COMP 322: Parallel and Concurrent Programming

Lecture 4: Lazy Computation

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Lazy computation

Lazy evaluation

Logging library
Log.*i*(**TAG**,   **"current input: "** + input.toString()); *// eager*
Log.*i*(**TAG**, ()->**"current input: "** + input.toString()); *// lazy*

Why lazy?

Maybe it's expensive to compute something.

Maybe you won't actually need it (e.g., if you disabled logging).

The idea: defer computation of a value until (and if) you need it.

Related idea: Memoization

The idea: compute a value <u>once</u> when you need it, then save it.

Deep, powerful idea in computer science (e.g., dynamic programming).

Lazy memo implementation (simplified)

```java
public class Lazy<T> {
  private T contents;
  private Supplier<T> supplier;

  private Lazy(Supplier<T> supplier) {
    contents = null;
    this.supplier = supplier;
  }

  public T get() {
    if (contents != null) {
      return contents;
    }

    if (supplier != null) {
      contents = supplier.get();
      supplier = null;
    }

    return contents;
  }
}
```

Private constructor (as usual) plus a factory method (Lazy.*of*)

If we've already computed the answer, return it.

Call the lambda (once), get the result, forget the lambda.

Lazy lists

A <u>lazy</u> list is:

A head value

A lambda returning another lazy list ("tail-value function" or "tail supplier")

Or:

An Empty list

```
static <T> LazyList<T> cons(T head, Supplier<LazyList<T>> tailSupplier) {
    return new LazyCons<>(head, tailSupplier);
}
```

Implementing lazy lists (simplified)

```java
class LazyCons<T> implements LazyList<T> {
    final T head;
    final Lazy<LazyList<T>> tail;

    LazyCons(T head, Supplier<LazyList<T>> tail) {
        this.head = head;
        this.tail = Lazy.of(tail);
    }

    public T head() {
        return head;
    }

    public LazyList<T> tail() {
        return tail.get();
    }
}
```

a lambda that will return the tail

Build a memo around the tail supplier so that we only call the lambda once

tail() hides the implementation details

The payoff? Infinite lists!

```java
// Make a LazyList of integers starting from i, skipping by step
public static LazyList<Integer> from(int i, int step) {
  return cons(i,()->from(i+step, step));
}

// Make a LazyList consisting of all the same elements
public static <T> LazyList<T> continually(T s) {
  return cons(s,() -> continually(s));
}

    var wholeNumbers = from(0, 1); // 0, 1, 2, 3…, runs quickly

    var evens = wholeNumbers.filter(x -> x % 2 == 0); // 0, 2, 4, 6…, runs quickly

    var squares = wholeNumbers.map(x -> x * x); // 0, 1, 4, 9…, runs quickly

    var zeros = continually(0); // 0, 0, 0, 0…, runs quickly

    var alsoEvens = from(0, 2); // runs quickly

    var yetAnotherEvens = wholeNumbers.map(x -> x * 2); // runs quickly
```

But be careful…

var evens = wholeNumbers.filter(x -> x % 2 == 0); *// 0, 2, 4, 6...*
var alsoEvens = *from*(0, 2);

assertEquals(evens, alsoEvens); *// never finishes!*

You can't do any operation that requires the entire list!

No length of an infinite list (will never terminate).

You can't fold an infinite list (will never terminate).

You can't test list equality (.equals() will go forever as well).

But other operations are just fine.

map, filter, etc.: run in *constant time*, return a new lazy list.

And if you *take()* a finite number of elements from an infinite list, you can do anything with it

fold, length, equality, etc.

Lazy Filter

```java
class LazyCons<T> implements LazyList<T> {
  public LazyList<T> filter(Predicate<T> predicate) {
    if (predicate.test(headVal)) {
      return cons(headVal, () -> tail().filter(predicate));
    } else {
      return tail().filter(predicate);
    }
  }
}
```

Lazy Map

```java
class LazyCons<T> implements LazyList<T> {
  public <R> LazyList<R> map(Function<T, R> f) {
    return cons(f.apply(headVal), ()->tail().map(f));
  }
}
```

How about *take()*? That's lazy too!

```java
class LazyCons<T> implements LazyList<T> {
  public LazyList<T> take(int n) {
    if (n < 1) {
      return empty();
    } else if (n == 1) {
      return cons(headVal, ()-> empty());
    } else {
      return cons(headVal, () -> tail().take(n - 1));
    }
  }
}
```

fold() cannot be lazy

```
public <U> U foldRight(U zero, BiFunction<T, U, U> operator) {
  return operator.apply(headVal, tail().foldRight(zero, operator));
}
```

Exact same implementation as *GList*

fold() is a ***terminal*** operation

When to be lazy? When to be eager?

- Laziness almost always wins (in big-*O*)

  - But the memoization does have a cost.

    - If your lists have millions of entries, this starts to matter.

    - But at that point, maybe you shouldn't be using lists.

- Some programming languages are *extremely* lazy (Haskell).

  - No value ever computed until it's ultimately needed.

  - Yet still your computation runs efficiently.

- Java is only lazy when you explicitly use lambdas.

Summary

Sometimes you want to be lazy

- If it's too expensive to be eager

- If it's possible that you'll never need the value

- If it's hard to keep track of which values you have already computed

Defer the computation to when (and if) you need it

Very easy to do in functional programming

Map, take, etc. run in ***constant*** time.

Filter is a bit more complicated, but runs in constant time most of the time

Allows you to create logically infinite data structures

Need to be careful with ***terminal*** operations (fold, length, etc.) on infinite data