COMP 322: Parallel and Concurrent Programming

Lecture 29: Dining Philosophers

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Liveness Recap

- Deadlock: task's execution remains incomplete due to it being blocked awaiting some condition
- Livelock: two or more tasks repeat the same interactions without making any progress
- Starvation: some task is repeatedly denied the opportunity to make progress
- Bounded wait (fairness): each task requesting a resource should only have to wait for a bounded number of other tasks to "cut in line"
- Non-concurrency: a task is prevented from making progress due to overly restrictive resource management

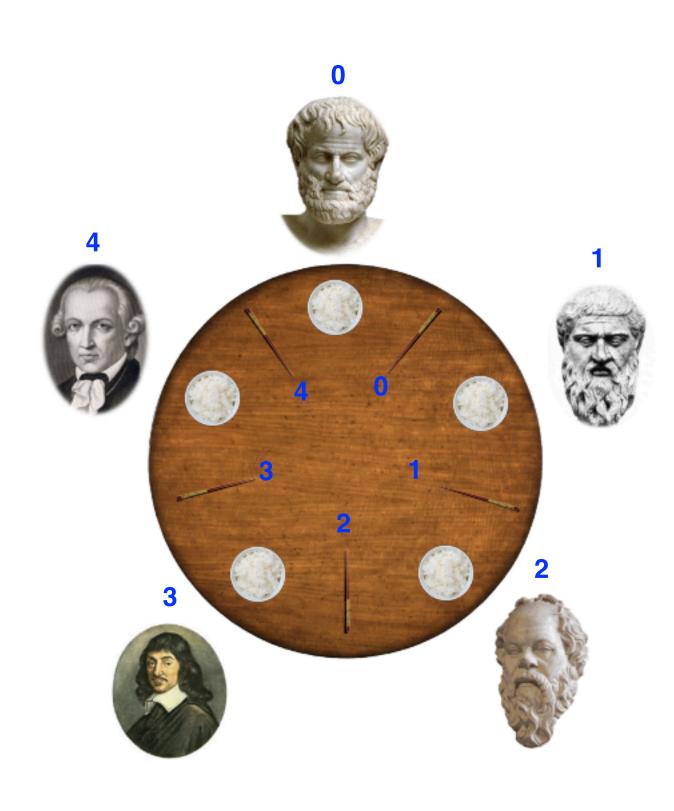


Deadlock Conditions

- Mutual Exclusion
 - At least one resource that must be held is in non-shareable mode
- Hold and wait
 - There exists a task holding a resource, and waiting for another
- No preemption
 - Resources cannot be preempted
- Circular wait
 - There exists a set of tasks {T₁, T₂, ... T_N}, such that
 - T_1 is waiting for T_2 , T_2 for T_3 , and T_N for T_1
- All four conditions must hold for deadlock to occur



The Dining Philosophers Problem



A classical Synchronization Problem devised by Dijkstra in 1965 Constraints

- Five philosophers either eat or think
- They must have two chopsticks to eat
- Can only use chopsticks on either side of their plate
- No talking permitted

Goals

- Progress guarantees
 - Deadlock freedom
 - Livelock freedom
 - Starvation freedom
 - Maximum concurrency (no one should starve if there are available forks for them)



General Structure of Dining Philosophers Problem: PseudoCode

```
1. int numPhilosophers = 5;
2. int numChops = numPhilosophers;
3. Chop[] chop = ...; // Initialize array of chopsticks
4. for(p in 0 .. numPhilosophers-1) {
   async(() \rightarrow \{
6.
     while(true) {
       Think ;
8.
      Acquire chopsticks;
9.
      // Left chopstick = chop[p]
10.
       // Right chopstick = chop[(p-1)%numChops]
        Eat;
11.
     } // while
13. }); // async
14.} // for
```



Solution 1: Using Java's Synchronized Statement

```
1. int numPhilosophers = 5;
2. int numChops = numPhilosophers;
3. Chop[] chop = ...; // Initialize array of chopsticks
4. for(p in 0 .. numPhilosophers-1) {
   async(() \rightarrow \{
6.
     while(true) {
7.
       Think ;
8.
       synchronized(chop[p]) { // get the left chopstick
        synchronized(chop[(p-1)%numChops]) { // get the right chopstick
9.
10.
          Eat;
11.
12.
13.
     } // while
14. }); // async
15.} // for
```



Problems?

- What if everyone picks up the left chopstick at the same time?
- Deadlock!
- Starvation due to deadlock
- No livelock
- Non-concurrency due to deadlock



Solution 2: Using Java's tryLock

```
1. int numPhilosophers = 5;
2. int numChops = numPhilosophers;
3. Chop[] chop = ...; // Initialize array of chopsticks
4. for(p in 0 .. numPhilosophers-1) {
   async(() \rightarrow \{
     int first = p; int second = (p - 1) % numChops;
     while(true) {
8.
       Think;
9.
      if (!chop[first].lock.tryLock()) continue;
10.
       if (!chop[second].lock.tryLock()) {
11.
        chop[first].lock.unLock(); continue;
12.
13.
       Eat;
       chop[first].lock.unlock();chop[second].lock.unlock();
14.
     } // while
16. }); // async
17.} // for
```



Problems?

- Everyone picks up the left chopstick at the same time, tries to pick up the right one, gives up, puts down the left one, and repeat
- Livelock!
- Starvation due to livelock!
- No deadlock
- Non-concurrency due to livelock



Solution 3: Using Global Isolated

```
1. int numPhilosophers = 5;
2. int numChops = numPhilosophers;
3. Chop[] chop = ...; // Initialize array of chopsticks
4. for(p in 0 .. numPhilosophers-1) {
   async(() -> {
     while(true) {
6.
       Think;
8.
      isolated {
9.
        Pick up left and right chopsticks;
10.
        Eat;
11.
12.
     } // while
13. }); // async
14.} // for
```



Problems?

- No deadlock or lovelock possible
- Starvation!
 - No guarantee that a philosopher will ever get to eat, if others are very hungry and "cut in line" all the time.
- Non-concurrency
 - Only one philosopher can eat at any time



Solution 4a: Impose Order

```
1. int numPhilosophers = 5;
2. int numChops = numPhilosophers;
3. Chop[] chop = ...; // Initialize array of chopsticks
4. for(p in 0 .. numPhilosophers-1) {
5. async(() -> {
     int first = (p == 0)? (p - 1) % numChops : p
     int second = (p == 0)? p : (p - 1) % numChops
     while(true) {
8.
      Think;
9.
10.
       synchronized(chop[first]) {
11.
        synchronized(chop[second]) {
12.
         Eat;
13.
14.
    } // while
16. }); // async
17.} // for
```



Preventing Deadlock by Ordering

It is not possible for all philosophers to have a chopstick

- Two philosophers, A and B, must share a chopstick, X, that is "bigger" than all other chopsticks
- 2. One of them, A, has to pick up X first
- 3. B can't pick up X at this point
- 4. B can't pick up the "smaller" chopstick until X is released
- 5. SO, 4 philosophers left, 5 chopsticks total
- 6. One philosopher must be able to have two chopsticks!



Solution 4b: Using tryLock

```
1. int numPhilosophers = 5;
2. int numChops = numPhilosophers;
3. Chop[] chop = ...; // Initialize array of chopsticks
4. for(p in 0 .. numPhilosophers-1) {
5. async(() -> {
    int first = (p == 0)? (p - 1) % numChops : p
    int second = (p == 0)? p : (p - 1) % numChops
8.
     while(true) {
9.
       Think;
10.
       if (!chop[first].lock.tryLock()) continue;
11.
       if (!chop[second].lock.tryLock()) {
12.
        chop[first].lock.unLock(); continue;
13.
       Eat;
14.
15.
       chop[first].lock.unlock();chop[second].lock.unlock();
     } // while
16.
17. }); // async
18.} // for
```



Solution 4c: Using Object-Based Isolation

```
1. int numPhilosophers = 5;
2. int numChops = numPhilosophers;
3. Chop[] chop = ...; // Initialize array of chopsticks
4. for(p in 0 .. numPhilosophers-1) {
   async(() \rightarrow \{
6.
     while(true) {
      Think;
8.
      isolated (chop[p], chop[(p-1)%numChops){
9.
        Eat;
10.
11. } // while
12. }); // async
13.} // for
```



Problems for 4a, 4b and 4c?

- No deadlock or lovelock possible
- Starvation!
 - No guarantee that a philosopher will ever get to eat, if others are very hungry and "cut in line" all the time.
- Concurrency
 - 4a: still have a non-concurrency problem. If philosopher 0 is eating, philosophers 1-3 could all be holding their left chopstick waiting
 - 4b and 4c: If a philosopher is hungry, and his chopsticks are not used for eating, he'll get to eat



Solution 5: Using Semaphores

"true" parameter creates a

semaphore that guarantees

```
1. int numPhilosophers = 5;
2. int numChops = numPhilosophers;
3. Chop[] chop = ...; // Initialize array of chopsticks
4. Semaphore table = new Semaphore(3, true);
5. for (i=0;i<numChops;i++) chop[i].sem = new Semaphore(1, true);
6. for(p in 0 .. numPhilosophers-1) {
7. async(() -> {
    while(true) {
9.
      Think;
10.
       table.acquire(); // At most 3 philosophers at table
11.
       p = empty place at the table that has nobody on the left
12.
       chop[p].sem.acquire(); // Acquire left chopstick
13.
       chop[(p-1)%numChops].sem.acquire(); // Acquire right chopstick
14.
       Eat;
15.
       chop[p].sem.release(); chop[(p-1)%numChops].sem.release();
16.
       table.release();
     } // while
18. }); // async
19.} // for
```

