# COMP 322: Fundamentals of Parallel Programming

# Lecture 14: Data-Driven Tasks and Data-Driven Futures

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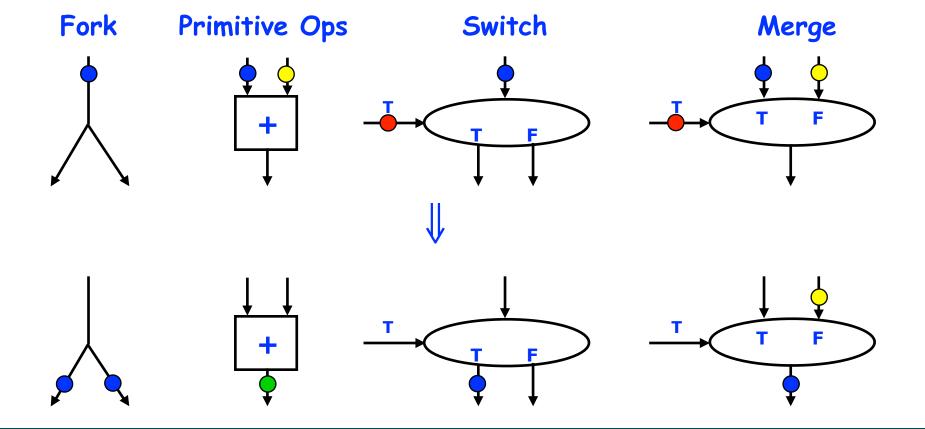
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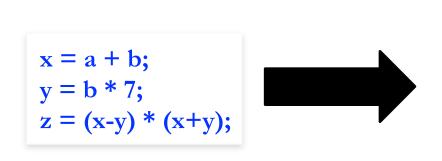
## **Dataflow Computing**

Original idea: replace machine instructions by a small set of dataflow operators

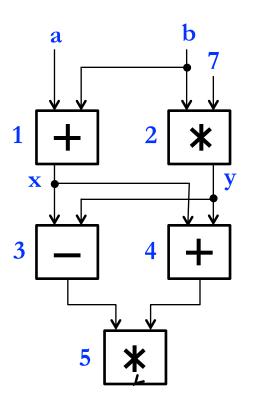




## Example instruction sequence and its dataflow graph



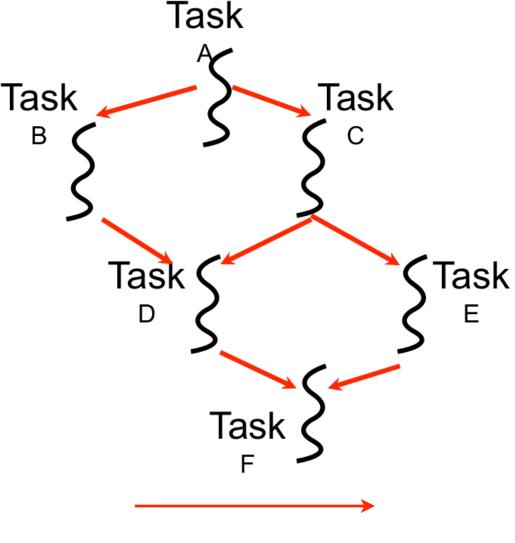
An operator executes when all its input values are present; copies of the result value are distributed to the destination operators.



No separate control flow



### **Macro-Dataflow Programming**



Communication via "single-assignment" variables

- "Macro-dataflow" = extension of dataflow model from instruction-level to task-level operations
- General idea: build an arbitrary task graph, but restrict all inter-task communications to single-assignment variables
  - Static dataflow ==> graph fixed when program execution starts
  - Dynamic dataflow ==> graph can grow dynamically
- Semantic guarantees: race-freedom, determinism
  - Deadlocks are possible due to unavailable inputs (but they are deterministic)



## Extending HJ Futures for Macro-Dataflow: Data-Driven Futures (DDFs) and Data-Driven Tasks (DDTs)

#### HjDataDrivenFuture<T1> ddfA = newDataDrivenFuture();

- Allocate an instance of a <u>data-driven-future</u> object (container)
- Object in container must be of type T1
- Used to implement "edges" in a computation graph

```
asyncAwait(ddfA, ddfB, ..., () -> Stmt);
```

- Create a new <u>data-driven-task</u> to start executing Stmt after all of ddfA, ddfB, ... become available (i.e., after task becomes "enabled")
- Used to implement "nodes" in a computation graph

```
ddfA.put(V) ;
```

- Store object V (of type T1) in ddfA, thereby making ddfA available
- Single-assignment rule: at most one put is permitted on a given DDF

```
ddfA.get()
```

- Return value (of type T1) stored in ddfA
- Throws an exception if put() has not been performed
  - Should be performed by async's that contain ddfA in their await clause, or if there's some other synchronization to guarantee that the put() was performed



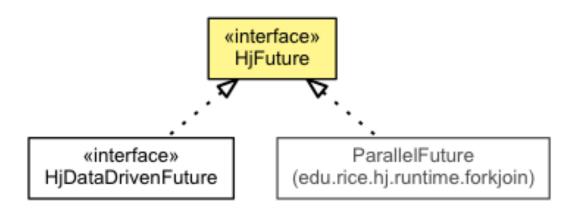
### Implementing Future Tasks using DDFs

#### Future version

#### DDF version



## HjFutures and HjDataDrivenFuture



- future.get()
  - Returns the value wrapped in the future.
- future.resolved()
  - Returns whether the future has been resolved, i.e. the value has been computed.
    - WARNING: use of resolved() can introduce nondeterminism



## Use of DDFs with dummy objects (like future<Void>)

```
В
1. finish(() -> {
2.
     HjDataDrivenFuture<Void> ddfA = newDataDrivenFuture();
3.
     HjDataDrivenFuture<Void> ddfB = newDataDrivenFuture();
4.
     HjDataDrivenFuture<Void> ddfC = newDataDrivenFuture();
5.
     HjDataDrivenFuture<Void> ddfD = newDataDrivenFuture();
6.
     HjDataDrivenFuture<Void> ddfE = newDataDrivenFuture();
7.
     async(() -> { ... ; ddfA.put(null); }); // Task A
8.
     asyncAwait(ddfA, () -> { ...; ddfB.put(null); }); // Task B
     asyncAwait(ddfA, () -> { ...; ddfC.put(null); }); // Task C
9.
10.
     asyncAwait(ddfB, ddfC, ()->{ ...; ddfD.put(null); }); // Task D
11.
     asyncAwait(ddfC, () -> { ...; ddfE.put(null); }); // Task E
12.
    asyncAwait(ddfD, ddfE, () -> { ... }); // Task F
13. }); // finish
```



### Differences between Futures and DDFs/ DDTs

- Consumer task blocks on get() for each future that it reads, whereas async-await does not start execution till all DDFs are available
- Future tasks cannot deadlock, but it is possible for a DDT to block indefinitely ("deadlock") if one of its input DDFs never becomes available
- DDTs and DDFs are more general than futures
  - Producer task can only write to a single future object, where as a DDT can write to multiple DDF objects
  - The choice of which future object to write to is tied to a future task at creation time, where as the choice of output DDF can be deferred to any point with a DDT
  - Consumer tasks can be created before the producer tasks
- DDTs and DDFs can be more implemented more efficiently than futures
  - An "asyncAwait" statement does not block the worker, unlike a future.get()



## Two Exception (error) cases for DDFs that do not occur in futures

- <u>Case 1:</u> If two put's are attempted on the same DDF, an exception is thrown because of the violation of the singleassignment rule
  - There can be at most one value provided for a future object (since it comes from the producer task's return statement)
- <u>Case 2:</u> If a get is attempted by a task on a DDF that was not in the task's await list, then an exception is thrown because DDF's do not support blocking gets
  - —Futures support blocking gets



### **Deadlock example with DDTs**

```
1. HjDataDrivenFuture left = newDataDrivenFuture();
2. HjDataDrivenFuture right = newDataDrivenFuture();
3. finish(() -> {
4.
     asyncAwait(left, () -> {
5.
        right.put(rightWriter()); });
6.
     asyncAwait(right, () -> {
7.
        left.put(leftWriter()); });
8. });

    HJ-Lib has deadlock detection mode

· Enabled using:
    - System.setProperty(HjSystemProperty.trackDeadlocks.propertyKey(), "true");
    - Reports an edu.rice.hj.runtime.util.DeadlockException when deadlock detected
```

