
COMP 322: Fundamentals of Parallel Programming

Lecture 27: InterruptedException, Advanced Locking

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Solution to Worksheet #26: Java Threads

1) Write a sketch of the pseudocode for a Java threads program that exhibits a data race using `start()` and `join()` operations.

```
1. // Start of thread t0 (main program)
2. sum1 = 0; sum2 = 0; // Assume that sum1 & sum2 are fields
3. // Compute sum1 (lower half) and sum2 (upper half) in parallel
4. final int len = x.length;
5. Thread t1 = new Thread(() -> {
6.     for(int i=0 ; i < len/2 ; i++) sum1+=x[i];});
7. t1.start();
8. Thread t2 = new Thread(() -> {
9.     for(int i=len/2 ; i < len ; i++) sum2+=x[i];});
10. t2.start();
11. int sum = sum1 + sum2; // data race between t0 & t1, and t0 & t2
12. t1.join(); t2.join();
```



Solution to Worksheet #26: Java Threads (contd)

2) Write a sketch of the pseudocode for a Java threads program that exhibits a data race using synchronized statements.

```
1. // Start of thread t0 (main program)
2. sum = 0; // static int field
3. Object a = new ... ;
4. Object b = new ... ;
5. Thread t1 = new Thread(() -> { synchronized(a) { sum++; } });
6. Thread t2 = new Thread(() -> { synchronized(b) { sum++; } });
1. t1.start();
7. t2.start(); // data race between t1 & t2
8. t1.join(); t2.join();
```



Objects and Locks in Java --- synchronized statements and methods (Recap)

- Every Java object has an associated lock acquired via:
 - **synchronized statements**
 - `synchronized(foo) { // acquire lock on object foo
// execute code while holding foo's lock
} // release lock on object foo`
 - **synchronized methods**
 - `public synchronized void op1() { // acquire lock on "this" object
// execute method while holding 'this' lock
} // release lock on "this" object`
- Java language does not enforce any relationship between object used for locking and objects accessed in isolated code
 - If same object is used for locking and data access, then the object behaves like a monitor
- Locking and unlocking are **automatic**
 - Locks are released when a synchronized block exits
 - By normal means: end of block reached, `return`, `break`
 - When an exception is thrown and not caught



Use of class objects in synchronized statements/methods

- A **class** object exists for every class
- **static synchronized** methods lock the **class** object
- **class** object can be locked explicitly:
 - `synchronized(Foo.class) { /* ... */ }`
- No connection between locking the **Class** object and locking an instance of the class
 - Locking the **Class** object **does not** lock any instance
 - Instance methods that use static variables must synchronize access to them explicitly by locking the **Class** object

Always use the class literal to get reference to **Class** object—
not `this.getClass()` as you may access a subclass object



Cancelling Threads: Interruption

- Problem: how do we shut down a thread like a web server?
- Need to communicate that shutdown has been requested
 - Could set a flag that is polled in the main loop
 - But main loop could be blocked waiting for a request
- Interruption provides a means of signalling a request to another thread
- Each **Thread** has an “interrupted status” which is
 - Set when **interrupt()** method is invoked on it
 - Queried by **isInterrupted()** method
- Many blocking methods respect interruption requests and return early by throwing checked **InterruptedException**
 - **Object.wait()**
 - Throwing **IE** usually clears interrupted status



Calling methods that may throw InterruptedException

- Many methods in Java thread libraries may throw an InterruptedException e.g., <thread>.join(), <object>.wait(),
- When calling any such method, you will either need to include each call to join() in a try-catch block, or add a “throws InterruptedException” clause to the definition of the method that includes the call to join()
- Try-catch code for InterruptedException in Bounded Buffer example:

```
while (count == BUFFER SIZE) {  
    try {  
        wait();  
    }  
    catch (InterruptedException e) { }  
}
```



Responses to Interruption

- **Re-throw IE**
 - So caller can handle interruption request
- **Cancel and return early**
 - Clean up and exit without signalling an error
 - May require rollback or recovery
- **Ignore interruption**
 - When it is too dangerous to stop
 - Should re-assert interrupted status before returning
- **Postpone interruption**
 - Remember that interrupt occurred
 - Finish what you are doing and then throw IE
- **Throw a general failure exception**
 - When interruption is one of many reasons method can fail



Example: Shutting Down the Web Server

```
1. public class WebServerWithShutdown {
2.     private final ServerSocket server;
3.     private Thread serverThread;
4.     public WebServerWithShutdown(int port) throws IOException {
5.         server = new ServerSocket(port);
6.         server.setSoTimeout(5000); // so we can check for interruption
7.     }
8.     public synchronized void shutdownServer() throws IE.,IOException {
9.         if (serverThread == null) throw new IllegalStateException();
10.        serverThread.interrupt();
11.        serverThread.join(5000); // wait 5s before closing socket
12.        server.close(); // to give thread a chance to cleanup
13.    }
14.    public synchronized void startServer() {
15.        if (serverThread == null) {
16.            (serverThread = new Thread() {
17.                public void run() {
18.                    while (!Thread.interrupted()) {
19.                        try { processRequest(server.accept()); }
20.                        catch (SocketTimeoutException e) { continue; }
21.                        catch (IOException ex) { /* log it */ }
22.                    }
23.                }
24.            }).start();
25.        }
26.    }
27. }
```

Note: shutdownServer can be harmlessly called more than once



Locks and Conditions in java.util.concurrent library

- Atomic variables
 - The key to writing lock-free algorithms
- Concurrent Collections:
 - Queues, blocking queues, concurrent hash map, ...
 - Data structures designed for concurrent environments
- Locks and Conditions
 - More flexible synchronization control
 - Read/write locks
- Executors, Thread pools and Futures
 - Execution frameworks for asynchronous tasking
- Synchronizers: Semaphore, Latch, Barrier, Exchanger
 - Ready made tools for thread coordination



Locks

- Use of monitor synchronization is just fine for most applications, but it has some shortcomings
 - Single wait-set per lock
 - No way to interrupt or time-out when waiting for a lock
 - Locking must be block-structured
 - Inconvenient to acquire a variable number of locks at once
 - Advanced techniques, such as hand-over-hand locking, are not possible
- Lock objects address these limitations
 - But harder to use: Need `finally` block to ensure release
 - So if you don't need them, stick with **synchronized**

Example of hand-over-hand locking:

- `L1.lock() ... L2.lock() ... L1.unlock() ... L3.lock() ... L2.unlock() ...`



java.util.concurrent.locks.Lock interface

```
interface Lock {
    void lock();
    void lockInterruptibly() throws InterruptedException;
    boolean tryLock(); // return false if lock is not obtained
    boolean tryLock(long timeout, TimeUnit unit)
                    throws InterruptedException;
    void unlock();
    Condition newCondition();
    // can associate multiple condition vars with lock
}
```

- `java.util.concurrent.locks.Lock` interface is implemented by `java.util.concurrent.locks.ReentrantLock` class



Simple ReentrantLock() example

- Used extensively within `java.util.concurrent`

```
final Lock lock = new ReentrantLock();  
...  
lock.lock();  
try {  
    // perform operations protected by lock  
}  
catch(Exception ex) {  
    // restore invariants & rethrow  
}  
finally {  
    lock.unlock();  
}
```

- **Must manually ensure lock is released**



java.util.concurrent.locks.condition interface

- Can be allocated by calling `ReentrantLock.newCondition()`
- Supports multiple condition variables per lock
- Methods supported by an instance of condition
 - void `await()` // NOTE: not wait
 - Causes current thread to wait until it is signaled or interrupted
 - Variants available with support for interruption and timeout
 - void `signal()` // NOTE: not notify
 - Wakes up one thread waiting on this condition
 - void `signalAll()` // NOTE: not notifyAll()
 - Wakes up all threads waiting on this condition
- For additional details see
 - <http://download.oracle.com/javase/1.5.0/docs/api/java/util/concurrent/locks/Condition.html>



BoundedBuffer implementation using two conditions, notFull and notEmpty

```
1. class BoundedBuffer {
2.     final Lock lock = new ReentrantLock();
3.     final Condition notFull = lock.newCondition();
4.     final Condition notEmpty = lock.newCondition();

6.     final Object[] items = new Object[100];
7.     int putptr, takeptr, count;
8.
9.     . . .
```



BoundedBuffer implementation using two conditions, notFull and notEmpty (contd)

```
10. public void put(Object x) throws InterruptedException
11. {
12.     lock.lock();
13.     try {
14.         while (count == items.length) notFull.await();
15.         items[putptr] = x;
16.         if (++putptr == items.length) putptr = 0;
17.         ++count;
18.         notEmpty.signal();
19.     } finally {
20.         lock.unlock();
21.     }
22. }
```



BoundedBuffer implementation using two conditions, notFull and notEmpty (contd)

```
23.    public Object take() throws InterruptedException
24.    {
25.        lock.lock();
26.        try {
27.            while (count == 0) notEmpty.await();
28.            Object x = items[takeptr];
29.            if (++takeptr == items.length) takeptr = 0;
30.            --count;
31.            notFull.signal();
32.            return x;
33.        } finally {
34.            lock.unlock();
35.        }
36.    }
```



Reading vs. writing

- Recall that the use of synchronization is to protect interfering accesses
 - Multiple concurrent reads of same memory: Not a problem
 - Multiple concurrent writes of same memory: Problem
 - Multiple concurrent read & write of same memory: Problem

So far:

- If concurrent write/write or read/write might occur, use synchronization to ensure one-thread-at-a-time

But:

- This is unnecessarily conservative: we could still allow multiple simultaneous readers

Consider a hashtable with one coarse-grained lock

- So only one thread can perform operations at a time

But suppose:

- There are many simultaneous lookup operations
- insert operations are very rare



java.util.concurrent.locks.ReadWriteLock interface

```
interface ReadWriteLock {  
    Lock readLock ();  
    Lock writeLock ();  
}
```

- Even though the interface appears to just define a pair of locks, the semantics of the pair of locks is coupled as follows
 - Case 1: a thread has successfully acquired `writeLock().lock()`
 - No other thread can acquire `readLock()` or `writeLock()`
 - Case 2: no thread has acquired `writeLock().lock()`
 - Multiple threads can acquire `readLock()`
 - No other thread can acquire `writeLock()`
- `java.util.concurrent.locks.ReadWriteLock` interface is implemented by `java.util.concurrent.locks.ReadWriteReentrantLock` class



Example code

```
class Hashtable<K,V> {
    ...
    // coarse-grained, one lock for table
    ReadWriteLock lk = new ReentrantReadWriteLock();
    V lookup(K key) {
        int bucket = hasher(key);
        lk.readLock().lock(); // only blocks writers
        ... read array[bucket] ...
        lk.readLock().unlock();
    }
    void insert(K key, V val) {
        int bucket = hasher(key);
        lk.writeLock().lock(); // blocks readers and writers
        ... write array[bucket] ...
        lk.writeLock().unlock();
    }
}
```



Worksheet #27: use of tryLock()

Name: _____

Netid: _____

Extend the transferFunds() method from Lecture 26 (shown below) to use j.u.c. locks with tryLock() instead of synchronized, and to return a boolean value --- true if it succeeds in obtaining in obtaining both locks and performing the transfer, and false otherwise. Assume that each Account object contains a reference to a dedicated ReentrantLock object. Sketch your answer below using pseudocode. Can you create a deadlock with multiple calls to transferFunds() in parallel?

```
1. public void transferFunds(Account from, Account to, int amount)
   {
2.     synchronized (from) {
3.         synchronized (to) {
4.             from.subtractFromBalance(amount);
5.             to.addToBalance(amount);
6.         }
7.     }
8. }
```

