# COMP 322: Fundamentals of Parallel Programming 

## Lecture 12: Barrier Synchronization

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## Solution to Worksheet \#11: One-dimensional Iterative Averaging Example

1) Assuming $n=9$ and the input array below, perform a "half-iteration" of the iterative averaging example by only filling in the blanks for odd values of j in the myNew[] array (different from the real algorithm).
Recall that the computation is "myNew[j] = (myVal[j-1] + myVal[j+1])/2.0;"

| index, $\mathbf{j}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| myVal | 0 | 0 | 0.2 | 0 | 0.4 | 0 | 0.6 | 0 | 0.8 | 0 | 1 |
| myNew | 0 | 0.1 | 0.2 | 0.3 | 0.4 | 0.5 | 0.6 | 0.7 | 0.8 | 0.9 | 1 |

2) Will the contents of myVal[] and myNew[] change in further iterations? No, this represents the converged value (equilibrium/fixpoint).
3) Write the formula for the final value of myNew[i] as a function of $i$ and
n . In general, this is the value that we will get if m (= \#iterations in sequential for-iter loop) is large enough.
After a sufficiently large number of iterations, the iterated averaging code will converge with myNew[i] = myVal[i] = i / $(\mathrm{n}+1)$

## Hello-Goodbye Forall Example (Pseudocode)

forall (0, m-1, (i) -> \{
int sq = i*i; // NOTE: video used lookup(i) instead
System.out.println("Hello from task with square = " + sq);
System.out.println("Goodbye from task with square = " + sq);
\});

Sample output for $m=4$ :
Hello from task with square $=0$
Hello from task with square $=1$
Goodbye from task with square $=0$
Hello from task with square $=4$
Goodbye from task with square $=4$
Goodbye from task with square $=1$
Hello from task with square $=9$
Goodbye from task with square $=9$

## Hello-Goodbye Forall Example (contd)

```
forall ( \(0, \mathrm{~m}-1\), (i) \(->\{\)
    int \(\mathbf{s q}=\mathbf{i}^{\boldsymbol{*}} \mathbf{i}\);
    System.out.println("Hello from task with square = " + sq);
    System.out.printIn("Goodbye from task with square = " + sq);
\});
```

- Question: how can we transform this code so as to ensure that all tasks say hello before any tasks say goodbye?
- Statements in red below will need to be moved to solve this problem

Hello from task with square $=0$
Hello from task with square $=1$
Goodbye from task with square $=0$
Hello from task with square $=4$
Goodbye from task with square $=4$
Goodbye from task with square $=1$
Hello from task with square $=9$
Goodbye from task with square $=9$

## Hello-Goodbye Forall Example (contd)

```
forall (0, m-1, (i) -> {
int sq = i*i;
System.out.println("Hello from task with square = " + sq);
System.out.printIn("Goodbye from task with square = " + sq);
});
```

- Question: how can we transform this code so as to ensure that all tasks say hello before any tasks say goodbye?
- Approach 1: Replace the forall loop by two forall loops, one for the hello's and one for the goodbye's
—Problem: Need to communicate local sq values from first forall to the second 1.// APPROACH 1

2. forall ( $\mathbf{0}, \mathrm{m}-1$, (i) $)>\{$
3. int $\mathrm{sq}=\mathrm{i}^{\mathrm{*}} \mathrm{i}$;
4. System.out.println("Hello from task with square = " + sq);
5. \});
6. forall ( $0, m-1$, (i) -> $\{$
7. System.out.println("Goodbye from task with square $=\mathbf{~}+\mathrm{sq}$ );
8. \});

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## Hello-Goodbye Forall Example (contd)

- Question: how can we transform this code so as to ensure that all tasks say hello before any tasks say goodbye, without having to change local ?
- Approach 2: insert a "barrier" ("next" statement) between the hello's and goodbye's

1. // APPROACH 2
2. forallPhased (0, m-1, (i) -> \{
3. $\operatorname{int} \mathbf{s q}=\mathbf{i}^{\star} \mathbf{i}$;
4. System.out.println("Hello from task with square =" $\boldsymbol{+} \mathbf{s q}$ );
5. next(); // Barrier
6. System.out.println("Goodbye from task with square = " +sq );
7. \});

- next $\rightarrow$ each forallPhased iteration waits at barrier until all iterations arrive (previous phase is completed), after which the next phase can start
-Scope of next is the closest enclosing forallPhased statement
-If a forallPhased iteration terminates before executing "next", then the other iterations don't wait for it


## Impact of barrier on scheduling forallPhased iterations


next() operation is modeled in the Computation Graph using signal and wait edges


- static void forallPhased(int s0, int e0, edu.rice.hj.api.HjProcedure<java.lang.Integer> body)
- static $<T>$ void forallPhased(java.lang.Iterable $<T>$ iterable, edu.rice.hj.api.HjProcedure<T> body)
- static void next()
- NOTE:
-All forallPhased API's include an implicit finish at the end (just like a regular forall)
-Calls to next() are only permitted in forallPhased(), not in forall()


## Observation 1: Scope of synchronization for "next" barrier is its closest enclosing forallPhased statement

1. forallPhased ( $0, \mathrm{~m}-1$, (i) -> \{
2. println("Starting forall iteration " + i);
3. next(); // Acts as barrier for forallPhased-i
4. forallPhased (0, n - 1, (j) -> \{
5. println("Hello from task (" + i + "," + j + ")");
6. next(); // Acts as barrier for forallPhased-j
7. println("Goodbye from task (" + i + "," + j + ")");
8. \} // forallPhased-j
9. next(); // Acts as barrier for forallPhased-i
10. println("Ending forallPhased iteration " + i);
11.\}); // forallPhased-i

Observation 2: If a forall iteration terminates before "next", then other iterations do not wait for it

1. forallPhased (0, m - 1, (i) -> \{
2. forseq ( $0, i,(j)->\{$
3. // forall iteration $i$ is executing phase $j$
4. System.out.println("(" + i + ", " + j + ")");
5. next();
6. \});
7. \});

- Outer forall-i loop has $m$ iterations, 0...m-1
- Inner sequential j loop has i+1 iterations, 0...i
- Line 4 prints (task,phase) $=(i, j)$ before performing a next operation.
- Iteration $\mathbf{i}=0$ of the forall-i loop prints ( 0,0 ), performs a next, and then terminates. Iteration $i=1$ of the forall-i loop prints ( 1,0 ), performs a next, prints (1,1), performs a next, and then terminates. And so on.


## Barrier Matching for previous example

- Iteration $\mathrm{i}=0$ of the forallPhased-i loop prints ( 0 , 0 ) in Phase 0 , performs a next, and then ends Phase 1 by terminating.
- Iteration $i=1$ of the forallPhased-i loop prints $(1,0)$ in Phase 0 , performs a next, prints $(1,1)$ in Phase 1, performs a next, and then ends Phase 2 by terminating.
- And so on until iteration $i=8$ ends an empty Phase 8 by terminating


## Observation 3: Different forallPhased iterations may perform "next" at different program points

1. forallPhased (0, m-1, (i) $->$ \{
2. if (i \% $2==1$ ) $\{/ /$ i is odd
3. oddPhase0(i);
4. next();
5. oddPhase1(i);
6. \} else $\{/ /$ i is even
7. evenPhase0(i);
8. next(); $\qquad$
Barriers are not statically scoped - matching barriers may come from different program points, and may even be in different methods!
9. evenPhase1(i);
10. \} // if-else
11. \}); // forall

- Barrier operation synchronizes odd-numbered iterations at line 4 with even-numbered iterations in line 8
- One reason why barriers are "less structured" than finish, async, future

