

COMP 322: Fundamentals of Parallel Programming

Lecture 18: Abstract vs Real Performance - An “under the hood” look at HJlib

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HJ-lib Compilation and Execution Environment

Java 8 IDE

Foo.java

`javac Foo.java`

Java compiler

Foo.class

`java Foo`

*HJ-lib Runtime Environment =
Java Runtime Environment +
HJ-lib libraries*

HJ-lib Program Output

HJ-lib source program is a standard Java 8 program

Java compiler translates *Foo.java* to *Foo.class*, along with calls to HJ-lib with lambda parameters (*async*, *finish*, *future*, etc)

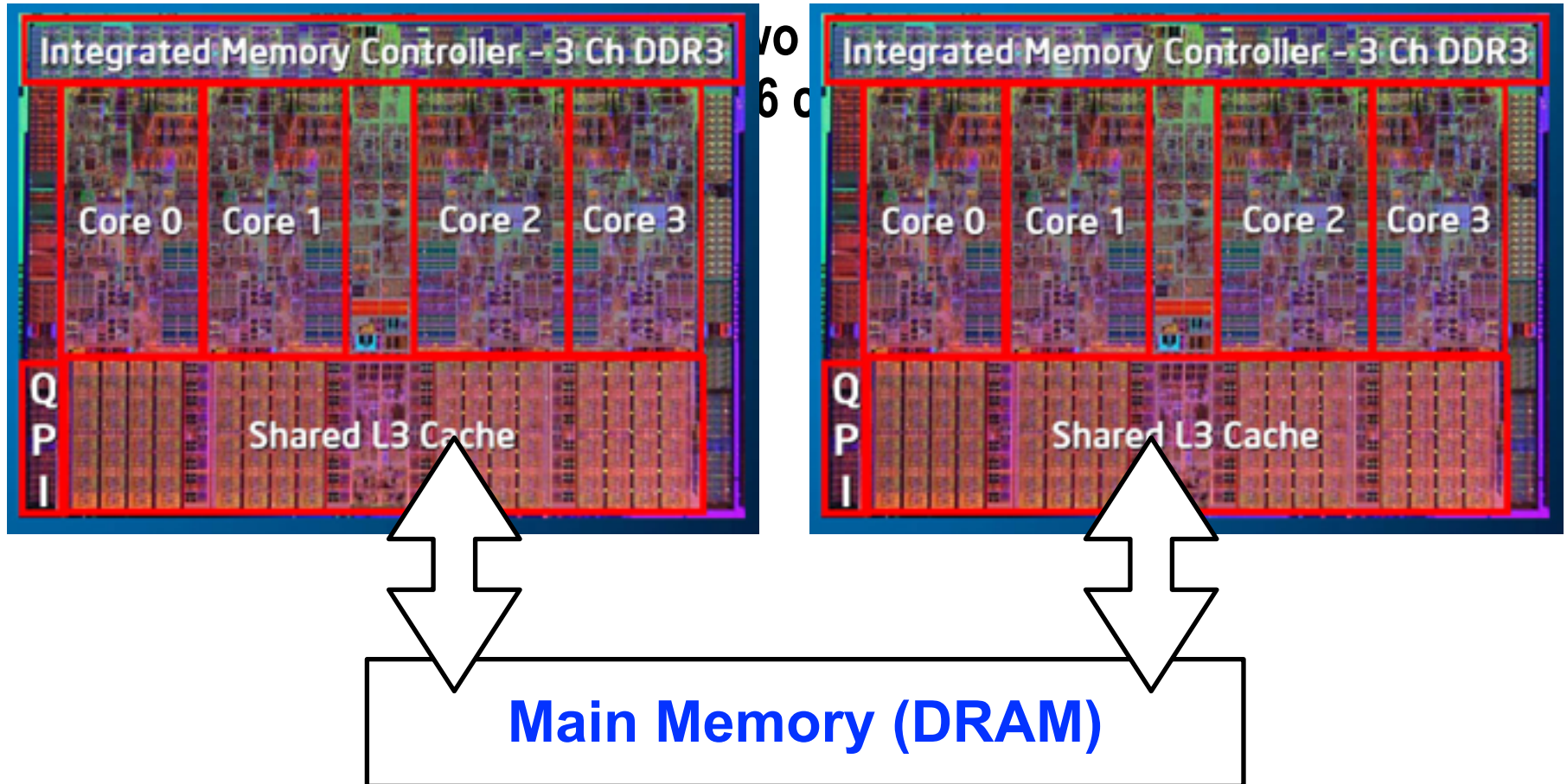
All the "magic" happens here!

HJ runtime initializes *m* worker threads
(value of *m* depends on options or default value)

*HJ Abstract Performance Metrics,
HJ-Viz output
(all enabled by appropriate options)*



Looking under the hood — let's start with the hardware

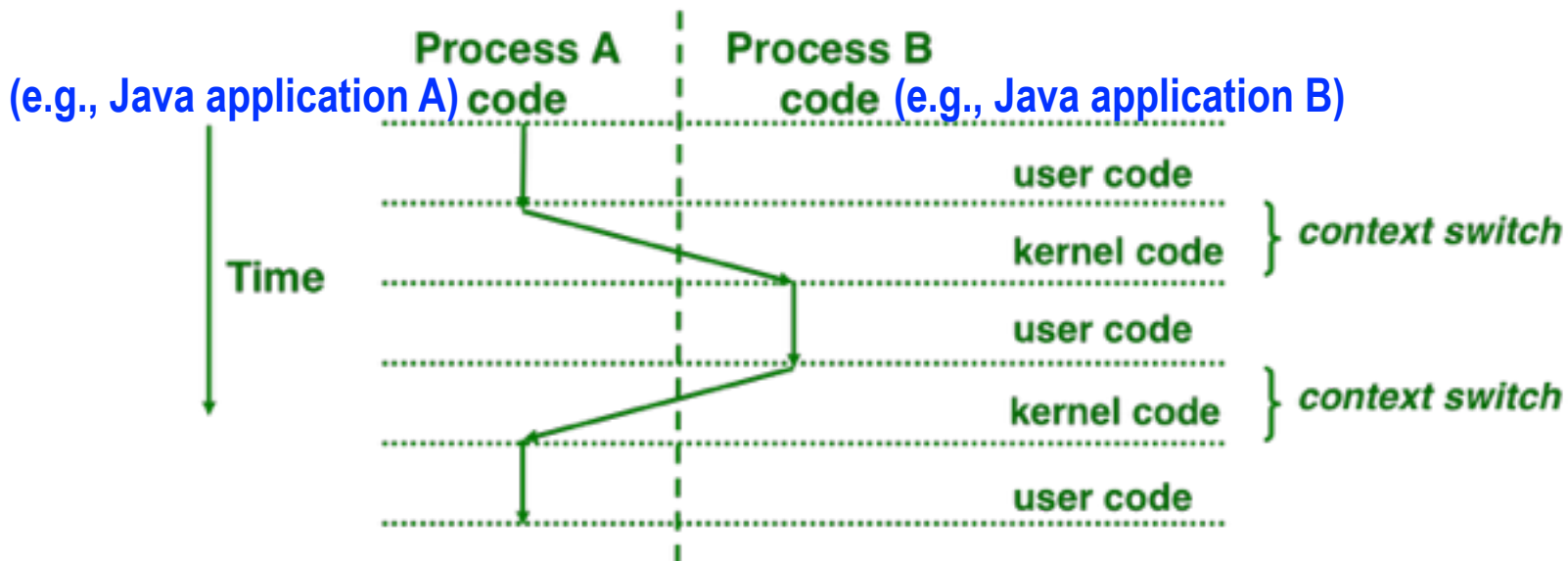


Next, how does a process run on a single core?

Processes are managed by OS kernel

- Important: the kernel is not a separate process, but rather runs as part of some user process

Control flow passes from one process to another via a context switch



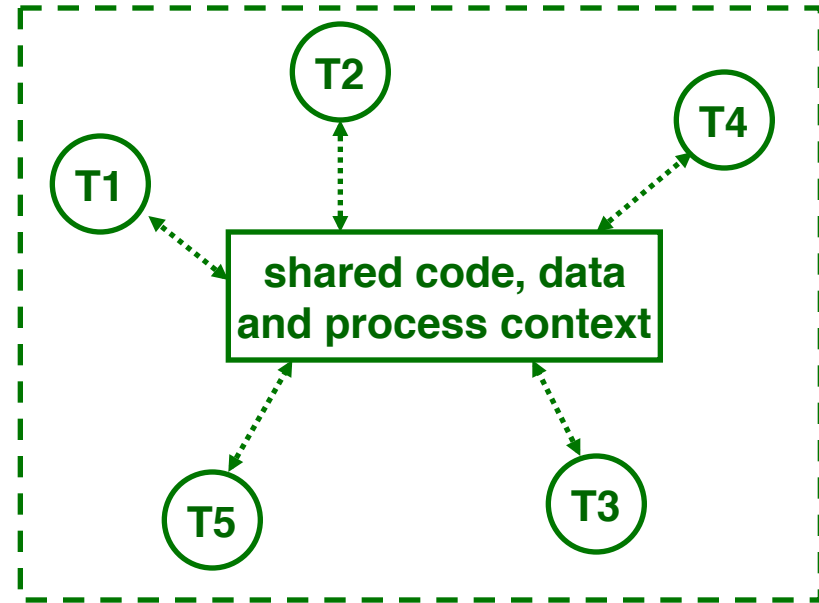
Context switches between two processes can be very expensive!

Source: COMP 321 lecture on Exceptional Control Flow (Alan Cox)



What happens when executing a Java program?

- A Java program executes in a single Java Virtual Machine (JVM) process with multiple threads
- Threads associated with a single process can share the same data
- Java main program starts with a single thread (T1), but can create additional threads (T2, T3, T4, T5) via library calls
- Java threads may execute concurrently on different cores, or may be context-switched on the same core

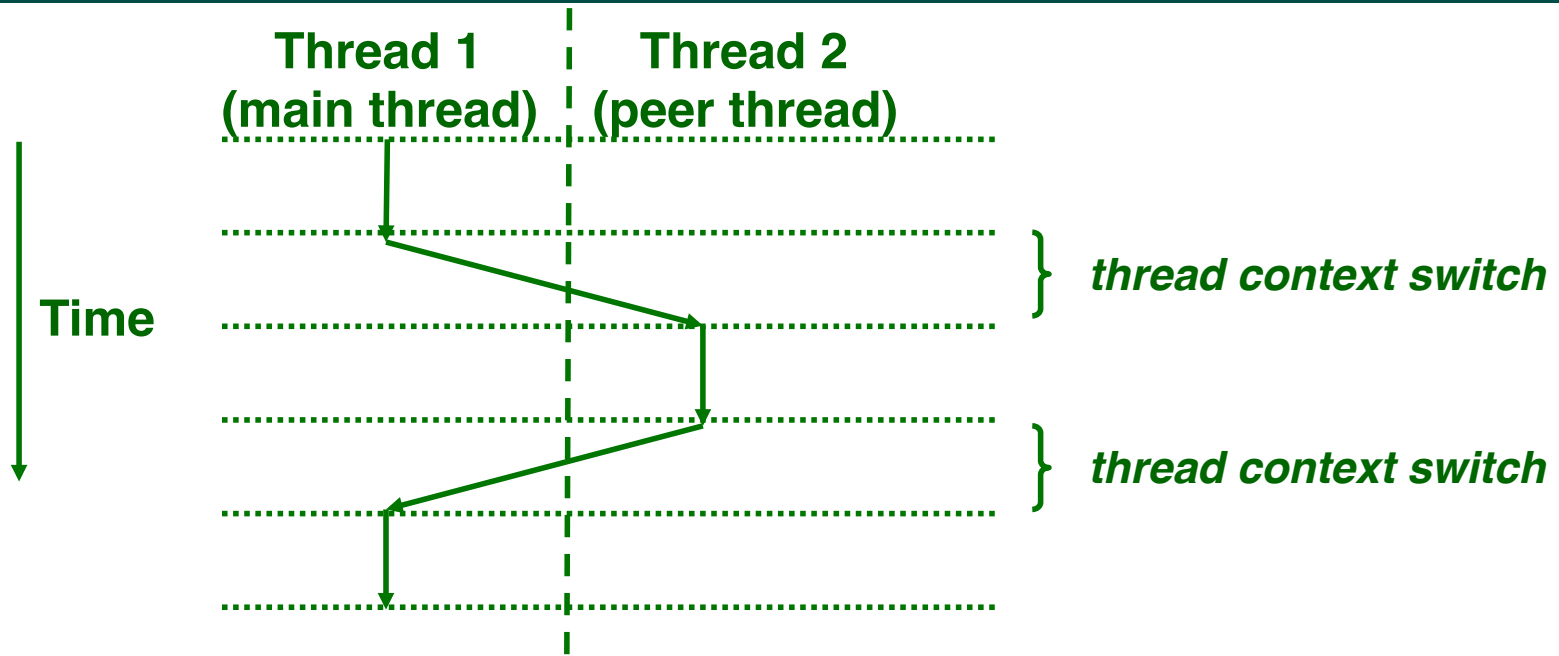


Java application with five threads — T1, T2, T3, T4, T5 — all of which can access a common set of shared objects

Figure source: COMP 321 lecture on Concurrency (Alan Cox)



Thread-level Context Switching on the same processor core



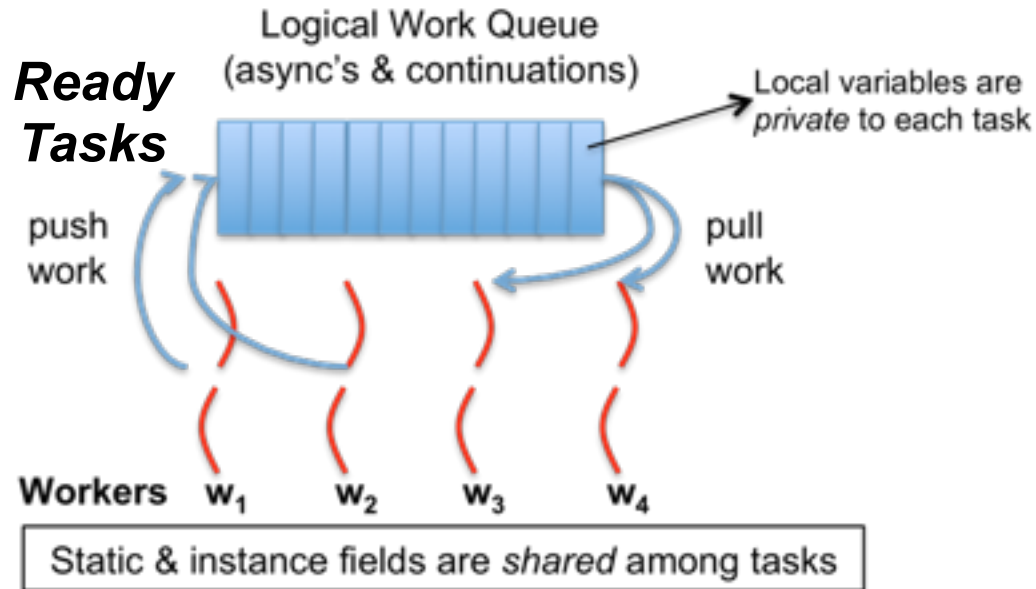
- Thread context switch is cheaper than a process context switch, but is still expensive (just not “very” expensive!)
- It would be ideal to just execute one thread per core (or hardware thread context) to avoid context switches

Figure source: COMP 321 lecture on Concurrency (Alan Cox)



Now, what happens in a task-parallel Java program (e.g., HJ-lib, Java ForkJoin, etc)

HJ-Lib Tasks & Continuations
Worker threads
Operating System
Hardware cores



- HJ-lib runtime creates a *small number of worker threads*, typically one per core
- Workers push new tasks and “continuations” into a logical work queue
- Workers pull task/continuation work items from logical work queue when they are idle (remember greedy scheduling?)



Task-Parallel Model: Checkout Counter Analogy



- Think of each checkout counter as a processor core

Image sources: <http://www.deviantart.com/art/Randomness-20-178737664>,
<http://www.wholefoodsmarket.com/blog/whole-story/new-haight-ashbury-store>

Task-Parallel Model: Checkout Counter Analogy



- Think of each checkout counter as a processor core
- And of customers as tasks

source: <http://www.deviantart.com/art/Randomness-20-178737664>

All is well until a task blocks ...



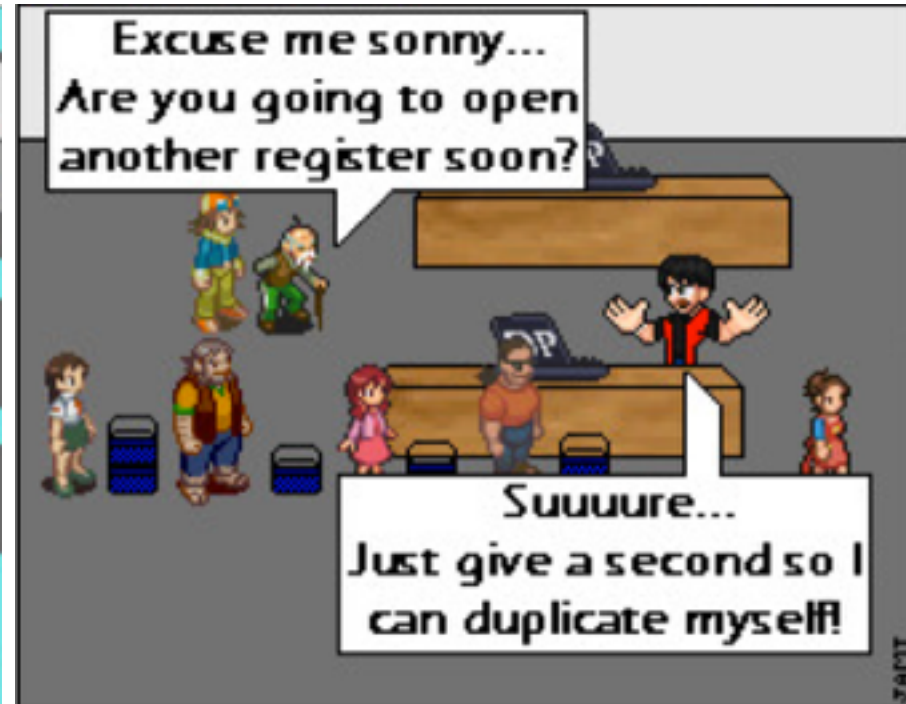
...



- A blocked task/customer can hold up the entire line
- What happens if each checkout counter has a blocked customer?

source: <http://viper-x27.deviantart.com/art/Checkout-Lane-Guest-Comic-161795346>

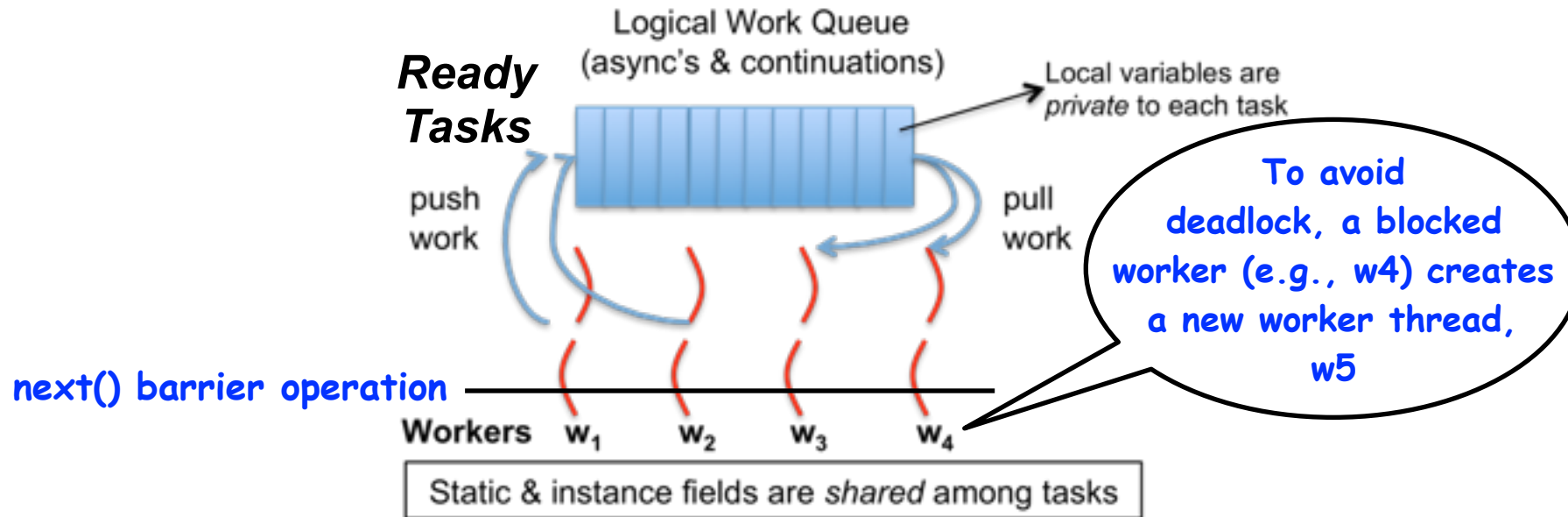
Approach 1: Create more worker threads (as in HJ-Lib's Blocking Runtime)



- Creating too many worker threads can exhaust system resources (OutOfMemoryError), and also leads to context-switch overheads when blocked worker threads get unblocked

source: <http://www.deviantart.com/art/Randomness-5-90424754>

Blocking Runtime (contd)



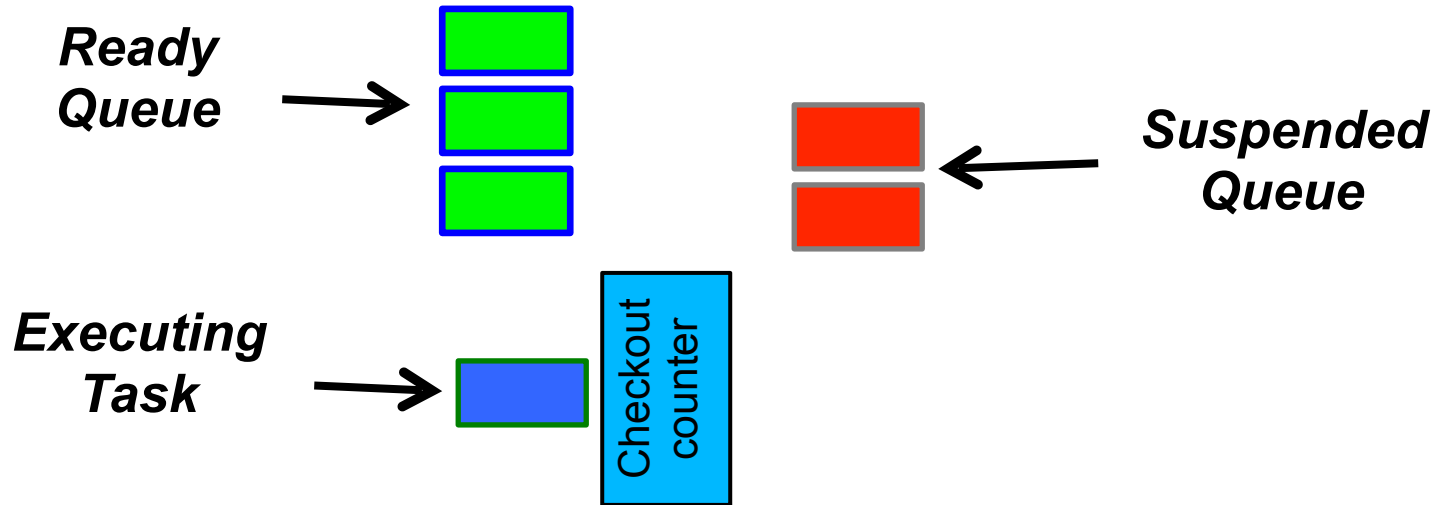
- Assume that five tasks (A1 ... A5) are registered on a barrier
- Q: What happens if four tasks (say, A1 ... A4) executing on workers w1 ... w4 all block at the same barrier?
- A: Deadlock! (All four tasks will wait for task A5 to enter the barrier.)
- Blocking Runtime's solution to avoid deadlock: keep task blocked on worker thread, and create a new worker thread when task blocks

Blocking Runtime (contd)

- Examples of blocking operations
 - End of finish
 - Future get
 - Barrier next
- Approach: Block underlying worker thread when task performs a blocking operation, and launch an additional worker thread
- Too many blocking operations can result in exceptions and/or poor performance, e.g.,
 - `java.lang.IllegalStateException: Error in executing blocked code! [89 blocked threads]`
- Maximum number of worker threads can be configured if needed
 - `HjSystemProperty.maxThreads.set(100);`



Approach 2: Suspend task continuations at blocking points (as in HJ-Lib's Cooperative Runtime)



- Upon a blocking operation, the currently executing tasks suspends itself and yields control back to the worker
- Task's *continuation* is stored in the suspended queue and added back into the ready queue when it is unblocked
- Pro: No overhead of creating additional worker threads
- Con: Need to create continuations (enabled by `-javaagent` option)

Continuations

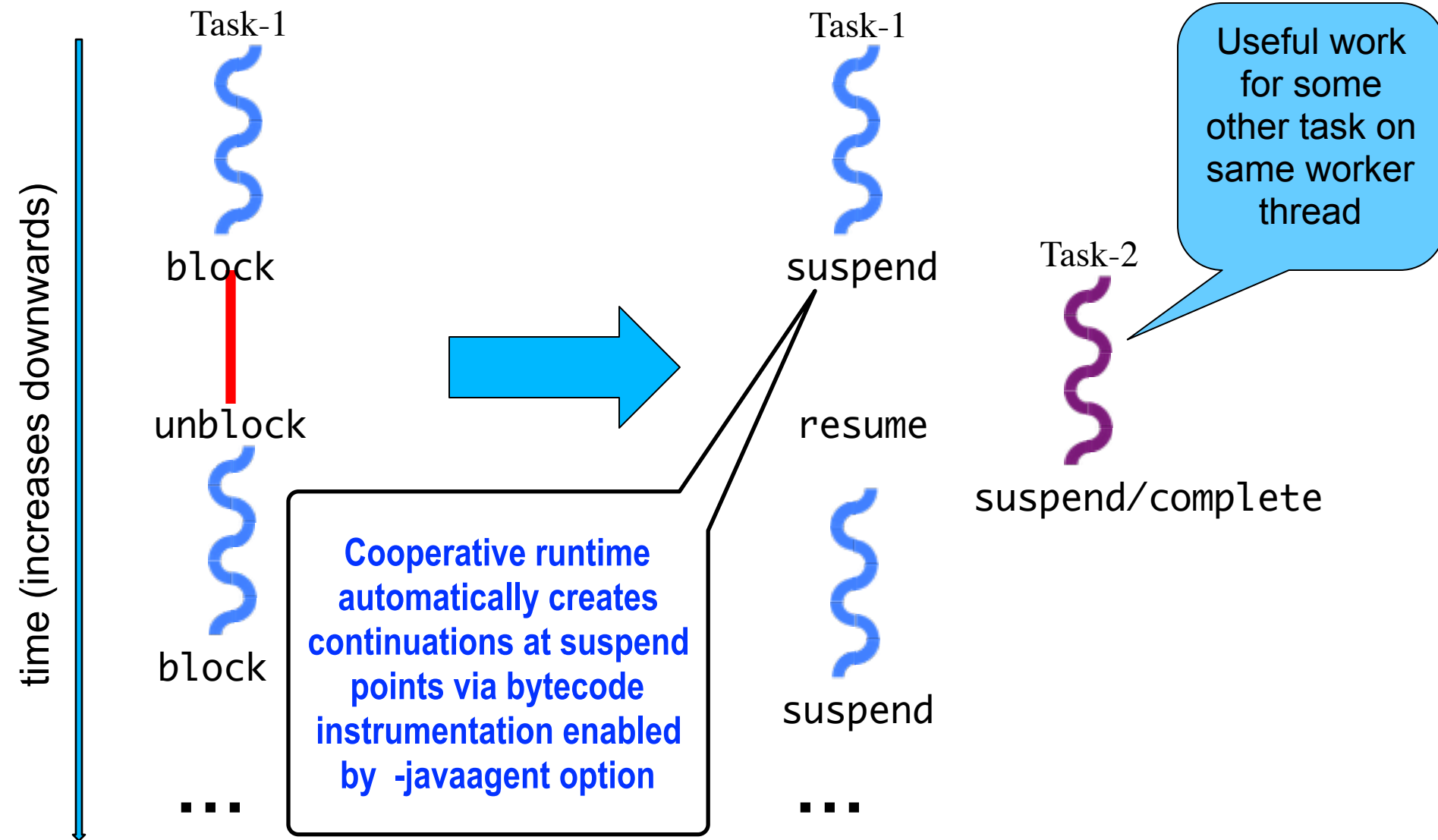
- A continuation is the point immediately following a *blocking* operation, such as an `end-finish`, future `get()`, barrier/phaser `next()`, etc.
- Continuations are also referred to as task-switching points
 - Program points at which a worker may switch execution between different tasks (depends on scheduling policy)

```
1. finish { // F1
2.  async A1;
3.  finish { // F2
4.    async A3;
5.    async A4;
6.  }
7.  S5;
8. }
```

Continuations

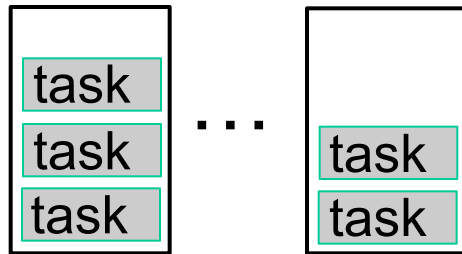


Cooperative Scheduling (view from a single worker)



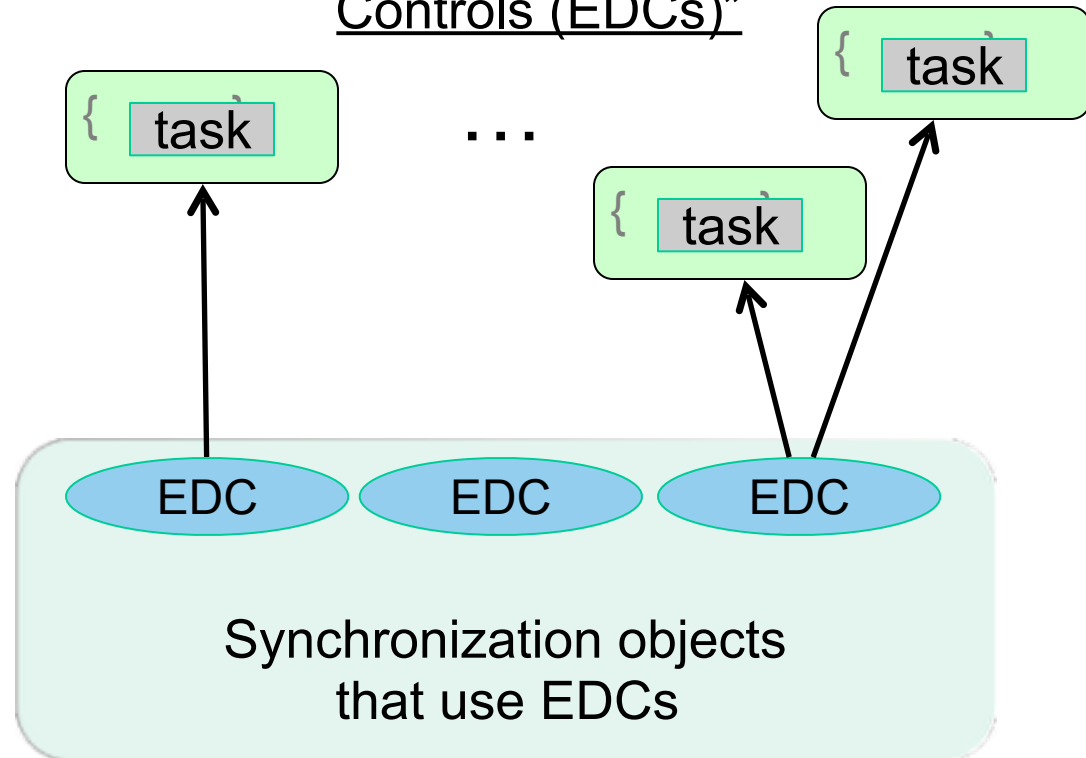
HJ-lib's Cooperative Runtime (contd)

Ready/Resumed Task Queues



Worker Threads

Suspended Tasks
registered with “Event-Driven
Controls (EDCs)””



Any operation that contributes to unblocking a task can be viewed as an event e.g., task termination in finish, return from a future, signal on barrier, put on a data-driven-future, ...



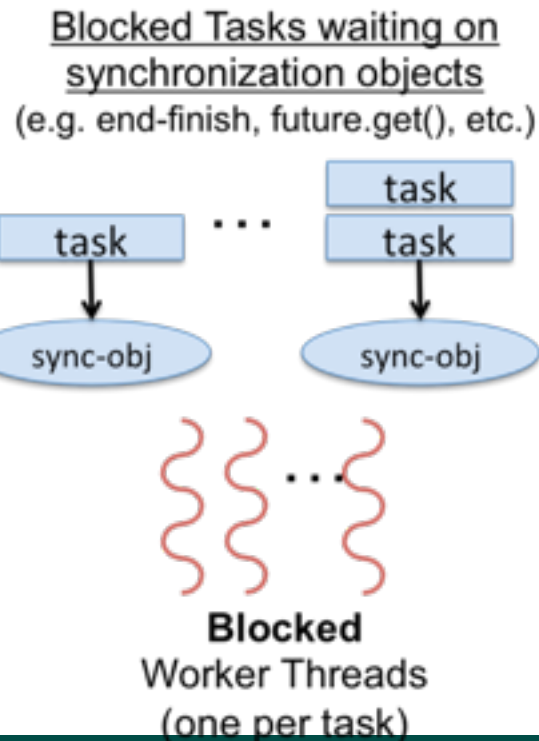
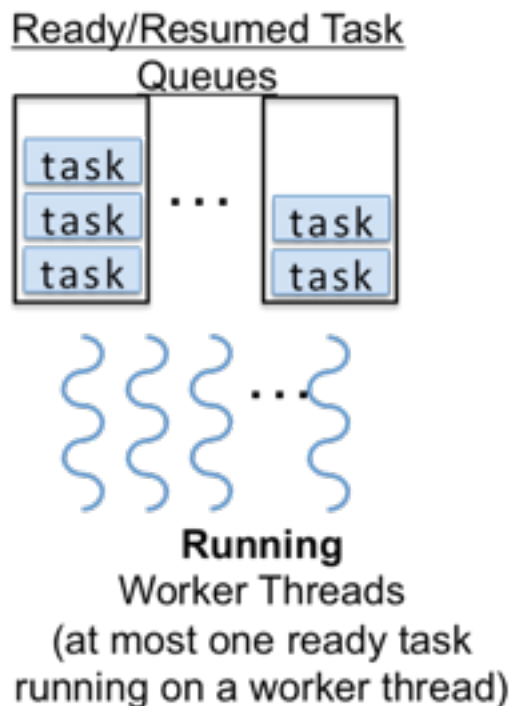
Why are Data-Driven Tasks (DDTs) more efficient than Futures?

- **Consumer task blocks on `get()` for each future that it reads, whereas `async-await` does not start execution till all Data-Driven Futures (DDFs) are available**
 - **An “`asyncAwait`” statement does not block the worker, unlike a `future.get()`**
 - **No need to create a continuation for `asyncAwait`; a data-driven task is directly placed on the `Suspended` queue by default**
- **Therefore, DDTs can be executed on a `Blocking Runtime` without the need to create additional worker threads, or on a `Cooperative Runtime` without the need to create continuations**



Summary: Abstract vs. Real Performance in HJlib

- **Abstract Performance**
 - Abstract metrics focus on operation counts for WORK and CPL, regardless of actual execution time
- **Real Performance**
 - HJlib uses ForkJoinPool implementation of Java Executor interface with **Blocking or Cooperative Runtime (option-controlled)**



Announcements & Reminders

- **HW3 CP 1 is available and due today by 11:59pm**
- **Watch the topic 4.1, 4.4 videos for the next lecture**
- **Use Piazza (public or private posts, as appropriate) for all communications re. COMP 322**
- **See Office Hours link on course web site for latest office hours schedule.**



Computation graph for async-finish program in Worksheet 18

