Scheme Primitives and Function Definitions

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Course Overview

- Functional program design in Scheme
  - Data-directed program design 2-10
  - Algorithm design 11-14
  - Applied functional programming 15-17
- OO program design in Java 18-45
  - ...

Outline of This Lecture

- Common basic types
- Common primitive operations
- Rules for reducing programs
- Simple programs =

  Variable definitions (Constants) +
  Function definitions
- The design recipe
- Errors
- Data definitions

Basic (primitive) types of data

1. numbers:

   naturals: **`0, 1, 2, …`** // number theory

   integers: **`…, -1, 0, 1, …`** // include negatives

   rational numbers: **`3/4, 0, -1/3, …`** // include fractions

   inexact numbers: **`#i0.123, #i0, …`** // floating point

   Operations: **`+`**, **`-`**, **`*`**, **`/`**, **`expt`**, **`remainder`**

   Scheme computes exact answers on exact inputs if possible

2. booleans: **`false`**, **`true`**
   Operations: **`not`**, **`and`**, **`or`**, …

3. Symbols: **`'A, 'a, 'Aa, 'Corky, …`**
   Operations: … // none important for now

4. Other basic types: strings, lists , … // none important now

Mixed-type Operations and Primitive Computation

- Basic relational operators
  - equal? // all data values
  - =, <, >, <=, >= // only on numbers
- Primitive computation ≡ application of a basic operation to constants (or primitive computations)
  - Basic operation ≡ basic function
  - Soon, we will see how to define our own (non-primitive) functions
- Function application in Scheme: parenthesized prefix notation
  - Scheme uses parenthesized prefix notation uniformly for *everything*
  - `(+ 2 2)`, `(sqrt 25)`, `(remainder 7 3)`
  - Bigger example: `(* (+ 1 2) (+ 3 4))`
  - How does this compare to writing `1+2*3+4` ?
- Scheme syntax is simple, uniform, and avoids potential ambiguity

Computation Is Repeated Reduction

- *Every Scheme computation is the evaluation of a given expression constructed from primitive or defined functions and variables (names for constants).*
- *Evaluation proceeds by repeatedly performing the leftmost possible reduction (simplification) until the resulting expression is a* ***value.***
- *A* ***value*** *is any constant.* We will identify all of the expressions that are values as we explicate the language. Numbers, booleans, symbols are all values.

Reduction for primitive functions

- A *reduction* is an atomic step in a computation that replaces some expression by a simpler expression as specified by a Scheme evaluation rule (law). Every application of a basic operation to values yields a value (where run-time error is a special kind of value).
- Example

```
(* (+ 1 2) (+ 3 4))
=> (reduces to) (* 3 (+ 3 4))
=>  (* 3 7) =>  21
```

- Always perform leftmost reduction
- The following is **not** an atomic step, and so **not** a reduction

```
(- (+ 1 3) (+ 1 3)) =>  0
```

Example cont.

```
     (- (+ 1 3) (+ 1 3))
=>   (- 4 (+ 1 3))
=>   (- 4 4)
=>   0
```

Programs = Variable Definitions + Function Definitions

- Variables are simply names for values (constants)
  - `pi, my-SSN, album-name, tax-rate, x`
- Variable definitions Have the form:
  - `(define freezing 32)`
  - `(define boiling 212)`
- Function definitions have the form:
  - `(define (area-of-box x) (* x x))`
  - `(define (half x) (/ x 2))`
- Function applications (just as we saw before) have the form:
  - `(area-of-box 2)`
  - `(half (area-of-box 3))`
- Almost **any** function f used in a program can be written in the form
  - `(define (f v1 … vn) <expression>)`

  where <*expression*> is constructed from constants, variables `v1, … vn`, function applications, and a few other constructs TBN.

Reductions for defined functions

- Assume we have declared the two functions
  - `(define (area-of-box x) (* x x))`
  - `(define (half x) (/ x 2))`
- Then Scheme can perform these reductions

```
   (half (area-of-box 3))
=> (half (* 3 3))
=> (half 9)
=> (/ 9 2)
=> 4.5
```

- Reduction stops when we get to a value or an error

Example: Solve quadratic equation

```
;; Contract solve-quadratic: number number number -> number
;; Purpose: (solve-quadratic a b c) finds the larger root of
   a*x*x + b*x + c = 0 given it has real roots and a != 0

;; Examples: (solve-quadratic 1 0 -25) = 5
;;           (solve-quadratic 5 0 -20) = 2
;;           (solve-quadratic 1 -10 25) = -4
;;           . . . and other examples

;; Template instantiation:  (degenerate)
;; (define (solve-quadratic a b c) ... )

;; Code
;; (define (solve-quadratic a b c)
;;   (/ (+ (- b) (sqrt (- (* b b) (* 4 a c)))) (* 2 a)))

;; Tests for solve-quadratic
;; (check-expect exp ans) reports error if exp != ans
   (check-expect (solve-quadratic 1 0 -25) 5)
   (check-expect (solve-quadratic 5 0 -20) 2)
   (check-expect (solve-quadratic 1 -10 25) -4)
```

Step 2

Step 3

Step 4

Step 5

Step 6

Syntax Errors

- A syntactically correct **expression** can be
  - An *atomic* expression, like
    - a number `17, 4.5, #i0.34`
    - a variable `radius`
  - A *compound* **expression**,
    - starting with `(`
    - followed by basic or program-defined operation such as `+` or `f`
    - one or more **expression**s separated by spaces
    - ending with **)**
- Syntax errors: `3 + 4` `+(3,4)` `3)` `(5`

Runtime Errors

- Happen when basic operations are applied to illegal arguments
- Consider the following examples:
  - `(sqrt 1 2 3 4) => error: sqrt applied to more than one argument`
  - `(18 17) => error: 18 applied as function  ;;`
  - `(/ 1 0) => error: division by zero`
  - `(+ 1 'a) => error: second argument in application of + is not a number`
- If a reduction produces an error, the computation is aborted and the error is returned as the result.
- Try things like that in DrScheme, and make a mental note of the error messages you get back.

Conditional Expressions

- An expression that distinguishes different forms of data
- Form:

```
(cond [question-1  result-1]
      [question-2  result-2]
      ...
      [question-n  result-n]
      [else        default-result])
```

- Square brackets are used above for clarity. In Scheme, they are synonymous with parentheses, but balancing brackets must match.
- **`else`** is optional. If omitted and none of the questions are true, the result is a run-time error (like division by zero).

Reduction of Conditional Expressions

```
   (cond [true    result]
         ... )
=> result

   (cond [false   result]
         ... )
=> (cond  ... )

   (cond [else    result])
=> result

   (cond [false result])
=> error: all question results were false
```

Conditional Expression Examples

```
    (cond [(> 12 0) 5] [else -5])
=> (cond [true 5] [else -5])
=> 5
```

Given

```
(define (abs x)
   (cond [(>= x 0) x]
         [else (- x)]))

    (abs -10)
=> (cond [(>= -10 0) -10] [else (- -10)])
=> (cond [false -10] [else (- -10)])
=> (cond [else (- -10)]) => (- -10) => 10
```

The Design Recipe

How should I go about writing programs?

1. Analyze problem and define any requisite data types.
2. State contract (type) and purpose for *function* that solves the problem.
3. Give examples of function use and result.
4. Select and instantiate a template for the function body.
5. Write the function itself.
6. Test it, and confirm that tests succeeded.

The order of the steps of the recipe is important

The Design Recipe (Big Picture)

- Encourages systematic problem solving
- Works best if keep our functions small
- We will learn how to repeatedly decompose problems into simpler problems until we reach problems that can be solved by simple expressions as in **`solve-quadratic`**
- Decomposition driven by structure of data being processed: *data-directed* design

Reminders

- New homework (HW1) is posted online
  - Due next Friday, so you will get to check it over in lab; don't wait until your lab to get started.
  - Sign up for mailing list to get any updates, discussions
  - Make absolutely sure you follow the **recipe** in writing Scheme programs.
  - Partners: Talk to people after class, at lab, etc.
  - For Scheme programs, follow format of the sample solution in the Scheme HW Guide.
  - For hand evaluations, follow the format of the hand evaluation problems posted in the Scheme HW Guide.
  - Submit your assignment using Owlspace.

Epilog

- Reminder: continue digesting chs. 1-10 in HTDP Section 8.3 is particularly important and it is not wordy.
- Next class
  - *Inductive Data* definitions
  - Amplified design recipe
- Challenge problem: What happens if we use rightmost reduction instead of leftmost? Can you devise a program using the Scheme subset given in this lecture such that some invocation of that program (expression composed from constants and and basic and program-defined operations defined in the program) behaves differently (either in terms the result produced by the computation or lack thereof) under rightmost evaluation than leftmost evaluation. Hint: focus on pathological behavior and note that two different errors are not equivalent.