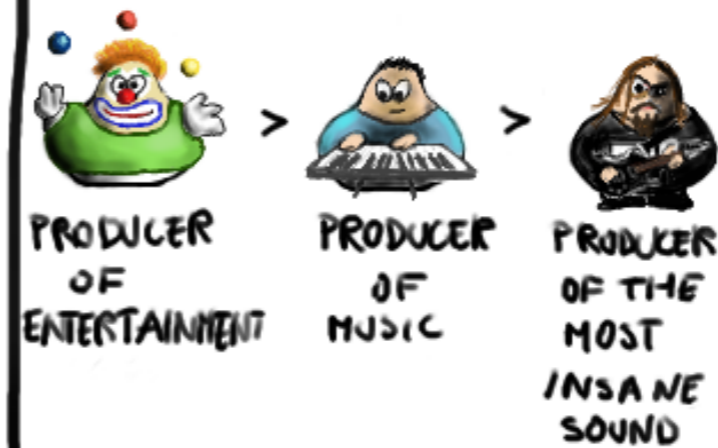


COVARIANCE:

HIERARCHY OF X:



PRODUCERS [+X]:



... YOU CAN GET IT FROM:



Image by Andrey Tyukin:

<https://stackoverflow.com/a/19739576/1427124>

See also: <https://www.cs.rice.edu/~javaplt/nv4/pecs/>

Generic Functions

Parametric Functions

- Just as we can add type parameters to a class definition, we can also add them to a function definition
- The type parameters are in scope in the header and body of the function

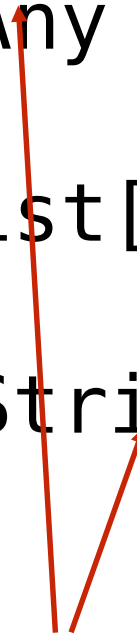
Map Revisited

```
abstract class List[+T] {  
  ...  
  def map[U](f: T => U): List[U]  
}
```

Why is this occurrence of T acceptable?

We Consider Specific Instantiations

```
abstract class List[Any] {  
  ...  
  def map[U](f: Any => U): List[U]  
}  
abstract class List[String] {  
  ...  
  def map[U](f: String => U): List[U]  
}
```



*Then List[String] is an acceptable subtype of List[Any]
provided that (String => U) >: (Any => U)
which requires that String <: Any.*

Generalizing Our Rules

- In our example, type parameter `T` occurs as the parameter of an arrow type:
 - `(String => U) >: (Any => U)`, provided:
 - `String <: Any`
 - `U <: U`
 - Consistent with `List[String] <: List[Any]`

An Example of How We Might Use Contravariant Type Parameters

```
abstract class Function1[-S,+T] {  
    def apply(x:S): T  
}
```

Map Revisited

```
case object Empty extends List[Nothing] {  
  ...  
  def map[U](f: Nothing => U) = Empty  
}
```

Map Revisited

```
case class Cons[+T](head: T, tail: List[T])
extends List[T] {
  ...
  def map[U](f: T => U) =
    Cons(f(head), tail.map(f))
}
```

Syntactic Sugar: Currying

- Scala provides special syntax for defining a function that immediately returns another function:

```
def f(x0: T0, ..., xN: TN) = (y0: U0, ..., yM: UM) => expr
```

can be rewritten as:

```
def f(x0: T0, ..., xN: TN)(y0: U0, ..., yM: UM) = expr
```

- Defining a function in this way is called “currying”, after the computer scientist Haskell Curry

Folding Revisited

```
abstract class List[+T] {  
  ...  
  def foldLeft[S](x: S)(f: (S, T) => S): S  
  def foldRight[S](x: S)(f: (T, S) => S): S  
}
```

Note that these functions are curried

