Comp 311 Functional Programming

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Announcement

Homework 2 was due today at 3pm

Midterm exam is on Tuesday Oct 29
in DH 1064 (next door) from 7pm – 10pm
(practice exam is on Piazza)

 Homework 3 is due on Thursday Nov 7 (posted to Piazza)

Scala Immutable Collections

Immutable Lists

- Behave much like the lists we have defined in class
- Lists are covariant
- The empty list is written Nil
- Nil extends List[Nothing]

Immutable Lists

The list constructor takes a variable number of arguments:

List(1,2,3,4,5,6)

Immutable Lists

 Non-empty lists are built from Nil and Cons (written as the right-associative operator ::)

1 :: 2 :: 3 :: 4 :: Nil

List Operations

- head returns the first element
- tail returns a list of elements but the first
- isEmpty returns true if the list is empty
- Many of the methods we have defined are available on the built-in lists

FoldLeft and FoldRight Written as Operators

• foldLeft:

```
(zero /: xs)(op)
```

foldRight:

```
(xs : \ zero)(op)
```

FoldLeft and FoldRight Written as Operators

• foldLeft:

(xs foldLeft zero)(op)

foldRight:

(xs foldRight zero)(op)

FoldLeft and FoldRight Written as Methods

• foldLeft:

```
xs.foldLeft(zero) { op }
```

foldRight:

```
xs.foldRight(zero) { op }
```

SortWith

```
List(1,2,3,4,5,6) sortWith (_ > _ )

\mapsto

List(6, 5, 4, 3, 2, 1)
```

Range

```
List.range(1,5)

→
List(1, 2, 3, 4)
```

Using Fill for Uniform Lists

```
List.fill(10)(0) \mapsto List(0,0,0,0,0,0,0,0,0)
```

Using Fill for Uniform Lists

```
List.fill(3,3)(0) →
List(List(0,0,0),
List(0,0,0))
```

Tabulating Lists

```
List.tabulate(3,3) { (m,n) =>
  if (m == n) 1 else 0
}

List(List(1,0,0),
      List(0,1,0),
      List(0,0,1))
```

Immutable Sets

Immutable Sets

- · Sets are unordered, unrepeated collections of elements
- Set[T] extends the function type T ⇒ Boolean
- Sets are parametric and invariant in their element type

Why *in-*variant?

Set Factory

Set(1,2,3,4,5)

Set Element Addition

Set(1,2,3) + 4
$$\mapsto$$
 Set(1,2,3,4)

Set Element Subtraction

Set Intersection

```
Set(1,2,3) intersect Set(2,4,5,3) → Set(2,3)
```

Set Union

Set Difference

Set Cardinality

Set Membership

```
Set(1,2,3).contains(2) → true
```

The *apply* method on sets is equivalent to the *contains* method.

Immutable Maps

Immutable Maps

- Maps are collections of key/value pairs
- They are parametric in both the key and value type
 - Covariant in their value type
 - Invariant in their key type



The -> Operator

• The infix operator -> returns a pair of its arguments:

 Note: Scala also allows *Unicode Operators*, and the infix "→" operator is one such example:

$$1 \rightarrow 2$$

$$\mapsto$$

$$(1,2)$$

The → Operator is Left Associative

```
> 1 \rightarrow 2 \rightarrow 3 \rightarrow 4
res8: (((Int, Int), Int), Int) = (((1,2),3),4)
```

The Map Factory

Map Addition

Map("a"
$$\rightarrow$$
 1, "b" \rightarrow 2, "c" \rightarrow 3) + ("d" \rightarrow 4) \mapsto Map(a -> 1, b -> 2, c -> 3, d -> 4)

Map Operations

The operators/methods are defined in the expected way:

- -
- ++
- - -
- size

Map Membership

```
Map("a" \rightarrow 1, "b" \rightarrow 2, "c" \rightarrow 3).contains("b") \mapsto true
```

Map Lookup

```
Map("a" → 1, "b" → 2, "c" → 3)("c")

\stackrel{\mapsto}{3}

Map("a" → 1, "b" → 2, "c" → 3).get("c")

\stackrel{\mapsto}{Some}(3)
```

Map Keys

 $Map("a" \rightarrow 1, "b" \rightarrow 2, "c" \rightarrow 3).keys$

```
Set(a, b, c): Iterable[String]

Map("a" \rightarrow 1, "b" \rightarrow 2, "c" \rightarrow 3).keySet

\mapsto

Set(a, b, c): Set[String]
```

Map Values

```
Map("a" → 1, "b" → 2, "c" → 3).values

\mapsto
Set(1,2,3)
```

Map Empty

Map("a" → 1, "b" → 2, "c" → 3).isEmpty
$$\mapsto$$
 false

Call-By-Value and Call-By-Name

- Thus far, the evaluation semantics we have studied (both with the substitution and environment models) is known as call-by-value:
 - To evaluate a function application, we first evaluate the arguments and then evaluate the function body

• We have seen several "special forms" where this evaluation semantics is not what we want:

 We could delay evaluation in these cases by wrapping arguments in function literals that take no parameters

```
def myOr(left: Boolean, right: () => Boolean) =
  if (left) true
  else right()
```

 We could delay evaluation in these cases by wrapping arguments in function literals that take no parameters

$$my0r(true, () => 1/0 == 2) \rightarrow true$$

 Functions that take no arguments are referred to as thunks

Call-By-Name

 Scala provides a way that we can pass arguments as thunks without having to wrap them explicitly

```
def myOr(left: Boolean, right: => Boolean) =
  if (left) true
  else right
```

We simply leave off the parentheses in the parameter's type

Call-By-Name

 Now we can call our function without wrapping the second argument in an explicit thunk:

$$my0r(true, 1/0 == 2) \rightarrow true$$

 The thunk is applied (to nothing) the first time that the argument is evaluated in a function

Call-By-Name

We can use by-name parameters to define new control abstractions:

```
def myAssert(predicate: => Boolean) =
  if (assertionsEnabled && !predicate)
  throw new AssertionError
```

Syntactic Sugar: Braces for Passing Arguments

 Any function that takes a single argument can be applied by passing the argument enclosed in braces instead of parentheses

```
myAssert {
   2 + 2 == 4
}
```

Syntactic Sugar: Braces for Passing Arguments

 Any function that takes a single argument can be applied by passing the argument enclosed in braces instead of parentheses

```
myAssert {
   def double(n: Int) = 2 * n
   double(2) == 4
}
```