COMP 322: Fundamentals of Parallel Programming

Lecture 3: Computation Graphs, Ideal Parallelism

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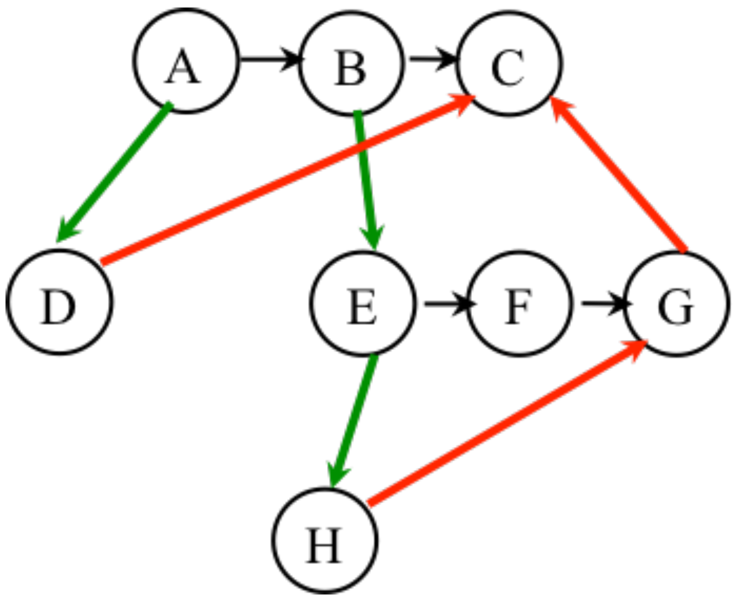
One Possible Solution to Worksheet 2 (Reverse Engineering a Computation Graph)



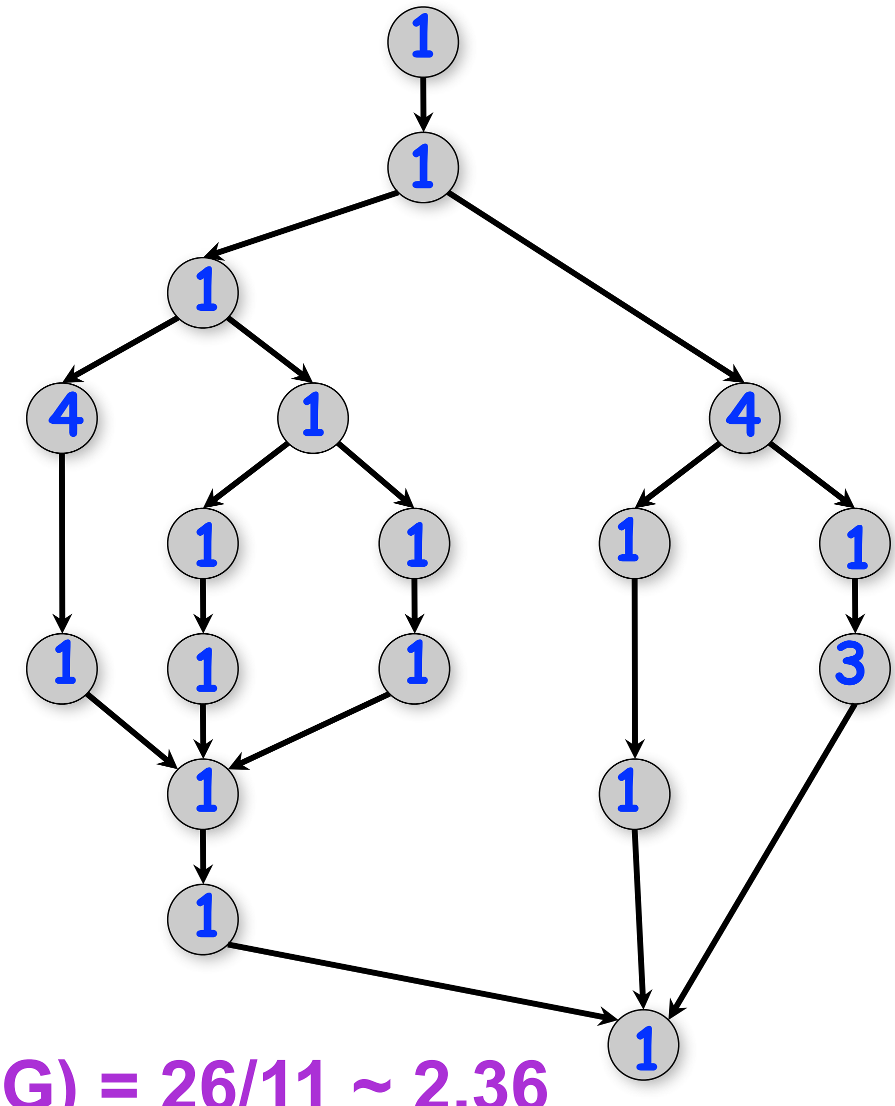

```
A();
finish { // F1
  async D();
  B();
  {
    E();
    finish { // F2
      async H();
       F();
    } // F2
    G();
  }
} // F1
C();
```

Observations:

- Any node with out-degree > 1 must be an async (must have an outgoing spawn edge)
- Any node with in-degree > 1 must be an end-finish (must have an incoming join edge
- Adding or removing transitive edges does not impact ordering constraints

Ideal Parallelism (Recap)

- Define ideal parallelism of Computation G Graph as the ratio, WORK(G)/CPL(G)
- Ideal Parallelism only depends on the computation graph, and is the speedup that you can obtain with an unbounded number of processors



Example:

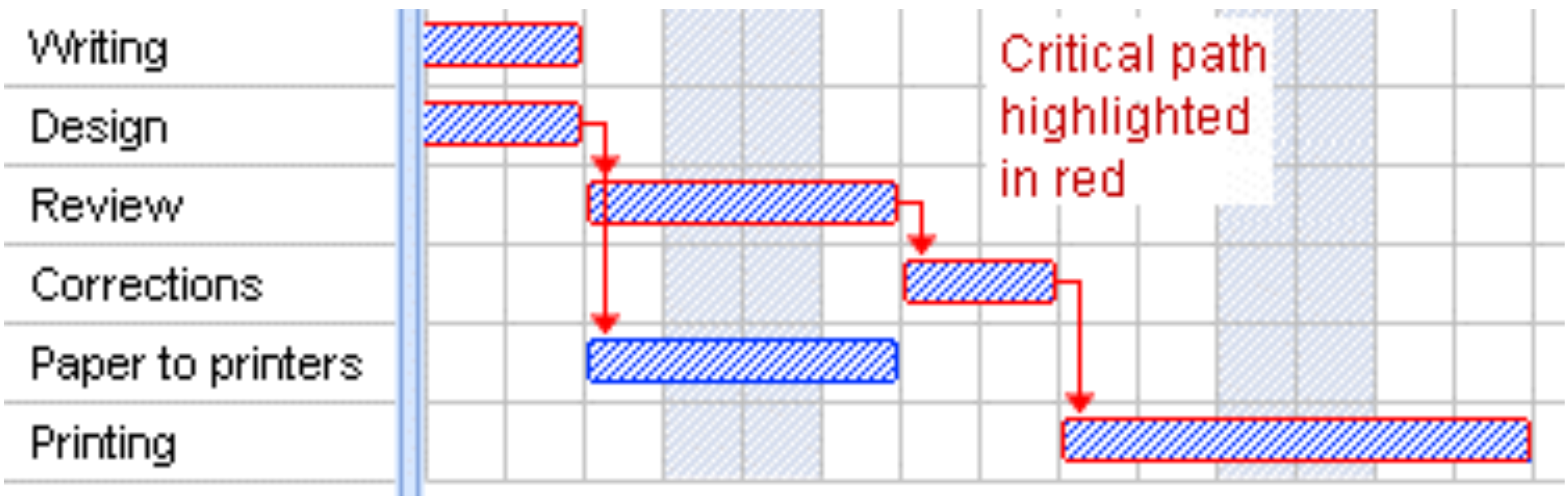
WORK(G) = 26

CPL(G) = 11

Ideal Parallelism = WORK(G)/CPL(G) = 26/11 ~ 2.36

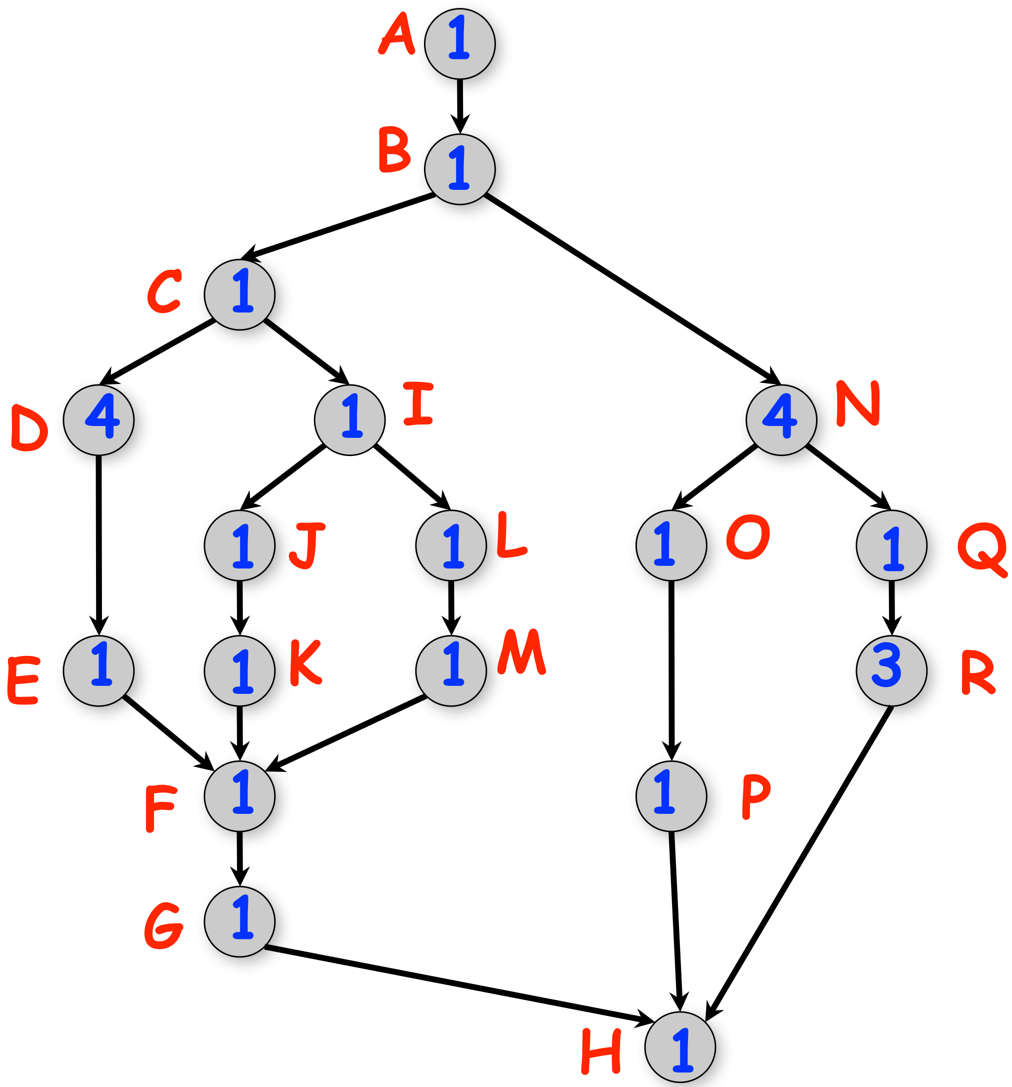
Computation Graphs are used in Project Scheduling as well

- Computation graphs are referred to as “Gantt charts” in project management
- Sample project for preparing a printed document
  - —Source: http://www.gantt.com/creating-gantt-charts.htm



Scheduling of a Computation Graph on a fixed number of processors: Example

node label = time(N), for all nodes N in the graph



Start time	Proc 1	Proc 2	Proc 3
0	A		
1	B		
2	C	N	
3	D	N	I
4	D	N	J
5	D	N	K
6	D	Q	L
7	E	R	M
8	F	R	O
9	G	R	P
10	H		
11	Completion time = 11		

NOTE: this schedule achieved a completion time of 11. Can we do better?

Scheduling of a Computation Graph on a fixed number of processors, P

- Assume that node N takes TIME(N) regardless of which processor it executes on, and that there is no overhead for creating parallel tasks
- A schedule specifies the following for each node
  - —START(N) = start time
  - —PROC(N) = index of processor in range 1...P

  such that
  - —START(i) + TIME(i) <= START(j), for all CG edges from i to j (Precedence constraint)
  - —A node occupies consecutive time slots in a processor (Non-preemption constraint)
  - —All nodes assigned to the same processor occupy distinct time slots (Resource constraint)

Greedy Schedule

- A greedy schedule is one that never forces a processor to be idle when one or more nodes are ready for execution
- A node is ready for execution if all its predecessors have been executed
- Observations
  - T_1 = WORK(G), for all greedy schedules
  - T_∞ = CPL(G), for all greedy schedules
- where T_P = execution time of a schedule for computation graph G on P processors

Lower Bounds on Execution Time of Schedules

- Let T_P = execution time of a schedule for computation graph G on P processors
  - Can be different for different schedules
- Lower bounds for all greedy schedules
  - Capacity bound: $T_P \geq WORK(G)/P$
  - Critical path bound: $T_P \geq CPL(G)$
- Putting them together
  - $T_P \geq \max(WORK(G)/P, CPL(G))$

Upper Bound on Execution Time of Greedy Schedules

Theorem [Graham '66]. Any greedy scheduler achieves

$$T_P \leq WORK(G)/P + CPL(G)$$

<u>Proof sketch</u>:

Define a time step to be complete if $\geq P$ nodes are ready at that time, or incomplete otherwise

complete time steps $\leq WORK(G)/P$

incomplete time steps $\leq CPL(G)$

Start time	Proc 1	Proc 2	Proc 3
0	A		
1	B		
2	C	N	
3	D	N	I
4	D	N	J
5	D	N	K
6	D	Q	L
7	E	R	M
8	F	R	O
9	G	R	P
10	H		
11			

Bounding the performance of Greedy Schedulers

Combine lower and upper bounds to get

$\max(\text{WORK}(G)/P, \text{CPL}(G)) \le T_P \le \text{WORK}(G)/P + \text{CPL}(G)$

Corollary 1: Any greedy scheduler achieves execution time T_P that is within a factor of 2 of the optimal time (since max(a,b) and (a+b) are within a factor of 2 of each other, for any $a \ge 0, b \ge 0$).

Corollary 2: Lower and upper bounds approach the same value whenever

- There's lots of parallelism, WORK(G)/CPL(G) >> P
- Or there's little parallelism, WORK(G)/CPL(G) << P

Abstract Performance Metrics

- Basic Idea
  - Count operations of interest, as in big-O analysis
  - Abstraction ignores many overheads that occur on real systems
- Calls to doWork()
  - Programmer inserts calls of the form, doWork(N), within a step to indicate abstraction execution of N application-specific abstract operation
    - e.g., in the Homework 1 programming assignment (Parallel Sort), we will include one call to doWork(1) in each call to compareTo(), and ignore the cost of everything else.
- Abstract metrics are enabled by calling
  - HjSystemProperty.abstractMetrics.set(true);
- If an HJ program is executed with this option, abstract metrics are printed at end of program execution with WORK(G), CPL(G), Ideal Parallelism = WORK(G) / CPL(G)

Reminders

- Send email to comp322-staff@rice.edu if you do not have access to Piazza site (otherwise use Piazza for class communications, as far as possible)
- Office hours today will be held during 2pm - 3pm in Duncan Hall 3092
- Watch videos and read handout for topic 1.5 for next lecture on Wednesday, Jan 20th
- Complete this week's assigned quizzes on edX by 11:59pm today (all quizzes for topics 1.1, 1.2, 1.3, 1.4 including last quiz titled "Multiprocessor Scheduling")
- HW1 will be assigned today, and is due by 12noon on Jan 29th
- See course web site for all work assignments and due dates
  - http://comp322.rice.edu