Data definitions and Conditionals

Prof. Robert “Corky” Cartwright

Dr. Stephen Wong

Rice University

Recap of Previous Lecture

- Primitive types and values
  - numbers, booleans, symbols
- Variable definitions (constants), function definitions
- Operators
  - Arithmetic, relational, function application
- Rules for reducing programs
  - Leftmost reduction
- Syntax Errors & Runtime Errors
- Conditional Expressions

Challenge Problem from Previous Lecture

Can you think of a Scheme program that exibits different behaviors with rightmost reduction instead of leftmost?

Consider the following example:

```
(+ (/ 1 0) (+ 'A 12))
```

Error conditions can make reasoning about programs more difficult than in naive mathematics *e.g.*, you may not preserve program behavior by replacing **`(* 0 (f x))`** by **`0`**. What if evaluating **`(f x)`** generates an error?

Another example: given a non-terminating function **`omega`**:

```
(+ (/ 1 0) (omega 0))
```

Goals of this lecture

- Defining compound data (Scheme *structs*)
- Template for processing *structs*
- Union (mixed) data definitions
- Conditionals
- Template for processing union data.
- Inductive (self-referential) data definitions
- Template for processing inductive data

Simple Data Definitions

- How do we define new forms of data in Scheme? For example, say we want to write a program for the registrar that maintains a directory of courses that can be searched …
- Informal rose description

  “A complex number is a pair with a real part and an imaginary part , which are both numbers ”
- Corresponding data definition in Scheme

```
;; Complex is a structure (make-Complex real imag)
;; where real and imag are numbers
;; NOTE: the type complex is primitive in Scheme so we
;;       capitalize the name to avoid syntax errors
(define-struct Complex (real imag))
```

- A Scheme *struct* is a tuple tagged with the struct name
- Scheme processes this definition by creating the following operations:
  - *constructor:* **make-Complex**,
  - *accessors:* **Complex-real**, **Complex-imag**
  - *recognizer:* **Complex?** (which checks the tag)

Structs Can Represent Compound Data

In the struct definition

```
(define-struct Complex (real imag))
```

`real` and `imag` are called **fields**.

If a struct has more than one field, it is a compound form of data because it more than one internal part. A struct with *k* fields can be thought of as a box with *k* compartments where each compartment is labeled with a distinct field name.

For example, the struct `Complex` has two fields (compartments) called `real` and `imag`.

Operations on Structures

Recall that the following operations are automatically generated from the **define-struct** declaration for **`Complex`**

- *constructor:* **`make-Complex`**
- *accessors:* **`Complex-real`**, **`Complex-imag`**
- *recognizer:* **`Complex?`**

Sample reductions for these field accessors and structure recognizers

```
(Complex-imag (make-Complex 1 2)) => 2
(Complex? (make-Complex 3 4)) => true
```

Structures are values

- In a program, the structure returned by a constructor is a value and its parts are values.
- Inside a structure, the parts *must* be values. The application of a struct constructor like **`make-Complex`** to some argument expressions evaluates these arguments to produce *values*. At this point, the struct application becomes a *value* because all of the parts in its compartments are values.
- For example:

  `(make-Complex 0 (+ 2 2))` is a constructor application--not a value because the argument expression **`(+ 2 2)`** is not a value.

  `(make-Complex 0 (+ 2 2)) => (make-Complex 0 4)`

  `(make-Complex 0 4)` is a *value*. Why?

  `(make-Complex x y)` is *not* a *value*. Why?

Evaluation Rules for Structures

Given the data definition

```
(define-struct Complex (real imag))
```

Scheme supports the following reduction rules:

```
(Complex-real (make-Complex Val1 Val2)) => Val1
(Complex-imag (make-Complex Val1 Val2)) => Val2
(Complex? (make-Complex Val1 Val2)) => true
(Complex? Val3) => false
```

where `Val1 Val2` are Scheme values and `Val3` is not of the form `(make-Complex V1 V2)`

The Design Recipe

How should I go about writing programs?

1. Analyze problem and **define** any requisite **data** types
2. State the **type contract** and **purpose** for *function* that solves the problem
3. Give **examples** of function use and result
4. Select and **instantiate** a **template** for the function body
5. Write the **code** for the function.
6. **Test** the code, and confirm that tests succeeded

The order of the steps of the recipe is important. In DrScheme, steps 3 and 6 can be collapsed because the examples can be presented as calls on **`check-expect`.** DrScheme does not evaluate these tests until the end of the program text.

Template for Defined Data Type

- We start from the data definition. Example:

```
;; A Complex is a structure (make-Complex real imag)
;; where real and imag are numbers
(define-struct Complex (real imag))
```

- General template for <u>any</u> function <u>processing</u> an argument of type **Complex**

```
;; (define (f c)
;;       ... (Complex-real c) ...
;;       ... (Complex-imag c) ...)
```

- Type contracts for some possible functions on **Complex**

```
;; magnitude : Complex -> number
;; Complex-zero? : Complex -> bool
;; conjugate: Complex -> Complex
```

Example: write conjugate function

Assume that we have already defined the **`Complex`** type include a template for functions that process inputs of type **`Complex`**.

```
;; Type contract
;; conjugate: Complex -> Complex

;; Examples:
(check-expect (conjugate (make-Complex 0 0)) (make-Complex 0 0))
(check-expect (conjugate (make-Complex 0 1)) (make-Complex 0 -1))
(check-expect (conjugate (make-Complex 0 -1)) (make-Complex 0 1))
(check-expect (conjugate (make-Complex 1 -1)) (make-Complex 1 1))
(check-expect (conjugate (make-Complex -1 1)) (make-Complex -1 -1))
```

Data Type → Template → Template Instantiation → Code

Instantiation of Complex template for magnitude?

```
;; Template instantiation
;; (define (conjugate c)
;;      ... (Complex-real c) ...
;;      ... (Complex-imag c) ...)
```

- This template instantiation is trivial but more complex examples are not.
- Templates help us write the code

```
 (define (conjugate c)
    (make-Complex (Complex-real c) (- (Complex-imag c)))
```

- Sophisticated types -> sophisticated templates …
  helping us write correct, sophisticated code

Union (Mixed) Data Definitions

How can we define data types that include more than one kind of data?

- Use the notion of *disjoint* set union from mathematics
- Example:

```
;; A shape is either:
;;   a square (make-square s) with side s,
;;   an equilateral triangle (make-triangle s) with
;;     side s, or
;;   a circle (make-circle s) with diameter s,
;; where s is a number and square, triangle, and circle
;; are structs defined as follows.
;; (define square (size))
;; (define triangle (size))
;; (define circle (size))
```

- This data definition can be abbreviated as follows:

```
;; shape ::= (make-square s) | (make-triangle s) |
;;           (make-circle s)
;; where s is a number and square triange, and circle...
```

Template for Union (Mixed) Data

For the type defined on the previous slide, the general template is:

```
; (define (f ... ashape ...)
;   (cond
;    [(square? ashape) ... (square-size ashape) ...]      ;; square case
;    [(triangle? ashape) ... (triangle-size ashape) ...] ;; triangle case
;    [(circle? ashape) ... (circle-size ashape) ...]))    ;; circle case
```

Processing mixed data requires a conditional to direct control to the appropriate case code. Note that **cond** is critical because it directs evaluation to the appropriate code, ignoring irrelevant clauses.

The template for an arbitrary union (assuming each construction is unary)

```
;; mixed-type ::= (make-S1 field) | ... | (make-SN field)
```

and struct definitions for ***S1*, . . . , *SN***, the general template for processing data of this type is:

```
; (define (f ... amt ...)
;   (cond
;    [(S1? amt) ... (S1-field amt) ... ]    ;; S1 case
;     ...                                    ;; cases 2, ..., N-1
;    [(SN? amt) ... (SN-field amt) ... ])) ;; SN case)
```

Inductive Data Definitions

How can we generate arbitrarily large data objects like lists?

- Use self-reference (induction/recursion) in typing fields
- Example:

```
;; A list-of-numbers is either
;;   empty, or
;;   (cons n lon)
;; where n is a number and lon is a list-of-
numbers
```

- If we assume that **`empty`** is a built-in (primitive) constant (like **`true`**), this definition can be implemented in Scheme by the *struct*

```
(define-struct cons (first rest))
```

where **`make-`** is elided from the constructor

Inductive Data Definitions

The **`cons`** *struct* definition is built-in to Scheme; it is *primitive*. For the sake of brevity, the constructor is simply called **`cons`** rather than **`make-cons`** and the accessors are called **`first`** and **`rest`** rather than **`cons-first`** and **`cons-rest`**. Note that a Scheme *struct* definition does not stipulate the types of the fields of the structure. Hence, the programmer is responsible for ensuring that **`cons`** is used correctly. In our dialects of Scheme, **`cons`** ensures that its second argument is a list. (In the standard dialect, **`first`** is called **`car`** and **`rest`** is called **`cdr`** for historical reasons.)

Template for Inductive Data Type

```
;; (define (f ... alon ...)
;;  (cond
;;    [(empty? alon ) ...]                 ;; empty case
;;    [(cons? alon ) ... (first alon) ... ;; cons case
;;       ... (f ... (rest alon) ...) ...]))
```

- Processing inductive (self-referential) data requires recursion (self-reference) in the computation.
- What is **cond** ?
- This template for processing inductive data is an extension of the on Slide 8 for processing is a degenerate form of this template he previous slide where there are multiple clauses (varieties) but no self-reference. The template is identical except for absence of the recursive call.

Extended Example: Insertion Sort

- Problem: given a list-of-numbers, sort it into ascending (non-decreasing) order.
- The solution that we will develop is the sample solution in the Scheme HW Guide.

If Expressions

- Simplified notation for common conditional expressions.
- Form:

  (if *question result-1 result-2*)

  abbreviates:

```
(cond [question  result-1]
      [else      result-2])
```

- Hence,

```
(if true result-1 result-2) => result-1
(if false result-1 result-2) => result-2
```

Epilog

- Reminder: work on HW01. Over the weekend, you should be able to complete problems from Section 8.3 and make substantial progress on the programs. They all process lists.
- Next class: data-directed design using other inductive types

Example of help function

```
; Type contract
; magnitude: Complex -> number
; Purpose: (magnitude c) computes the magnitude of the Complex number c, i.e.
; the L2 norm of (x,y) where c = (make-Complex x y)
; Examples
 (check-expect (magnitude (make-Complex 0 0)) 0)
 (check-expect (magnitude (make-Complex 3 4)) 5)
 (check-expect (magnitude (make-Complex 4 3)) 5)
; Code
(define (magnitude c) (2norm (Complex-real c) (Complex-imag c))
 ; Type contract
 ; 2norm: number number -> number
 ; Purpose: (2norm x y) returns the L2 norm of the vector (x,y), i.e.
 ;    (sqrt (+ (* x x) (* y y)))
 ; Examples:
 (check-expect (2norm 0 0) 0)
 (check-expect (2norm 3 4) 5)
 (check-expect (2norm 4 3) 5)
 ; Template intsantion: trival
 ; Code:
 (define (2norm x y) (sqrt (+ (* x x) (* y y))))
```