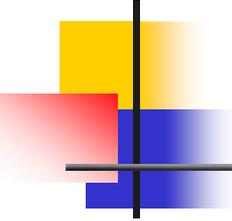


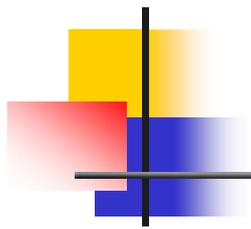
Functional Abstraction and Polymorphism

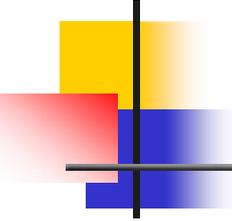
Corky Cartwright
Department of Computer Science
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(with thanks to John Greiner)



Abstracting Designs

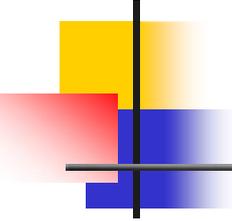
- The elimination of repetitions is the most important step in the (program) editing process – HTDP
- Software engineering term for revising a program to make it better or accommodate an extension: *refactoring*.
- Repeated code should be avoided at almost all costs. Why? Revisions involving repeated code are almost impossible to get right.
- *Abstractions* help us avoid this problem.





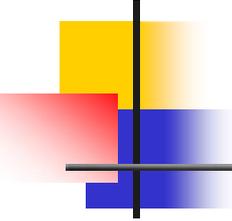
The Need for Abstractions

```
;; contains-doll? : los -> boolean
;; to determine whether alos contains
;; the symbol 'doll
(define (contains-doll? alos)
  (cond
    [(empty? alos) false]
    [else (or (symbol=? (first alos) 'doll)
              (contains-doll? (rest alos)))]))
```



The Need for Abstractions

```
;; contains-car? : los -> boolean
;; to determine whether alos contains
;; the symbol 'car
(define (contains-car? alos)
  (cond
    [(empty? alos) false]
    [else (or (symbol=? (first alos) 'car)
              (contains-car? (rest alos)))]))
```

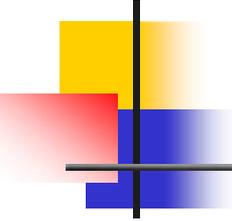


Creating Abstractions

How can we write one function that replaces

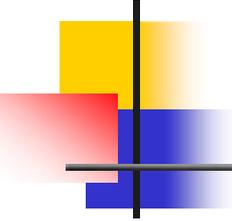
- `contains-doll?`
- `contains-car?`
- `contains-pizza?`
- `contains-number?`
- ...

As we design code that embodies an abstraction (such as containing a particular object), performing explicit type checks **blocks** abstraction. Explicit type checking expressed as executable code is a terrible software engineering idea! Why? The type information to should be confirmed in a given invocation depends on the context of the invocation!



Creating Abstractions

```
;; contains? : symbol, los -> boolean
;; to determine whether alos contains
;; the symbol s
(define (contains? s alos)
  (cond
    [(empty? alos) false]
    [else (or (symbol=? (first alos) s)
              (contains? s (rest alos)))]))
```



Creating Abstractions, cont.

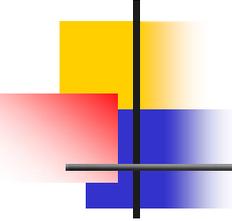
```
;; contains? : any list -> boolean
;; (contains? v alist) determines whether
;; alist contains the value v
(define (contains? v alist)
  (cond
    [(empty? alist) false]
    [else (or (equals? (first alist) v)
              (contains? v (rest alist)))]))
```

Using generic (parametric typing), the correct type contract is:

```
;; contains? : (list-of T) -> boolean
```

Note that the “scope” of the type parameter T is the entire type expression

```
(list-of T) -> boolean
```



Using Abstractions

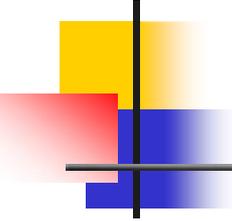
- How do we use `contains`?

```
(contains? 'doll (list ...))  
(contains? 'car (list ...))
```

- How can we better define `contains-doll?`, `contains-car`?

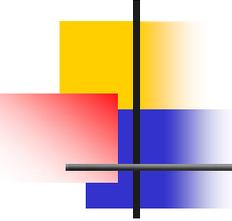
```
(define (contains-doll? alos) (contains? 'doll alos))  
(define (contains-car? alos) (contains? 'car alos))
```

- This idea is called *reuse*. Let's run with it!



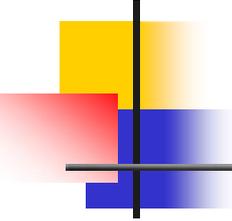
A more complex example

```
;; below : lon number -> lon
;; contract: to construct a list of those numbers
;; in alon that are less than or equal to t
(define (below alon t)
  (cond [(empty? alon) empty]
        [else
         (cond [(<= (first alon) t)
                (cons (first alon)
                      (below (rest alon) t))]
               [else (below (rest alon) t)]))]))
```



A more complex example ...

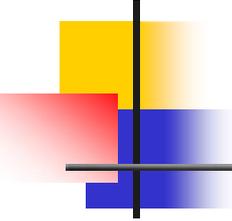
```
;; above : lon number -> lon
;; contract: to construct a list of those numbers
;; in alon that are greater than n
(define (above alon n)
  (cond [(empty? alon) empty]
        [else
         (cond [(> (first alon) n)
                  (cons (first alon)
                        (above (rest alon) n))]
               [else (above (rest alon) n)]))]))
```



Creating Abstractions

How can we write one function that replaces the functions

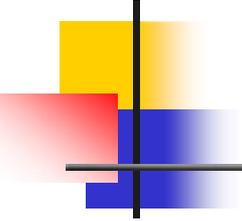
- `below`
- `above`
- `equal`
- `same-sign-as`
- ...



Creating Functional Abstractions

```
;; filter1: (number number -> boolean) (list-of number) number -> lon
;; contract: (filter1 test alon n) returns a list of those numbers e in
;; alon such that (test e n) is true
(define (filter1 test alon n)
  (cond [(empty? alon) empty]
        [else
         (cond [(test (first alon) n)
                 (cons (first alon)
                       (filter1 test (rest alon) n))]
               [else (filter1 test (rest alon) n)]))]))
```

What did we do? Use a function as an argument!

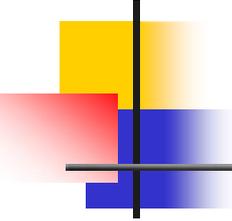


Using Abstractions

- How do we denote (express) function *values*? Two ways: names and **lambda**-abstractions. We will only use the first way for now: write the *name* of a defined function (primitive, library, or program-defined) as in:

```
(filter1 <= (list ...) 0)  
(filter1 > (list ...) 0)
```
- How can we define the functions above, below without code duplication?

```
(define (below alon n) (filter1 <= alon n))  
(define (above alon n) (filter1 > alon n))
```
- Both functions will work just as before!
- Can we do better? The **filter1** example is warped by assumptions that the input is a (list-of number) and that the abstracted function is the partial application (to the value of the argument *n*) of a binary relational operator. The standard **filter** operation that takes a unary predicate (equivalent to the partial application). This approach is simpler, but it requires **lambda**-notation to express the function arguments in our examples!



Repetition in Types

Repetition also happens in type definitions.

A `lon` is one of:

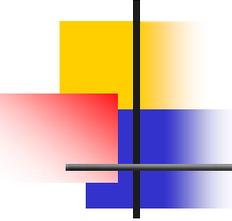
- `empty`
- `(cons n alon)`

where `n` is a number and `alon` is a `lon`.

A `los` is one of:

- `empty`
- `(cons s alos)`

where `s` is a `symbol` and `alos` is a `los`.



Abstracting Types

A `(list-of T)` is one of:

- `empty`
- `(cons t alot)`

where `t` is a `T` and `alot` is an `(list-of T)`.

A variable at the type level.

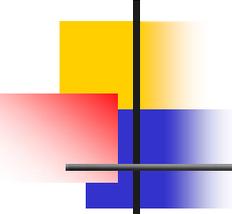
In FP, called **parametric polymorphism**

In OOP, called **genericity (generic types)**

Abstracting Types

Type	Example(s)
<code>(list-of number)</code>	<code>(list 1 2 3)</code>
<code>(list-of symbol)</code>	<code>(list 'a 'b 'pizza)</code>
<code>(list-of any)</code>	<code>(list 1 2 3)</code> <code>(list 'a 'b 'pizza)</code> empty <code>(list 1 'a +)</code>

Important! `(list-of x)` is NOT the same type as `(list-of any)`. It is an *abstracted* type. In some contexts (when X is not bounded by a supertype), `(list-of any)` is a *parametric instantiation* of `(list-of X)`.



Improving `filter1`

Can we generalize the type of `filter1`?

```
;; filter1: (number number -> boolean) (list-of number) number ->  
;;           (list-of number
```

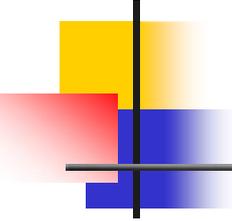
What is special about `number`? Does `filter1` rely on any of the properties of `number`?

No. It could be any type `X`.

```
;; filter1 : (X X -> boolean) (list-of X) -> (list-of X)
```

Comment: `filter1` is still lame. It should be unary:

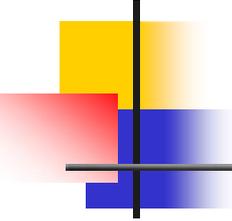
```
;; filter : (X -> boolean) (list-of X) -> (list-of X)
```



A Better Filter Function

```
;; filter: (T -> boolean) (list-of T) -> (list-of T)
;; contract: (filter f alot) constructs the (list-of T) consisting of all
;; elements e in alot such that (f e) is true. The ordering of elements
;; is preserved
(check-expect (filter number? empty) empty)
(check-expect (filter number? (list 1 false 2 true 3)) (list 1 2 3))

(define (filter f alot)
  (cond [(empty? alot) empty]
        [else
         (cond [(f (first alot))
                  (cons (first alot)
                        (filter f (rest alot)))]
               [else (filter f (rest alot))])]))
```



Final thoughts

- Function abstraction adds *expressiveness* to a programming language
- Type abstraction (polymorphism) does the same for type annotations.
- They work well together, e.g. the ML family (OCAML, Haskell).
- Core Racket does not have static type system but we still use types (homogeneous subsets of the data domain) informally in stating contracts. A type declaration is a restricted form of contract!
- Function abstraction is very lightweight in Racket and other functional languages. In Java, it is rather clumsy and heavy-weight (both in notation and implementation cost), but still important; single inheritance is not a general mechanism for expressing function abstraction but it is notationally simpler in the cases where it is applicable. Since Java 8, Java supports a clean form of multiple inheritance via interfaces with concrete methods (code!). The `default` method nomenclature is unfortunate.