Lecture 18: Abstract vs Real Performance - An “under the hood” look at HJlib

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HJ-lib Compilation and Execution Environment

Java 8 IDE

```
javac Foo.java
```

Java compiler translates Foo.java to Foo.class, along with calls to HJ-lib with lambda parameters (async, finish, future, etc)

```
java Foo
```

HJ-lib Runtime Environment = Java Runtime Environment + HJ-lib libraries

HJ-runtime initializes m worker threads (value of m depends on options or default value)

HJ-lib source program is a standard Java 8 program

All the “magic” happens here!

HJ Abstract Performance Metrics, HJ-Viz output (all enabled by appropriate options)
Looking under the hood — let’s start with the hardware

- Two quad-core Intel Xeon (CPUs), for a total of 8 cores/node (NOTS has 16 cores/node)
- Main Memory (DRAM)
Next, how does a process run on a single core?

Processes are managed by OS kernel

- Important: the kernel is not a separate process, but rather runs as part of some user process

Control flow passes from one process to another via a context switch

(e.g., Java application A) (e.g., Java application B)

Context switches between two processes can be very expensive!

Source: COMP 321 lecture on Exceptional Control Flow (Alan Cox)
What happens when executing a Java program?

- A Java program executes in a single Java Virtual Machine (JVM) process with multiple threads.
- Threads associated with a single process can share the same data.
- Java main program starts with a single thread (T1), but can create additional threads (T2, T3, T4, T5) via library calls.
- Java threads may execute concurrently on different cores, or may be context-switched on the same core.

Figure source: COMP 321 lecture on Concurrency (Alan Cox)
Thread-level Context Switching on the same processor core

- Thread context switch is cheaper than a process context switch, but is still expensive (just not “very” expensive!)
- It would be ideal to just execute one thread per core (or hardware thread context) to avoid context switches

Figure source: COMP 321 lecture on Concurrency (Alan Cox)
Now, what happens in a task-parallel Java program (e.g., HJ-lib, Java ForkJoin, etc)

<table>
<thead>
<tr>
<th>HJ-Lib Tasks &amp; Continuations</th>
</tr>
</thead>
<tbody>
<tr>
<td>Worker threads</td>
</tr>
<tr>
<td>Operating System</td>
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<tr>
<td>Hardware cores</td>
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</table>

- HJ-lib runtime creates a *small number of worker threads*, typically one per core
- Workers push new tasks and "continuations" into a logical work queue
- Workers pull task/continuation work items from logical work queue when they are idle (remember greedy scheduling?)

![Diagram of logical work queue and worker threads]

Local variables are *private to each task*

Static & instance fields are *shared among tasks*
Task-Parallel Model: Checkout Counter Analogy

- Think of each checkout counter as a processor core

Task-Parallel Model: Checkout Counter Analogy

- Think of each checkout counter as a processor core
- And of customers as tasks

source: http://www.deviantart.com/art(Randomness-20-178737664)
All is well until a task blocks …

• A blocked task/customer can hold up the entire line
• What happens if each checkout counter has a blocked customer?

source: http://viper-x27.deviantart.com/art/Checkout-Lane-Guest-Comic-161795346
Approach 1: Create more worker threads (as in HJ-Lib’s Blocking Runtime)

- Creating too many worker threads can exhaust system resources (OutOfMemoryError), and also leads to context-switch overheads when blocked worker threads get unblocked.
• Assume that five tasks (A1 … A5) are registered on a barrier

• Q: What happens if four tasks (say, A1 … A4) executing on workers w1 … w4 all block at the same barrier?

• A: Deadlock! (All four tasks will wait for task A5 to enter the barrier.)

• Blocking Runtime’s solution to avoid deadlock: keep task blocked on worker thread, and create a new worker thread when task blocks
Blocking Runtime (contd)

- Examples of blocking operations
  - End of finish
  - Future get
  - Barrier next

- Approach: Block underlying worker thread when task performs a blocking operation, and launch an additional worker thread

- Too many blocking operations can result in exceptions and/or poor performance, e.g.,
  - java.lang.IllegalStateException: Error in executing blocked code! [89 blocked threads]
  - Maximum number of worker threads can be configured if needed
    - HjSystemProperty.maxThreads.set(100);
Approach 2: Suspend task continuations at blocking points (as in HJ-Lib’s Cooperative Runtime)

- Upon a blocking operation, the currently executing tasks suspends itself and yields control back to the worker.
- Task’s *continuation* is stored in the suspended queue and added back into the ready queue when it is unblocked.
- **Pro:** No overhead of creating additional worker threads.
- **Con:** Need to create continuations (enabled by -javaagent option).

Continuations

- A continuation is the point immediately following a blocking operation, such as an end-finish, future get(), barrier/phaser next(), etc.

- Continuations are also referred to as task-switching points—Program points at which a worker may switch execution between different tasks (depends on scheduling policy)

```java
1. finish { // F1
2.   async A1;
3.   finish { // F2
4.     async A3;
5.     async A4;
6.   }
7.   S5;
8. }  
```

Continuations
Cooperative Scheduling (view from a single worker)

Task-1

block

unblock

block

... 

Task-1

suspend

resume

suspend

... 

Task-2

suspend/completed

Useful work for some other task on same worker thread

Cooperative runtime automatically creates continuations at suspend points via bytecode instrumentation enabled by -javaagent option
Any operation that contributes to unblocking a task can be viewed as an event e.g., task termination in finish, return from a future, signal on barrier, put on a data-driven-future, ...
Why are Data-Driven Tasks (DDTs) more efficient than Futures?

- Consumer task blocks on get() for each future that it reads, whereas async-await does not start execution till all Data-Driven Futures (DDFs) are available
  — An “asyncAwait” statement does not block the worker, unlike a future.get()
  — No need to create a continuation for asyncAwait; a data-driven task is directly placed on the Suspended queue by default
- Therefore, DDTs can be executed on a Blocking Runtime without the need to create additional worker threads, or on a Cooperative Runtime without the need to create continuations
Summary: Abstract vs. Real Performance in HJlib

- **Abstract Performance**
  - Abstract metrics focus on operation counts for WORK and CPL, regardless of actual execution time

- **Real Performance**
  - HJlib uses ForkJoinPool implementation of Java Executor interface with Blocking or Cooperative Runtime (option-controlled)
Announcements & Reminders

- HW3 CP 1 is available and due today by 11:59pm
- Watch the topic 4.1, 4.4 videos for the next lecture
- Use Piazza (public or private posts, as appropriate) for all communications re. COMP 322
- See Office Hours link on course web site for latest office hours schedule.