
COMP 322: Fundamentals of Parallel Programming

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Lecture 15: Point-to-point Synchronization, Pipeline Parallelism, Phasers (contd)

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Announcements

- Homework 4 due by 5pm on Wednesday, Feb 16th
 - We will try and return graded homeworks by Feb 23rd
- Guest lecture on Bitonic Sort by John Mellor-Crummey on Friday, Feb 18th
- Feb 23rd lecture will be a Midterm Review
- No lecture on Friday, Feb 25th since midterm is due that day
 - Midterm will be a 2-hour take-home written exam
 - Closed-book, closed-notes, closed-computer
 - Will be given out at lecture on Wed, Feb 23rd
 - Must be handed in by 5pm on Friday, Feb 25th



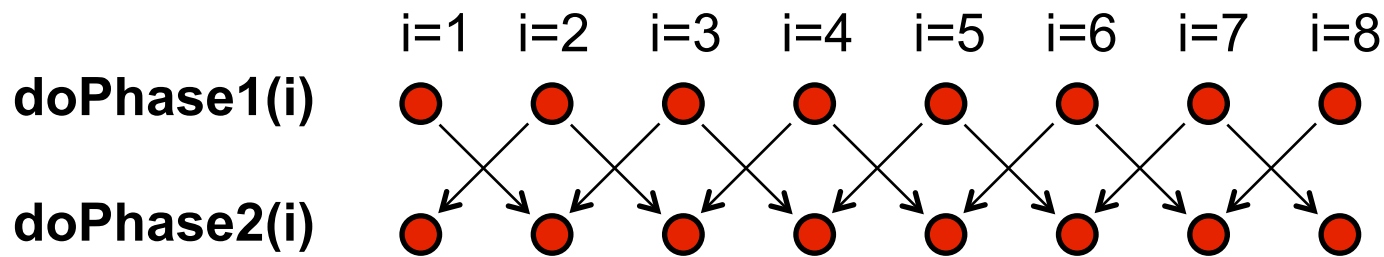
Acknowledgments for Today's Lecture

- [1] "X10: an object-oriented approach to non-uniform computing". Philippe Charles et al. OOPSLA 2005.
- [5] Phasers: a unified deadlock-free construct for collective and point-to-point synchronization. Jun Shirako et al. ICS '08
- Handout for Lectures 14 and 15



Point-to-Point Synchronization: Example 1 (Left-Right Neighbor Synchronization)

1. **finish** { // Expanded finish-for-async version of forall
 2. **for** (point[i] : [1:m])
 3. **async** {
 4. doPhase1(i);
 // Iteration i waits for i-1 and i+1 to complete Phase 1
 5. doPhase2(i);
 6. }
 7. }
- Need synchronization where iteration i only waits for iterations i-1 and i+1 to complete their work in doPhase1() before it starts doPhase2(i)? (Less constrained than a barrier)



Summary of Phaser Construct

- Phaser allocation
 - `phaser ph = new phaser(mode);`
 - Phaser `ph` is allocated with registration mode
 - *Phaser lifetime is limited to scope of Immediately Enclosing Finish (IEF)*
- *Registration Modes*
 - `phaserMode.SIG`
 - `phaserMode.WAIT`
 - `phaserMode.SIG_WAIT`
 - `phaserMode.SIG_WAIT_SINGLE`
- Phaser registration
 - `async phased (ph1<mode1>, ph2<mode2>, ...) <stmt>`
 - *Spawned task is registered with `ph1` in `mode1`, `ph2` in `mode2`, ...*
 - *Child task's capabilities must be subset of parent's*
 - `async phased <stmt>` propagates all of parent's phaser registrations to child
- Synchronization
 - `next;`
 - *Advance each phaser that current task is registered on to its next phase*
 - *Semantics depends on registration mode*



Capability Hierarchy

$SIG_WAIT_SINGLE = \{ \text{signal}, \text{wait}, \text{single} \}$

$SIG_WAIT = \{ \text{signal}, \text{wait} \}$

$SIG = \{ \text{signal} \}$

$WAIT = \{ \text{wait} \}$

- At any point in time, a task can be registered in one of four modes with respect to a phaser: SIG_WAIT_SINGLE , SIG_WAIT , SIG , or $WAIT$. The mode defines the set of capabilities — signal, wait, single — that the task has with respect to the phaser. The subset relationship defines a natural hierarchy of the registration modes.



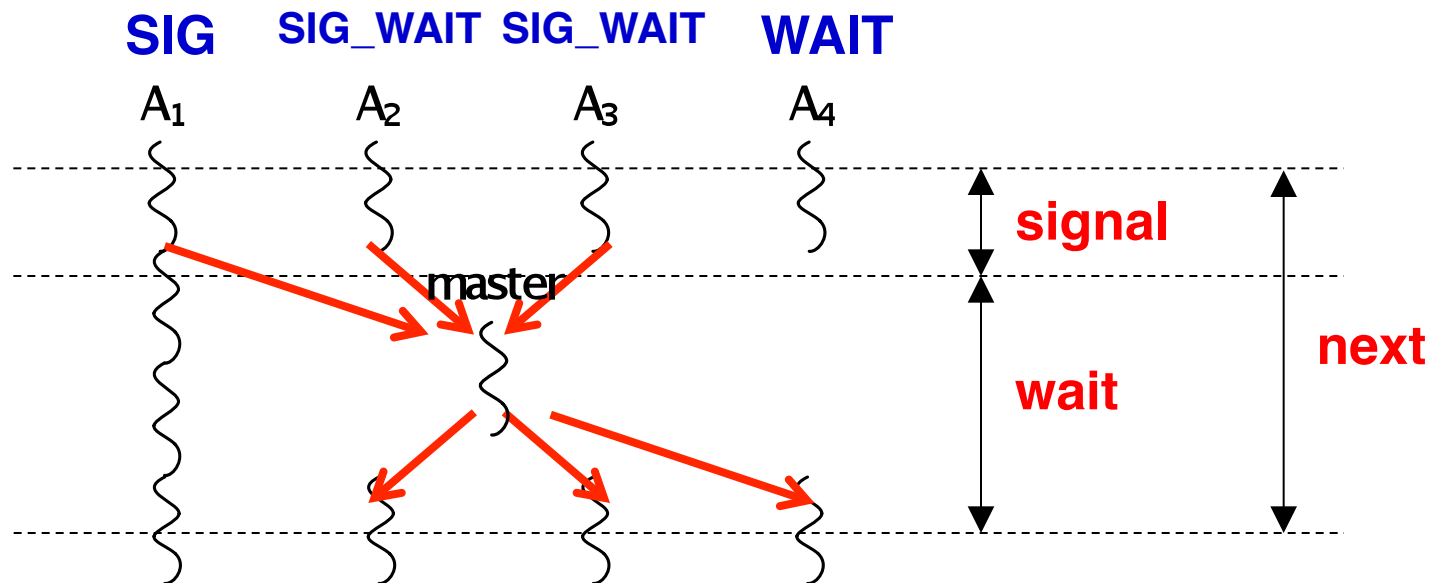
next operation

Semantics of **next** depends on registration mode

SIG_WAIT: next = signal + wait

SIG: next = signal (Don't wait for any task)

WAIT: next = wait (Don't disturb any task)



A master task **receives all signals and broadcasts a barrier completion**



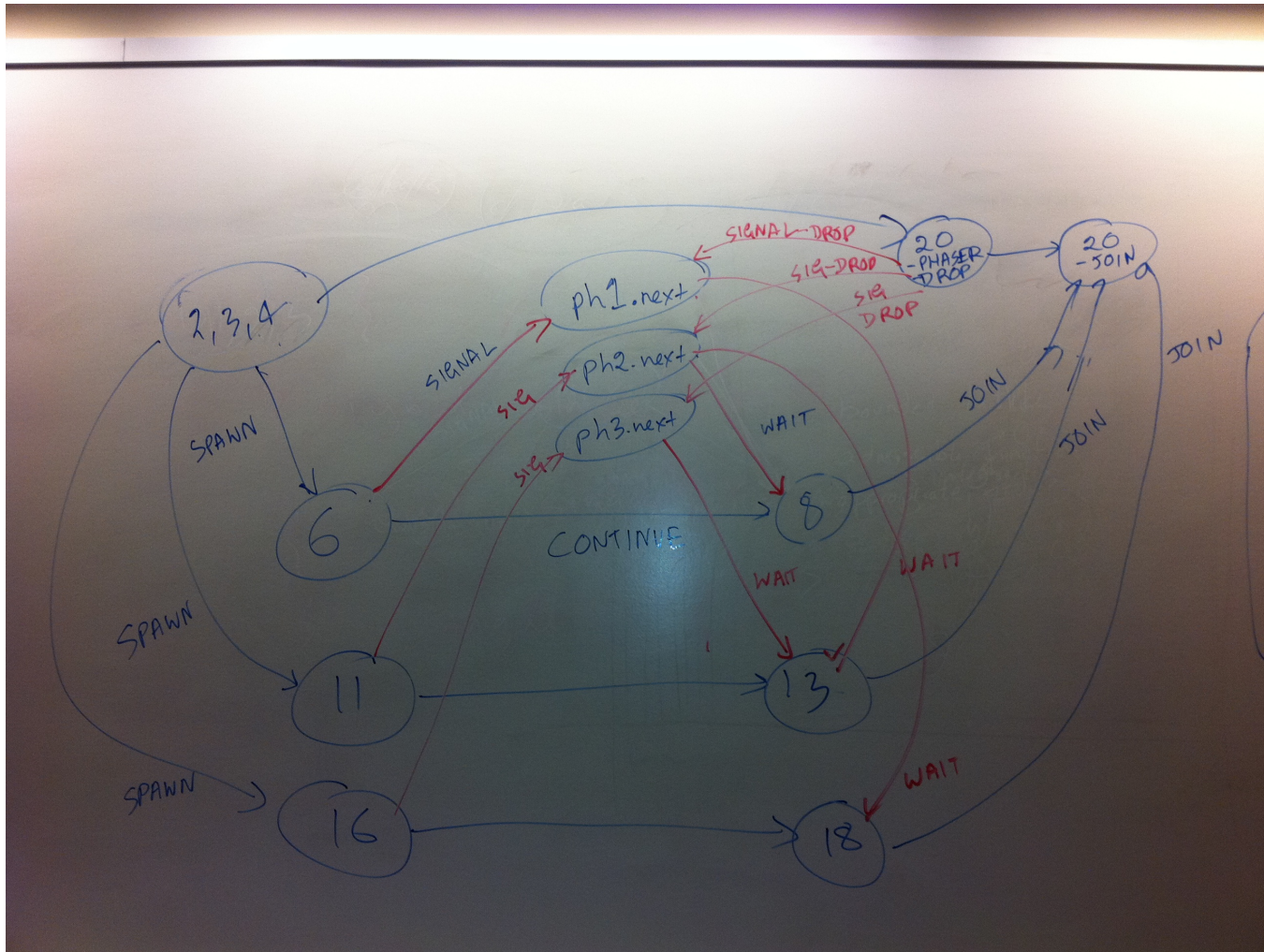
Left-Right Neighbor Synchronization Example for $m=3$ using Phasers

```
1  finish {
2    phaser ph1 = new phaser(); // Default mode is SIG_WAIT
3    phaser ph2 = new phaser(); // Default mode is SIG_WAIT
4    phaser ph3 = new phaser(); // Default mode is SIG_WAIT
5    async phased(ph1<SIG>, ph2<WAIT>) { // i = 1
6      doPhase1(1);
7      next; // Signals ph1, and waits on ph2
8      doPhase2(1);
9    }
10   async phased(ph2<SIG>, ph1<WAIT>, ph3<WAIT>) { // i = 2
11     doPhase1(2);
12     next; // Signals ph2, and waits on ph1 and ph3
13     doPhase2(2);
14   }
15   async phased(ph3<SIG>, ph2<WAIT>) { // i = 3
16     doPhase1(3);
17     next; // Signals ph3, and waits on ph2
18     doPhase2(3);
19   }
20 }
```

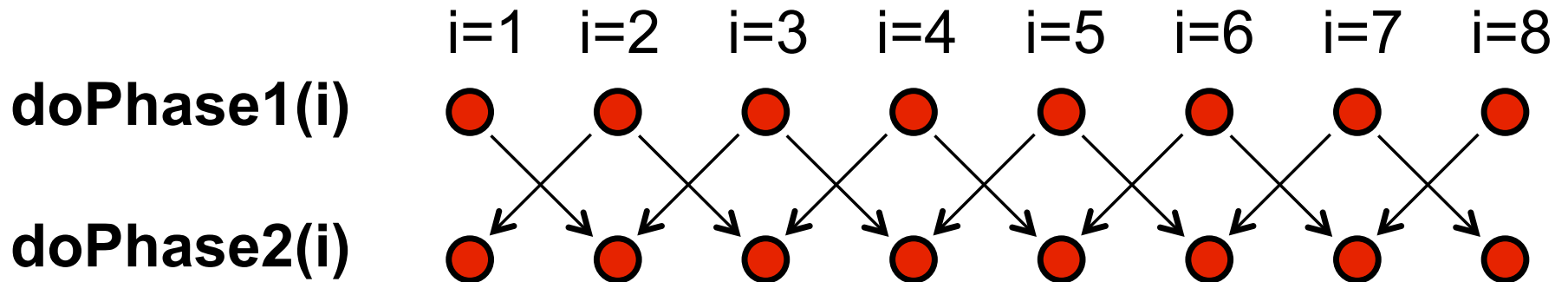
Listing 3: Extension of example in Listing 1 with three phasers for $m = 3$



Whiteboard picture from lecture (Computation Graph for previous slide)



Left-Right Neighbor Synchronization Example for General m



```
1 finish {
2   phaser ph = new phaser[m+2];
3   forall(point [i]:[0:m+1]) ph[i]=new phaser(); //Default mode is SIG_WAIT
4   for (point [i] : [1:m])
5     async phased(ph[i]<SIG>, ph[i-1]<WAIT>, ph[i+1]<WAIT>) {
6       doPhase1(i);
7       next; // Signals ph[i], and waits on ph[i-1] and ph[i+1]
8       doPhase2(i);
9     }
10 }
```

Listing 4: Extension of example in Listing 1 with array of $m + 2$ phasers for general m

